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Automatic Meshing of Free-Form Deformation Solids

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AUTOMATIC MESHING OF FREE-FORM DEFORMATION SOLIDS

A Dissertation
Presented to the
Department of Civil Engineering
Brigham Young University

In Partial Fulfillment
of the Requirement for the Degree
Doctor of Philosophy

by
Ssuta S. Hsu
April 1989

This dissertation, by Ssuta S. Hsu, is accepted in its present form by the Department of Civil Engineering of Brigham Young University as satisfying the dissertation requirement for the degree of Doctor of Philosophy.


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20 Mar 89
Date


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TABLE OF CONTENTS

ACKNOWLEDGEMENTS.....	iii
LIST OF FIGURES.....	vi
SECTION	
1.0 INTRODUCTION.....	1
2.0 BACKGROUND INFORMATION.....	4
2.1 Node generation approach.....	4
2.2 Mapping approach.....	5
2.3 Grid-based approach.....	6
2.4 Decomposition approach.....	7
2.5 Drag method.....	7
2.6 Automatic meshing of 3-D solids.....	8
2.7 Free-form deformation.....	10
2.8 Mesh smoothing.....	13
3.0 AUTOMATIC MESHING WITH HEXAHEDRON ELEMENTS.....	17
3.1 Inital cuboid primitive.....	17
3.2 Free-form deformation.....	18
3.3 Solid subdivision.....	19
3.3 Transition method.....	20
3.4 Smoothing procedure.....	22
4.0 EXAMPLES OF AUTOMATIC MESHING.....	24
4.1 Transition block.....	24
4.2 Hour glass model.....	26
4.3 Hammer head.....	28
5.0 CONCLUSION AND DISCUSSION.....	30
BIBLIOGRAPHY.....	31
APPENDICES.....	35
A. Results analysis and future development.....	36
B. Complete disscussion of algorithm.....	40
C. Continuity of elements.....	45
D. Bézier theory.....	49
E. Source code of free-form deformation mesh generator.....	59

LIST OF FIGURES

Figure

1.1	Free-form deformation grid model.....	2
1.2	Automatic Free-Form deformation model.....	3
2.1	Node generation approach.....	5
2.2	Mapping approach.....	5
2.3	Grid-based approach.....	6
2.4	Decomposition approach.....	7
2.5	Drag method model.....	8
2.6	Four hexahedra from a tetrahedron.....	9
2.7	Impossible for hexahedron element.....	10
2.8	Original plastic model with control points.....	12
2.9	Control points in deformed position.....	13
2.10	Deformed plastic model.....	13
2.11	Laplacian smoothing.....	15
2.12	Centroid area pull smoothing.....	16
3.1	Cuboid primitive.....	18
3.2	Deformed model.....	19
3.3	Transition case 1.....	21
3.4	Transition case 2.....	22
3.5	Transition case 3.....	22
3.6	Final result of finite element mesh.....	23
4.1	Original cuboid mesh.....	24
4.2	Control points and Deformation of transition block.	25
4.3	Uniform mesh of transition block.....	25
4.4	Smoothed mesh of transition block.....	26
4.5	Uneven element size of hour glass model.....	27

4.6	Near equal element size of hour glass model.....	27
4.7	Smoothed elements of hour glass model.....	28
4.8	Original mesh of hammer model.....	28
4.9	Mesh result of hammer model.....	29
A.1	Hammer model after smoothing.....	37
A.2	Finite element analysis result.....	39
C.1	Discontinue surface elements.....	46
C.2	Transition priorities.....	48
D.1	Center of mass of four points.....	50
D.2	Cubic Bézier Blending functions.....	51
D.3	Cubic Bézier curve.....	52
D.4	Examples of cubic Bézier curves.....	52
D.5	Subdivision of a cubic Bézier curve.....	54
D.6	Change the curve shape by the weights.....	55
D.7	Bicubic Bézier surface.....	58

1.0 INTRODUCTION

Computer-aided techniques are established as an effective way to approach modern engineering design and analysis. One aspect of Mechanical Computer-Aided Engineering (MCAE) consists of geometric modeling, mesh generation, finite element analysis, and result post-processing. Geometric modelers provide the capability to define the geometric model. Mesh generators create the finite element grid of the geometric model to be used for analysis. The finite element analysis software generates a solution to the boundary conditions applied to the geometric model. Post-processing provides the capability to display the analysis results using high performance graphics.

Mesh generation is the bridge between the geometry and finite element analysis procedure. Although the geometric and finite element analysis procedures are well developed, the process of discretizing a geometry into discrete numerical models for analysis is generally time consuming and tedious. There have been many programs written to generate finite element grids, but their applications are limited. Current research is directed toward automatic mesh generation. Such a procedure can produce a valid finite element mesh for an object with limited user input.

Solid free-form deformation geometric modeling, developed by Sederberg and Parry²⁹, is a very powerful tool in MCAE. It has the ability to create a geometric model by

sculpting surfaces with flexibility akin to clay in a sculptor's hand. In addition, any predefined grid can be modified by the free-form deformation process to create any shape defined by the user. Free-form deformation involves mapping by using a trivariate tensor product Bernstein polynomial²⁹. Figure 1.1 is an example of a model created by solid free-form deformation.

Often, the free-form deformation process distorts the initial elements so that they are unacceptable for finite element analysis. To correct this difficulty, a new method of mesh generation needs to be developed to create acceptable finite elements for the model.

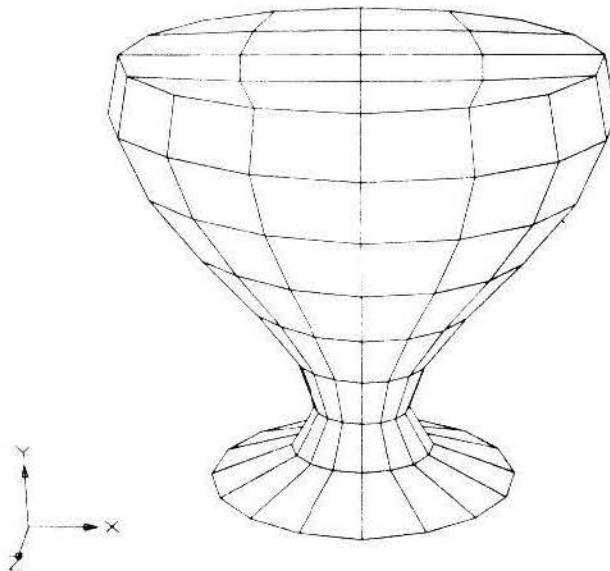


Figure 1.1 Free-form deformation grid model

The objective of this work is to develop a three-dimensional solid finite element mesh generator that will

automatically create acceptable hexahedral finite elements in conjunction with free-form deformation solid modeling. This automatic free-form deformation mesh generator combines the creation of the geometric model and automatic generation of hexahedral finite elements as shown in figure 1.2.

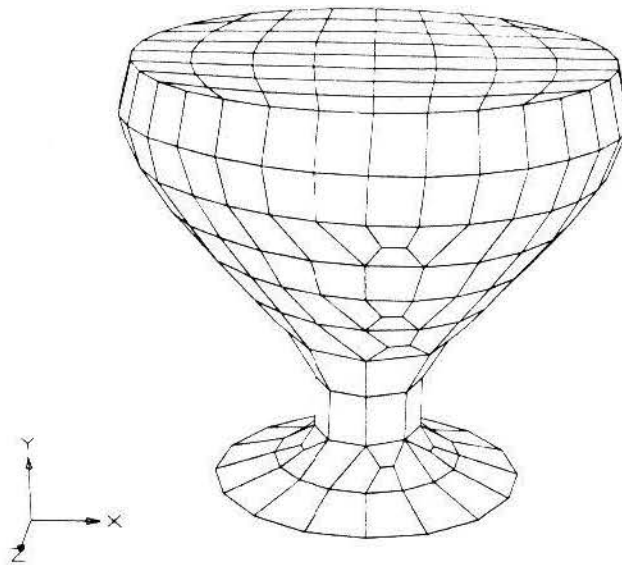


Figure 1.2 Automatic Free-form deformation model

Section two of this paper provides the background information on both the free-form deformation solid modeling and the automatic mesh generator. A description of the development of a new automatic mesh generator based on free-form deformation solid modeling is given in section three. Section four presents some examples which have been meshed by this new method and explains the user interaction required. Section five gives the conclusions and discussion.

2.0 BACKGROUND INFORMATION

Meshes were generated manually until in the early 1960s when computer assistance began to simplify the task of discretizing more complex objects. For more than twenty years, researchers have been developing methods to reduce the time and cost necessary to construct finite element models by computer assistance.

Buell and Bush¹ reviewed mesh generation procedures prior to 1972. Those methods, though rudimentary, have become the foundation of some current meshers. From 1972 to 1980 mesh generation techniques emphasized constructing irregular grids to better model curved boundaries. Thacker² has briefly reviewed the different techniques.

Although mesh generation is a very broad and diverse area of research, its general approach can be divided into the following groups:

2.1 NODE GENERATION APPROACH

In this procedure, the interior nodes are generated first. They are then connected to form triangular or quadrilateral elements. This is a very basic method which has been used by many researchers, such as: Nguyen³, Cavendish, Field and Frey⁴, Cendes, Shenton and Shahnasser⁵, Sadek⁶ and Lo⁷. Such an approach is useful for automatic generation because it is very easy to generate triangular, tetrahedral and quadrilateral elements; however, it is unable to adequately generate the hexahedron elements.

Figure 2.1 shows the basic idea of the node generation approach.

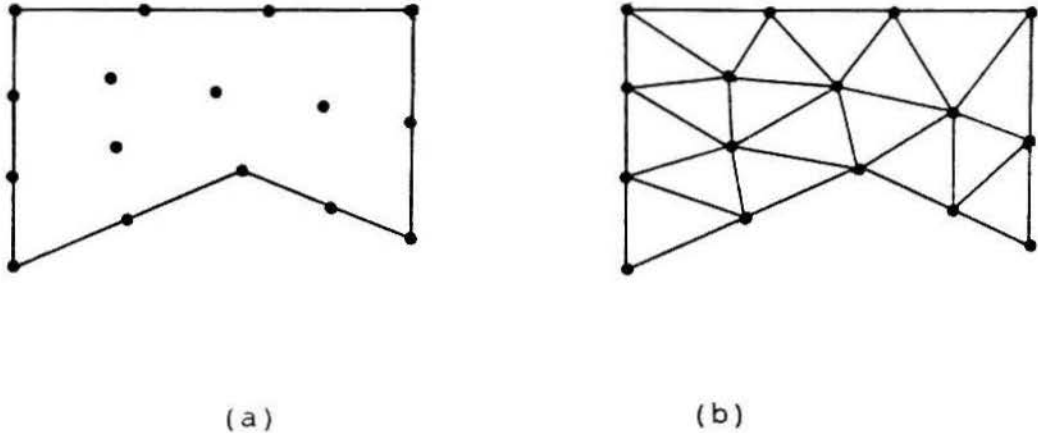


Figure 2.1 Node generation method: (a) all nodes are generated, then (b) connected to form elements.

2.2 MAPPING APPROACH

The mapping approach begins by developing the mesh in a unit square in a reference coordinate parametric space. It is then mapped to the final coordinate frame by blending functions.

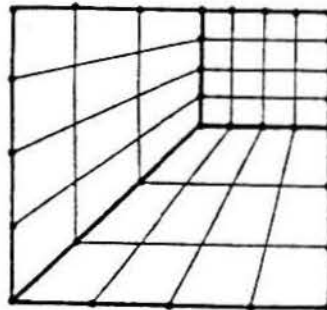


Figure 2.2 Mapping approach

For effective use of this procedure, an arbitrary object must first be subdivided into four-sided regions, and then each region is meshed separately. Zienkiewicz and Phillips⁸, Cohen⁹, Perucchio, Ingraffea and Abel¹⁰, Cook¹¹ and Gordon and Hall¹² have developed this approach. Figure 2.2 shows a model generated by this mapping technique.

2.3 GRID-BASED APPROACH

This method begins with an infinite grid superimposed on the object. The grid cells that fall outside the object are removed, then the grid cells that intersect the object boundary are adjusted or trimmed to fit the object. Thacker, Gonzalez and Putland³² and Heighway and Biddlecombe³³ use this approach to create finite element models. Yerry and Shephard²²⁻²⁴ use a modified quadtree and octree method to generate grids and to develop the final mesh. Figure 2.3 shows a model generated by grid-based approach.

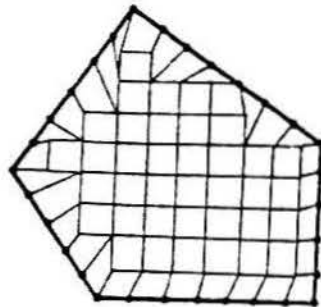


Figure 2.3 Grid-based approach

2.4 DECOMPOSITION APPROACH

The decomposition approach is a technique that continuously subdivides the model, until certain element sizes or limits are reached. Bykat¹⁴, Yeung and Hsu¹⁵, Sluiter and Hansen¹⁶, Sluiter¹⁷, Talbert¹⁸, Woo and Thomasma²⁰, Moscardini, Lewis, and Cross²¹ have used this approach to subdivide a surface to create quadrilateral elements or triangular elements, and subdivided a solid to create tetrahedral elements. An example is given in figure 2.4.

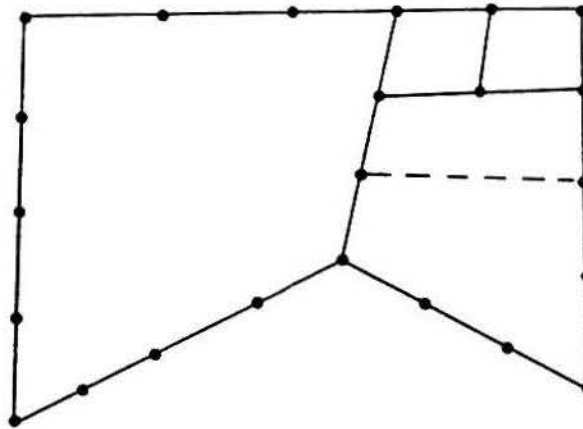


Figure 2.4 Decomposition approach

2.5 DRAG METHOD

For this procedure the basic elements are created in one dimension then "drag" into two dimensions. Similarly, simple elements in two dimensions can be "drag" to generate a three dimensional mesh as discussed by Park and Washam²⁵

and Pissanetzky²⁶. However this method is useful for symmetric models only, because the size of elements are different inside the solid. Figure 2.5 shows a model created by the drag method.

K Ho-Le³⁴ gives a more detailed classification of most of the method.

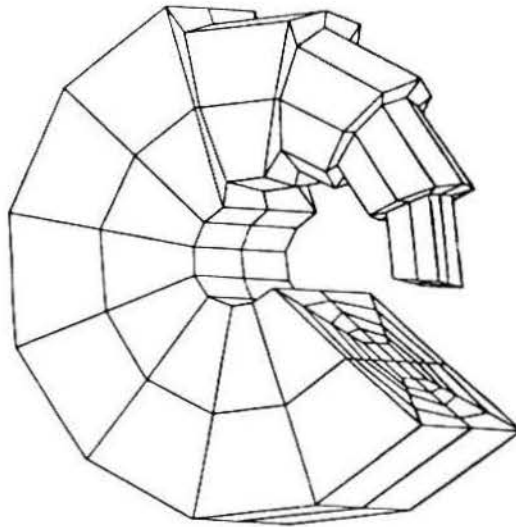


Figure 2.5 Model created by the drag method

2.6 AUTOMATIC MESHING OF THREE DIMENSIONAL SOLIDS

The methods described above generally have the capability to generate either quadratic, or triangular two dimensional elements. Drag method and coordinate blending functions have the capability to generate three dimensional hexahedral elements, however, they are limited to uniform and regular shaped geometries. The other methods can generate only tetrahedral elements.

Although the tetrahedral meshes are not as accurate

numerically in analysis¹³, they are ideal for automatic decomposition of complex polyhedra²⁰. This is based on the fact that every polyhedron can be divided into an arbitrary number of tetrahedra. Other elements can be formed by combining^{19,21} or subdividing¹⁶ tetrahedra, however, the final elements may have very irregular shapes.

A possible subdivision of a tetrahedron into four hexahedral is shown in figure 2.6. The aspect ratio of the elements is still suitable for finite element analysis.

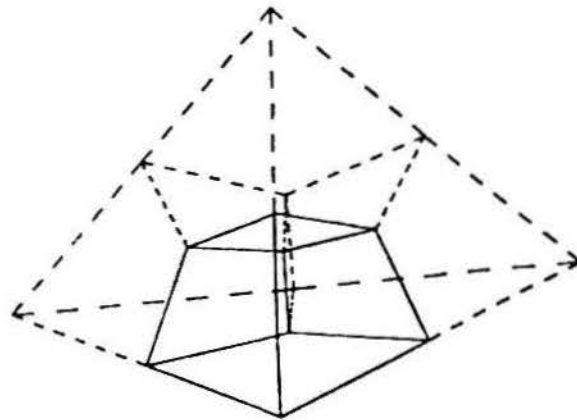


Figure 2.6 Four hexahedra from a tetrahedron

The subdivision method¹⁶, divides the six surfaces to get the surface elements first, then continues the subdivision for the volume to get the final solid elements. This method is very good for generating tetrahedral elements, however problem arise when creating three dimensional hexahedral elements in the volume, i.e. an impossible subdivision case may develop in the final volume.

Figure 2.7 shows this impossible case.

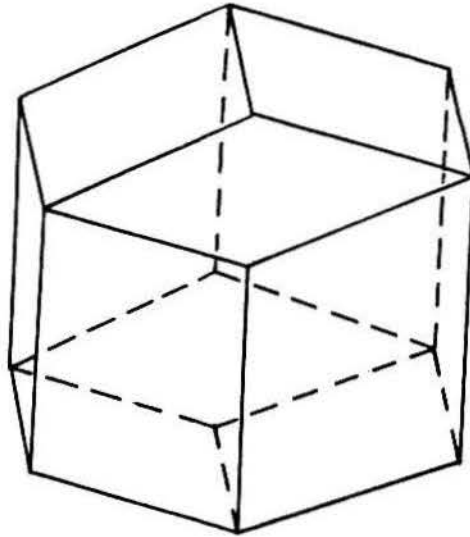


Figure 2.7 Impossible hexahedron elements

Such difficulties in creating hexahedral elements with an automatic procedure is the reason tetrahedral elements are most often used.

Pissanetzky²⁶ and Park and Washam²⁵ have successfully implemented the drag method to create hexahedral elements. Now quadrilateral elements are first created on a boundary surface then the elements are swept through the region to form solid hexahedral elements. This procedure is a good for regular shapes, however it is poor for irregular shapes. Furthermore, the element size may be different from point to point. Automatic meshing of irregular solid models with hexahedron elements remains a research challenge.

2.7 FREE-FORM DEFORMATION

The contents of this section assumes that the reader is familiar with the Bézier curves and surfaces patches³¹.

Free-form deformation is defined in terms of a tensor product trivariate rational Bernstein polynomial²⁸⁻³¹ which takes the following form.

$$P(s,t,u) = \frac{\sum_{i=0}^1 \sum_{j=0}^m \sum_{k=0}^n w_{ijk} B_i^1(s) B_j^m(t) B_k^n(u) P_{ijk}}{\sum_{i=0}^1 \sum_{j=0}^m \sum_{k=0}^n w_{ijk} B_i^1(s) B_j^m(t) B_k^n(u)}$$

where

$$B_i^1(s) = \frac{1!}{i!(1-i)!} (1-s)^{1-i} s^i, \quad i = 0, 1, \dots, 1.$$

$$B_j^m(t) = \frac{m!}{j!(m-j)!} (1-t)^{m-j} t^j, \quad j = 0, 1, \dots, m.$$

$$B_k^n(u) = \frac{n!}{k!(n-k)!} (1-u)^{n-k} u^k, \quad k = 0, 1, \dots, n.$$

$$s = \frac{X - X_{\min}}{X_{\min} - X_{\max}}, \quad t = \frac{Y - Y_{\min}}{X_{\min} - X_{\max}}, \quad u = \frac{Z - Z_{\min}}{X_{\min} - X_{\max}}$$

$B_i^1(s)$, $B_j^m(t)$, and $B_k^n(u)$ are called Bézier blending functions, and s, t, u are the local coordinates of a point with respect to the deformation region. P_{ijk} is the x, y, z coordinates of the displacement control points i, j, k in a solid that form an $(m+1) \times (n+1) \times (l+1)$ array of points. w_{ijk} is the weight of each control point. M, n, l are the total numbers of the control points to be used in each i, j, k direction. Any node that falls inside the control points boundary box $0 \leq s \leq 1$, $0 \leq t \leq 1$, $0 \leq u \leq 1$ will be calculated according to the Bernstein polynomial and moved to a new

location. Many more important properties of Bézier curves can be found in the article written by Bohm, Farin and Kahmann³¹.

The following example explains the basic procedure in creating the free-form deformation model. First consider a simple cuboid flexible plastic mesh model in the space which is to be deformed. Figure 2.8 shows this simple flexible plastic model.

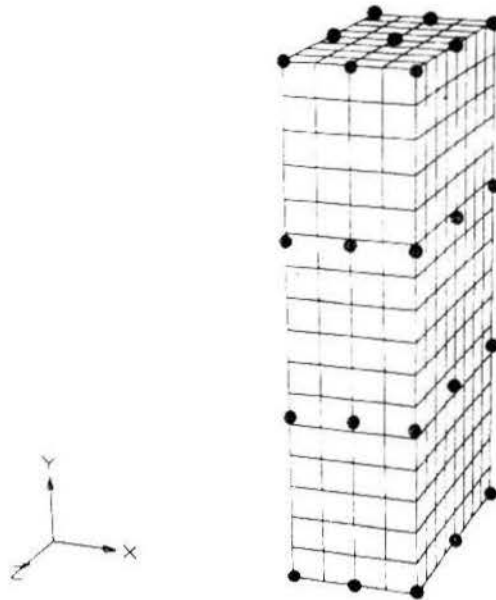


Figure 2.8 Original plastic model with control points

Next create free-form deformation control points (the black dots as shown in the figure 2.8) for the model. The model is now ready for deformation. The next step is to move the control points to new locations as shown in figure 2.9. The flexible plastic model inside the control points has been deformed to the shape as shown in figure 2.10.

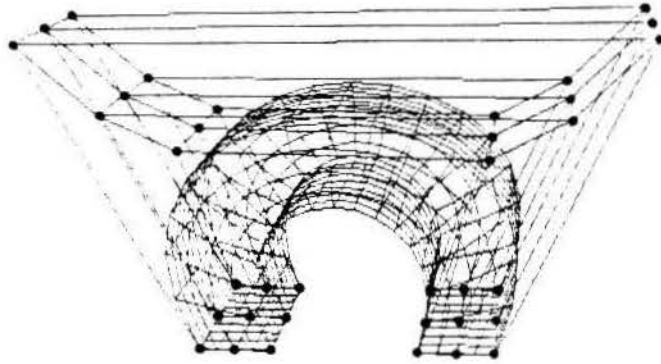


Figure 2.9 Control points in deformed position

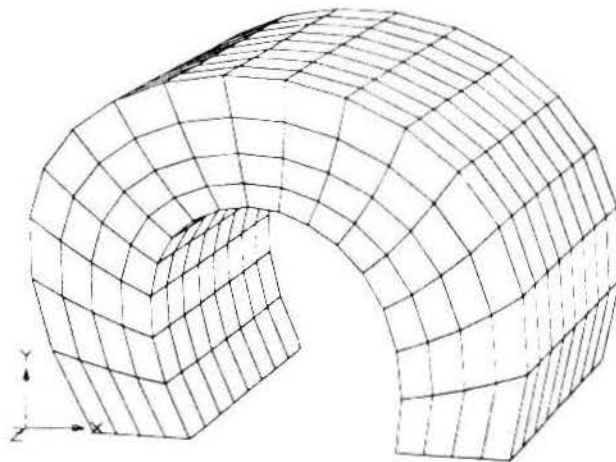


Figure 2.10 Deformed plastic model

2.8 MESH SMOOTHING

Once a solid model has been completely divided into hexahedral elements, the mesh is ready for a finite element

analysis. However, to insure optimum finite element results there is an element shape requirement²⁷. A good element should have an aspect ratio, i.e. max-side/min-side, that does not exceed roughly 7 for good displacement results and roughly 3 for good stress results. Generally an element that has an internal angle of less than 20 degrees or greater than 130 degrees is considered inadequate.

Smoothing is used to improve the shapes of the elements in a mesh by moving the interior nodes according to some rule. Effective smoothing provides the following capabilities:

- (1) Nearly equal areas or size of adjacent elements,
- (2) Gradually changing sizes of elements,
- (3) Gradually changing shapes of elements,
- (4) Nearly right angle corners in each element.

The Laplacian method is the most commonly used mesh smoothing technique. It changes the coordinates of each mesh node, in turn, to the average of the coordinates of all the immediate neighbors of that node. As shown in figure 2.11, the new position V .

$$V = 1/n \sum_{i=1}^n V_i$$

where $V_i = (x_i, y_i)$, $i = 1$ to n are the directly connected neighbors of this node.

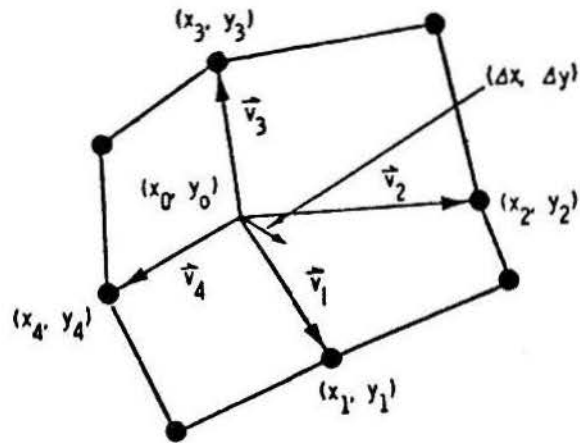


Figure 2.11 Laplacian smoothing

In order to get as close to the same element size as possible, the Laplacian method needs to be modified by using the area of the surrounding element to pull the node.

$$V = \frac{\sum A_i C_i}{\sum A_i}$$

where A_i is the area of the surrounding i -th elements and C_i is the vector from the corner node to the centroid of that element. The centroid of a element is defined by the average of the coordinates of four corner nodes:

$$C = \left(\frac{1}{4} \sum_{i=1}^4 x_i, \frac{1}{4} \sum_{i=1}^4 y_i \right)$$

where (x_i, y_i) , $i = 1, 2, 3, 4$, are the positions of the corners

of the element. If we set all these A_i 's to 1.0, the resulting smoothing technique is similar to Laplacian smoothing. Thus, the method is called "area-weighted Laplacian" or "centroid area pull." Figure 2.12 shows the centroid area pull smoothing method. This centroid area pull method gives better equal size results than the Laplacian method²⁷.

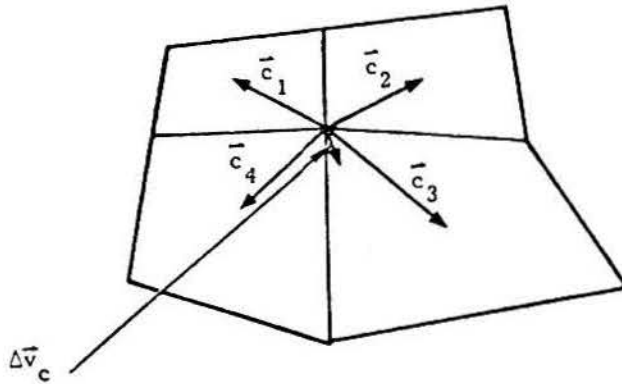


Figure 2.12 Centroid area pull smoothing

3.0 AUTOMATIC MESHING WITH HEXAHEDRON ELEMENTS

Generating an acceptable three dimensional finite element mesh of irregular shapes using hexahedral elements, with only a minimum of human interaction, is the goal of this work. This section describes the method developed by this work to generate such a mesh. The procedure couples both free-form deformation solid modeling and mesh smoothing as described in section 2.0.

The complete automatic meshing procedure consists of several steps. The first step is to create an initial cuboid primitive by eight points in the space. This primitive should consist of six surfaces and twelve lines as shown in figure 3.1. The next step applies free-form deformation to the initial cuboid such that the desired final geometrical shape is created. This step often severely distorts the mesh such that it is not acceptable for finite element calculations as shown in figure 3.2.

To establish an acceptable mesh on the deformed shape, a solid subdivision and mesh transition step is invoked. This step creates hexahedral elements from the solid model. An important aspect of this step is to insure that element continuity is completed. The final step invokes a smoothing process to bring all elements into an acceptable shape.

3.1 INITIAL CUBOID PRIMITIVE

At the beginning of the procedure, a cuboid primitive and a specific element size is defined. This cuboid

primitive is defined in the space by eight corner points. The user can choose any size and orientation of the cuboid primitive he desires, however a cuboid primitive relatively close to the final shape is easier for adjusting the control points of free-form deformation. After the eight points and element size have been chosen the model is ready for geometric deformation. A simple surface mesh is created in figure 3.1 for a view of the primitive model.

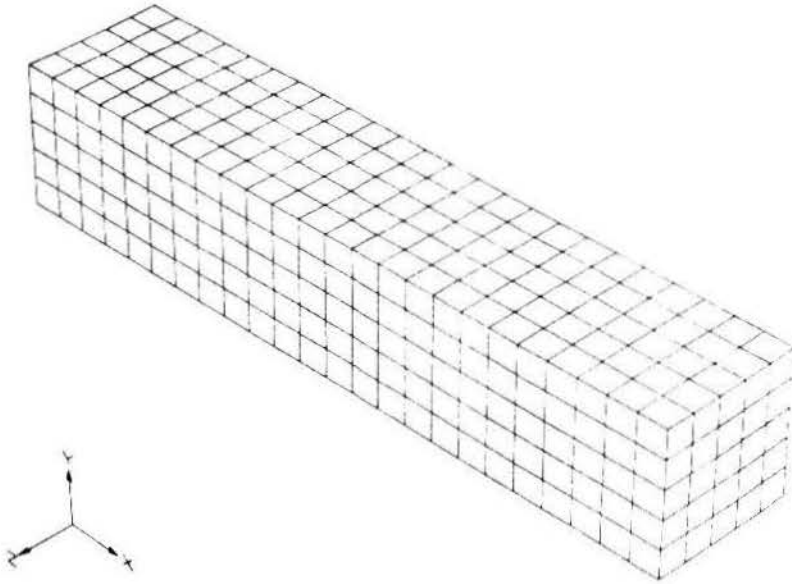


Figure 3.1 Cuboid primitive

3.2 FREE-FORM DEFORMATION

By moving the control points to new locations, as described in section 2 the model is deformed to the shape desired. However, the elements in this deformed model may have been distorted to very high aspect ratios. Figure 3.2

gives such a situation in the mesh which is deformed from the cuboid primitive created in figure 3.1.

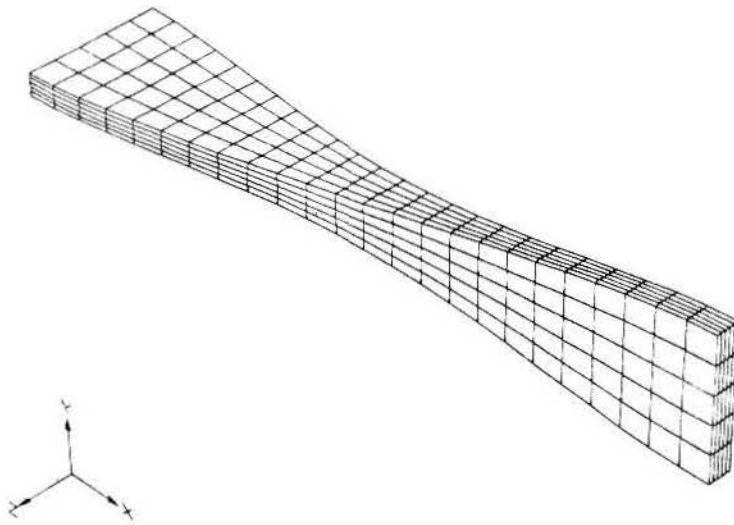


Figure 3.2 A deformed model with high aspect ratio elements

3.3 SOLID SUBDIVISION

The appropriate sized elements are generated by first locating the nodes on each of the boundary edges of the solid to set the initial nodal locations. These locations are chosen such that the element size will be near the size initially set by the user. This is followed by a subdivision process on the solid. This process is accomplished as follows:

1. Initially dividing the existing solid into two solids.
2. Selecting the number of nodes on the lines that define the joining boundary of these two solids. The number of nodes along this boundary is chosen such that the element

size will be near the size initially set by the user.

3. Force the condition that an even number of nodes exist on the lines that define the joining boundary of the two solids.

4. At this point, all edges of each solid have predefined nodes. Each solid is then checked to see if a) further subdivision is required, b) a transition situation exists, or c) the meshing is complete.

5. If subdivision is required, return to step 1.

6. If transition is required, proceed to the algorithm defined in Section 3.4.

7. When all subdivisions and transitions are completed, all elements have eight nodes and the mesh is ready for smoothing as described in Section 3.5.

A complete description of the above algorithm is outlined in Reference 36.

3.4 TRANSITION METHOD

If, during the processing by the above defined subdivision method, the number of nodes on one of the edges at one parameter axis equals two, and one or more of the remaining edges is greater than two, a transition region must be created. This means it is necessary to add nodes on the surfaces and inside the volume to create more new elements. In making these transitions, the three cases that may develop are explained below.

1) Transition case 1 - one of the four lines in a

parallel direction contains more than two nodes and the remaining three lines have only two nodes. This region will be processed by adding two nodes on each surface adjacent to the line containing the four nodes. These new nodes are labeled with an A in figure 3.3. Two more nodes, labeled with a B in figure 3.3, are created inside the cuboid to create five hexahedral elements.

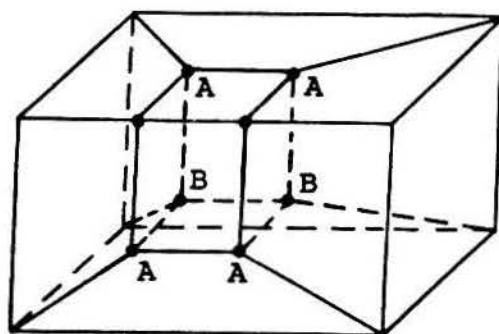


Figure 3.3 Transition case 1

2) Transition case 2 - two adjacent lines in a parallel direction contain more than two nodes and another pair of the lines have only two nodes. Adding two nodes (shown labeled with an A in figure 3.4) on both of the surfaces, creates four hexahedral elements to accommodate this transition.

3) Transition case 3 - three lines in a parallel direction contain more than two nodes and a single line contains two nodes. There is no way to directly accomplish the transition for this case, therefore an indirect method

must be used. This indirect method adds a cube inside the cuboid to create seven elements, as shown in figure 3.5. These new elements can be processed by the previous two methods to produce the eight node hexahedral elements.

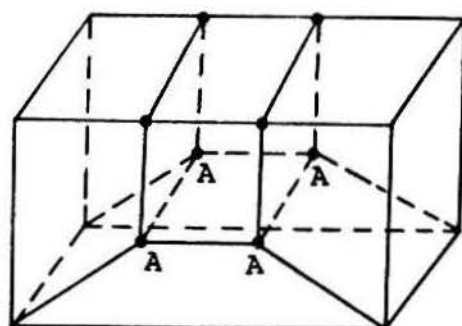


Figure 3.4 Transition case 2

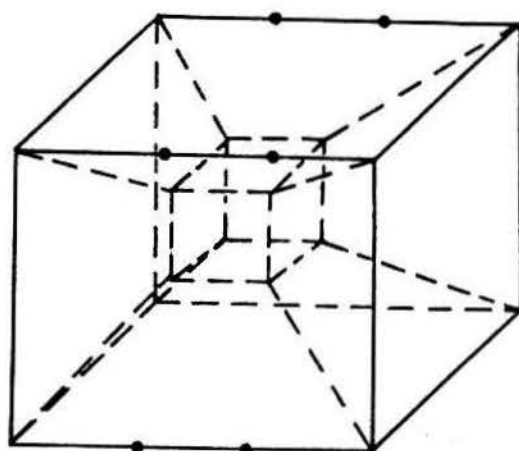


Figure 3.5 Transition case 3

3.5 SMOOTHING PROCEDURE

As described in the previous section, a smoothing procedure is required to provide proper elements for the

finite element analysis. After extensive testing on the samples, the centroid area pull method gave satisfactory results in element size. After running through all the steps of the process, the final result is shown in figure 3.6.

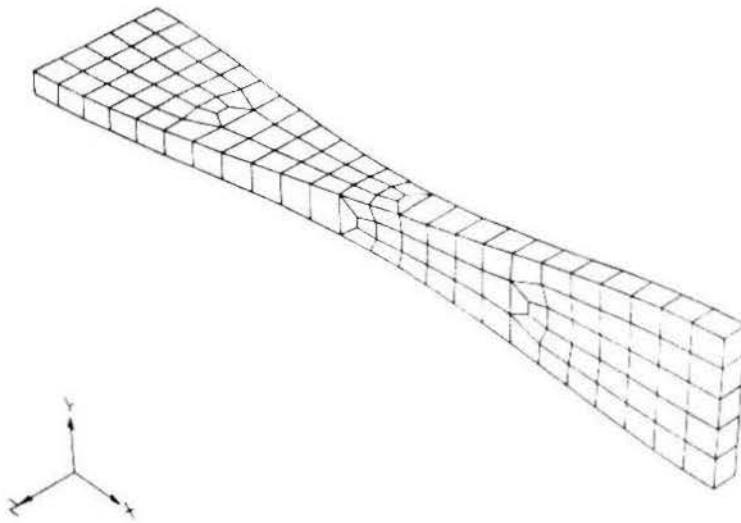


Figure 3.6 Final result of finite element mesh

4.0 EXAMPLES OF AUTOMATIC MESHING

This section presents some examples and provides some explanation for the basic technique for creating the free-form finite element models.

4.1 Transition head

The first example gives a step by step explanation of the capability of this procedure by modeling a part that transforms from a rectangular surface to a circular surface. The initial step defines the cuboid by specifying the information for the eight vertex points and the required size for the elements. This original cuboid with user's defined element size is shown in figure 4.1.

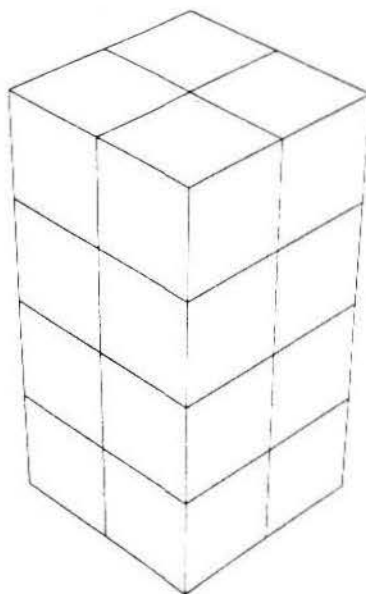


Figure 4.1 Original cuboid mesh

Next the control points for the free-form deformation are input. The location of control points are then adjusted so that all the nodes inside the bounding box form the shape

of the transition region. Figure 4.2 shows the control point net as well as the first deformation of this example.

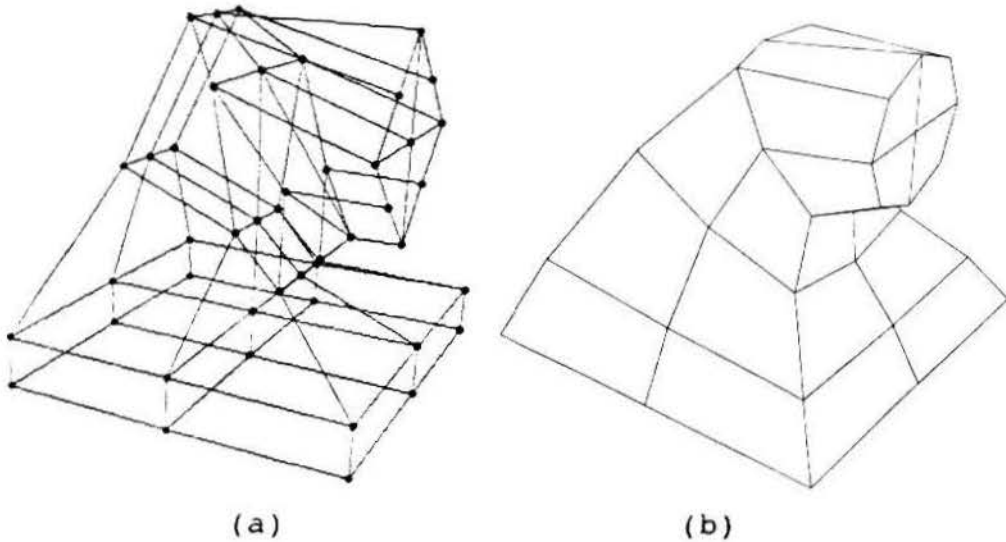


Figure 4.2 (a) Control point net and (b) deformation of the transition head

The next step is to use the mesh generation procedure to create the hexahedral elements using the solid subdivision and transition methods. Figure 4.3 shows the results of the finite element mesh generated by this procedure.

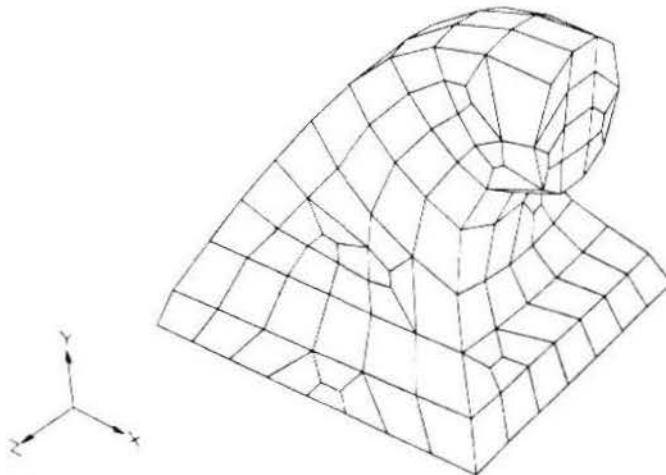


Figure 4.3 Uniform mesh of transition head

After the finite element mesh is generated, the smoothing procedure is applied to readjust the shape of the elements. Figure 4.4 shows the results after the centroid area pull smoothing method has been used.

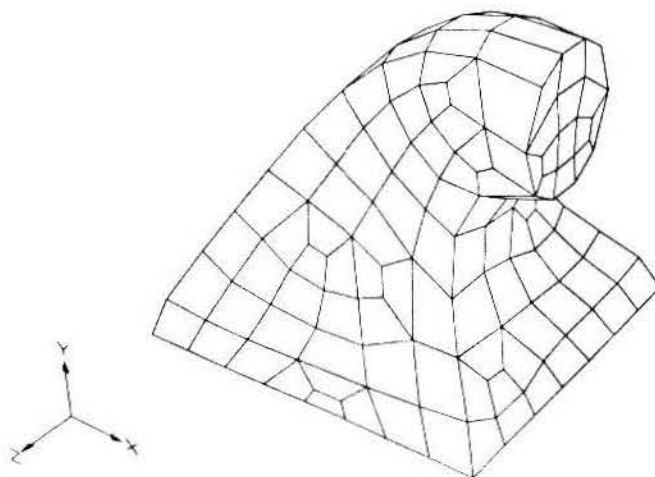


Figure 4.4 Smoothed mesh of transition head

4.2 Hour glass model

This example demonstrates the capability to create uniformly sized elements in a model with large dimensional variations. As before, the initial model begins from the regular cuboid similar to that shown in figure 4.1. Figure 4.5 shows the hour glass model after free-form deformation. The sizes of the elements at the middle of the model are now significantly different from those elements on both ends. These middle elements have very high aspect ratios which are obviously not well formed for finite element analysis.

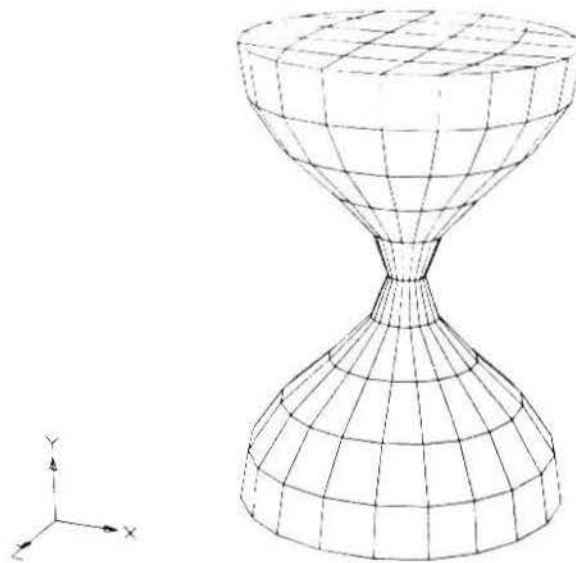


Figure 4.5 Uneven element size of hour glass model

Figure 4.6 shows the results of the model after the subdivision and transition is applied. The elements at the middle of the model are becoming uniform in size and have better aspect ratios. Finally the smoothed result is shown in figure 4.7.

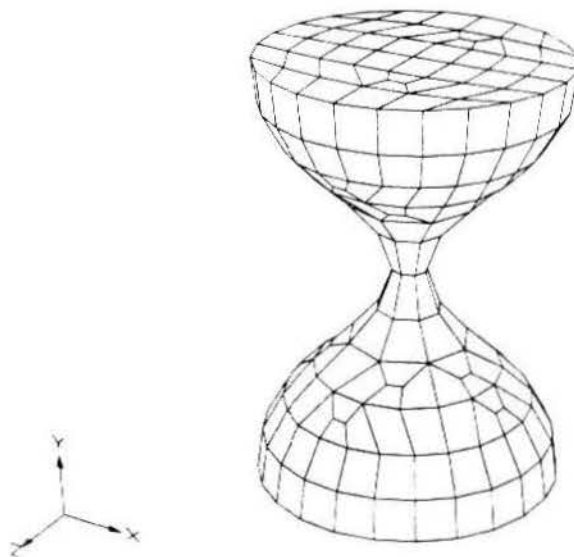


Figure 4.6 Uniform mesh of hour glass model

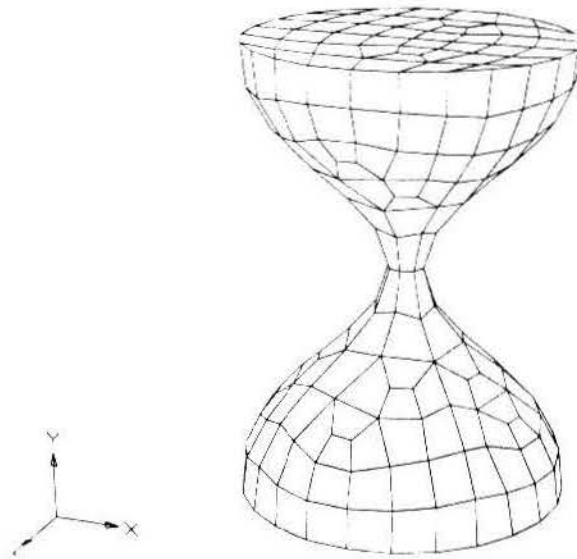


Figure 4.7 Smoothed elements of hour glass model

4.3 Hammer head

A significant capability of the free-form deformation procedure is the ability to generate arbitrary shapes. The hammer head model created here uses $2 \times 2 \times 9$ degrees of free-form deformation to define the shape. Figure 4.8 shows the original mesh of the model.

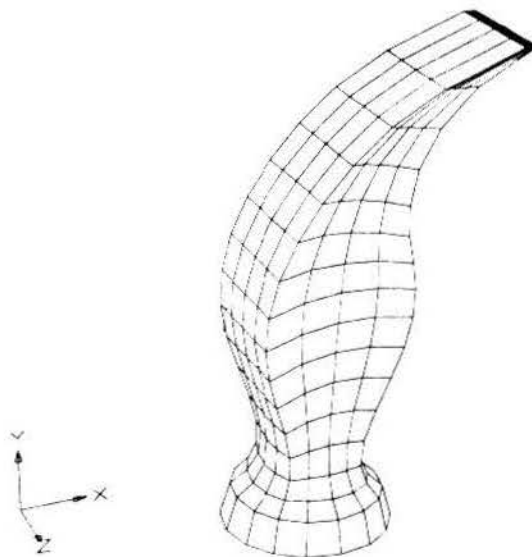


Figure 4.8 Original mesh of hammer model

Note that the model is not symmetric, therefore it is tedious for a traditional geometry modeler to create a model like this one. Figure 4.9 shows a mesh of the object.

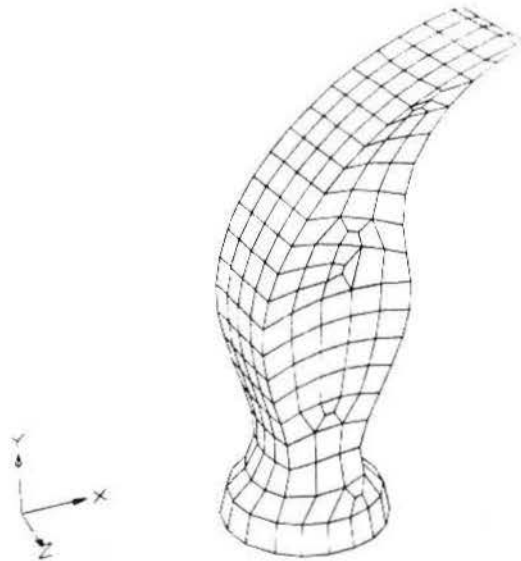


Figure 4.9 Mesh result of hammer head

5.0 CONCLUSIONS

The free-form deformation method is a very powerful tool to use in creating solid models of complex objects. Hexahedral (i.e. brick) elements are generally the element of choice for performing three dimensional finite element analysis. The work presented here couples the ability to create complex solid models using the free-form deformation method and the generation of acceptable finite elements for analysis purposes. The method is automated because a user is required to provide very little information to the algorithm that defines the procedure.

The conclusions that can be drawn from this work are:

1. The free-form deformation solid generation procedure can create solid geometric models appropriate for hexahedral finite element modeling.
2. A subdivision method that automatically generates hexahedral elements is appropriate to initially create elements within the deformed solid.
3. A transition scheme is necessary to allow the elements to effectively model severe changes in geometry.
4. Centroid area pull smoothing is adequate to properly reform the elements into acceptable shapes for finite element analysis.
5. Proper use of the method developed here will allow very complex finite element models of solid objects to be generated with a minimum of input from a user.

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APPENDICES

APPENDIX A
RESULTS ANALYSIS AND FUTURE DEVELOPMENT

RESULTS ANALYSIS

Because this mesh generator is supported by free-form deformation (whenever a subdivision line is created it needs to be deformed first to obtain the final element size), the processing costs are high.

Figure A.1 shows the smoothed result of the hammer head model. It can be seen that the model near the bottom of the hammer is squeezed in. This is because the smoothing processor relocated the nodes to an unsymmetric location on the surface. It can be solved by increasing the number of elements in the model to gradually change the surface in that area, or by creating two models separately at that critical area and by using the G2 continuity technique²⁸ to connect them into a continuous model.

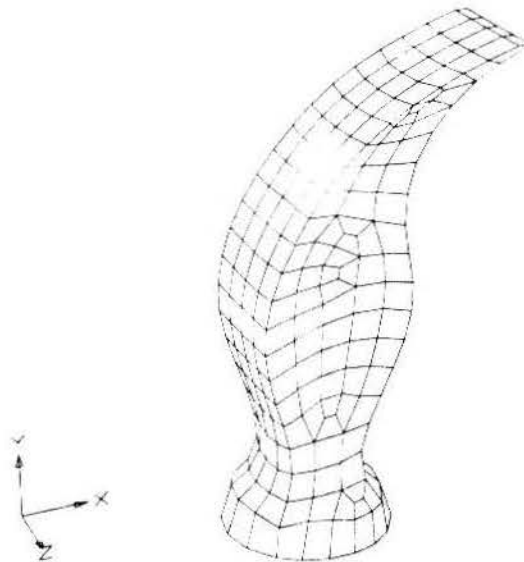


Figure A.1 Hammer model after smoothing

FUTURE DEVELOPMENT

The free-form deformation technique used in this research is based on the Bézier method which is a global control type of deformation. Therefore, it is easy to create the global smooth models, but models with irregular shapes requiring many local deformations (and therefore many control points) are difficult to do. future refinements can be made to enhance its capabilities.

1. The development of free-form deformation with the B-spline base³⁵ will better handle irregular shape.

2. The improvement of the subdivision processor to access many local deformations would also increase performance on handle irregular shapes.

3. Developing the variable element densities in certain locations may meet the requirements for the change of different local stresses or stains and produce better results in finite element analysis. This development can also be used to solve some uneven shapes.

CONNECTION WITH MOVIESTAR

MOVIESTAR, which was developed by Civil Engineering Department, is a very powerful pre/post processor which is connected with some very famous finite element software such as: ABAQUS, NASTRAN and MARC. A connection with MOVIESTAR would be a final destination for this automatic free-form deformation mesher.

At this time only a CAEDOS neutral file is created for

MOVIESTAR to do the pre-processing with the finite element software. Connecting directly with the MOVIESTAR database can make the best usage of the model created by this free-form deformation mesher. The meshed models can be run through the OPTIMIZER in MOVIESTAR to concatenate with other models into a continuous consolidated model for best simulation of the real object.

An ABAQUS finite element analysis result is shown in figure A.2, it can be proven that this automatic mesh generator is working properly in modeling and meshing.

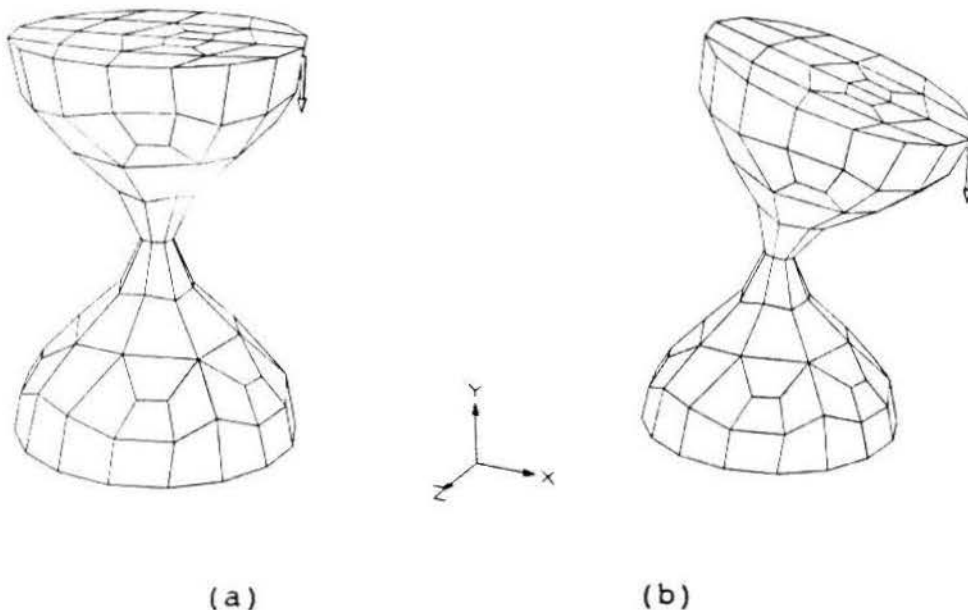


Figure A.2 A finite element analysis result (a) before testing, and (b) after testing.

APPENDIX B
COMPLETE DISCUSSION OF ALGORITHM

PREPARE ANALYSIS

Before running the program two files need to be prepared: a cuboid data file, and a free-form deformation control points file. All the input files need to be prepared with the same name but with different extensions. The cuboid data file is created with the extension *.sol; it contains the user defined length of the elements and the coordinates of the eight corner points. The free-form deformation file needs to be created with extension *.ffd; it contains the relative movements of all the control points and their weights.

The mesh program will begin by reading the cuboid data file and free-form deformation file. A geometry model will then be created by the defined control points. At this time a surface deformation file with the extension *.def is created for MOVIE.BYU to view the picture. The nodes on all the edges of the geometry will then be readjusted to have uniform length all around the edges. This nodal information will then be mapped back to the original cuboid as a reference for the subdivision process.

AUTOMATIC MESHING

After all the new boundary nodal information is defined and stored, the automatic meshing is ready to process for hexahedral elements. The steps of the entire algorithm are described as below:

1. A flag check is performed to check the number of

nodes on the edges of the solid to decide which processing step to choose next. For each of i, j, k direction, if the number of nodes on the edges of a direction are all greater than two, which means this solid can be subdivided, then the subdivision is performed. If the number of nodes on all the edges are all equal two then the subdivision process is done, and hexahedral element is created. If the number of the nodes on one of the edges is equal two and at least one of the rest of the edge has the number of nodes greater than two then there is a transition required to break the solid down into several hexahedral elements.

2. In cases where a subdivision is required, the subdivision priority is set to j -direction, then the k -direction, and last the i -direction. The subdivision starts with the surface containing the edge which has the least number of nodes, by defining a subdivision line between two opposite edges on the surface to be the best line for division. While looking for the best subdivision line, a search for the existing old line between those two edges needs to be performed first. If there is an old line existing then use it, otherwise find the shortest distance to be the best line and store it for next line check. The purpose of checking the existing old lines is to ensure continuity between two adjacent elements.

3. This subdivision line then passes through the free-form deformation to determine the real length in the

final state, and new nodes are then added to this line to create the length for the elements. A check for an even number of nodes on the surface is required to creating hexahedral elements. If an odd number appears on the surface then one more node need to be added to that subdivision line to make it even.

4. After one surface is divided then the node on that edge will be used for the adjacent surface. The subdivision process will continue until three surfaces have been processed. The last surface will be divided by using the nodes passed in through both opposite edges, and without checking the best subdivision line.

5. After four surfaces in a parametric direction have been divided, two solids are then created and pushed into a stack for further processing.

6. In case a transition is required, then it is performed according to the transition rules explained in section three and stored in the stack for further processing.

7. While performing a transition procedure a check for old lines is also necessary to ensure continuity with adjacent elements.

8. After a subdivision or transition is performed the movement will then return to step 1 for the next appropriate processing. All the processes will then repeat until the nodes on each edge are equal to two, then a linear

hexahedral element is created.

9. After all the elements are created, smoothing will then be performed to readjust the nodes so they will have better aspect ratios for the elements.

10. Three files which will be created at the end of the program under the same name but with different extensions (*.geo, *.mov, *.nav) for further analysis.

APPENDIX C
CONTINUITY OF ELEMENTS

To insure that adjacent elements have common nodes, the order of subdivision and the transition process is critical. Figure C.1 shows the problem of such a discontinuity between two adjacent elements. This possibility can be avoided by checking the subdivision order.

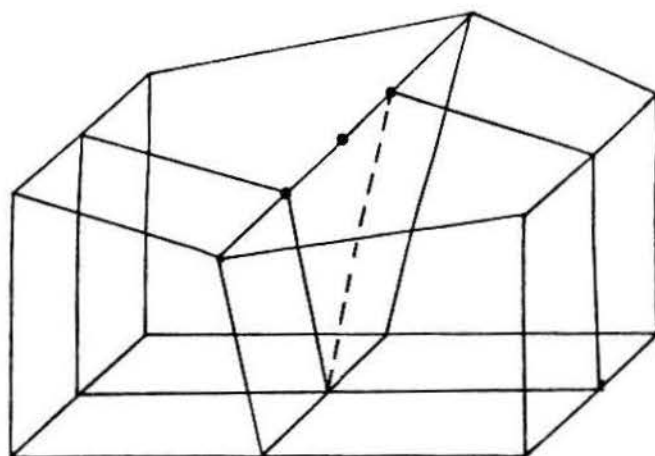


Figure C.1 Discontinuous surface elements

1. Checking old lines: While the new lines and nodes are created, the old information needs to be stored in a reference area. Before a new line is going to be created first check this new line with the old reference lines to see if this new line has already been created. Since a surface can only be used by just two solids, the line on that common surface can only be used twice. Therefore, if this line was an existing line, then use the old line data for this new subdivision line and delete that old line data from the reference area. If this new line was not an old

line then create it as a brand new line and store its information for the next reference line check.

2. Checking priorities: The order of subdivision priority must be set. If the direction of all the subdivisions are available at the same time then a processing order must be established. In this algorithm the j direction had been set to the first priority, k direction set to the second priority, and i direction set to the last priority. So after the meshing was done and ready to plot, the element at the bottom layer will be plotted first, and plotted from left to right as defined in the inverse of that mesh priority order.

For the case of transition 1 and 2 an old reference lines bucket needs to be set up for checking. For the case of priority in transition, an area need to be set up too, however it is a little different from the subdivision situation. The transition case 3 had been set to the first priority so that whenever a complicated case occurs it can process this case into the simple transition casees for further processing.

The transition case 2 had been set to the second priority and several conditions need to be met to qualify this case. Figure C.2 shows the example of the transition priority. The k direction transition seems to have the higher priority, however if this k direction transition had been processed first, as shows on the dashed line, then it

will cause discontinuity of the elements in the i direction on the top surface. Therefore, the i direction transition must be processed first.

Finally the transition case 1 has been set for last priority, because there is only one line to be processed.

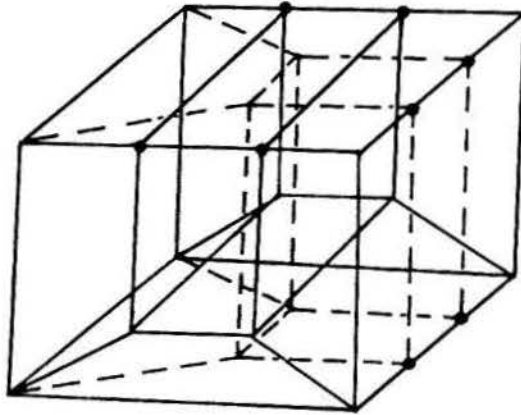


Figure C.2 Transition priorities

APPENDIX D**BÉZIER THEORY**

(As presented by Professor T. S. Sederberg)

BÉZIER CURVES

Bézier curves are named after their inventor, Dr. Pierre Bézier. Bézier was an engineer with the Renault car company and set out in the early 1960's to develop a curve formulation which would lend itself to shape design²⁸.

Bézier curves can be understood by thinking in terms of the center of mass of a set of point masses. For example, consider the four masses m_0, m_1, m_2, m_3 located at points P_0, P_1, P_2, P_3 (Figure D.1).

The center of mass of these four point masses is given by the equation

$$P = \frac{m_0P_0 + m_1P_1 + m_2P_2 + m_3P_3}{m_0 + m_1 + m_2 + m_3}$$

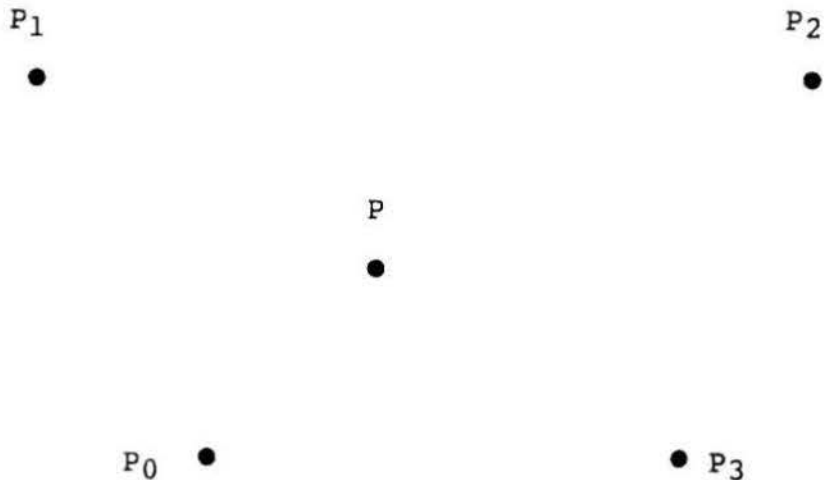


Figure D.1 Center of mass P , in four points masses.

Now, instead of having a fixed constant value of the masses at each point, each mass varies as a function of some

parameter t . Specifically let:

$$m_0 = (1-t)^3, \quad m_1 = 3t(1-t)^2, \quad m_2 = 3t^2(1-t), \quad m_3 = t^3.$$

Then the value of these masses as a function of t is shown in figure D.2.

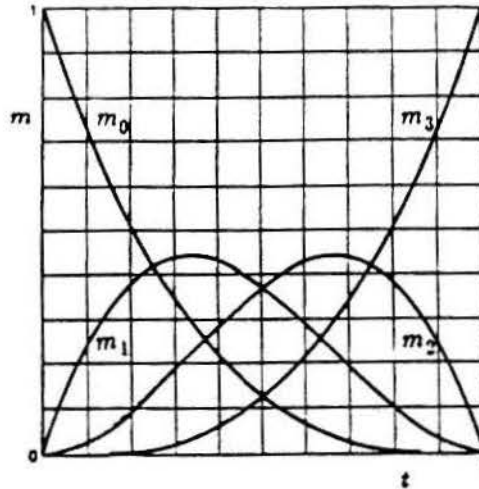


Figure D.2 Cubic Bézier blending functions

As t varies between 0 and 1, assume the masses have different weights, then their center of mass changes continuously. And a curve is swept out by the center of masses. This curve is called cubic Bézier curve, because the mass equations are cubic polynomials in t . For any value of t , $m_0 + m_1 + m_2 + m_3 \equiv 1$ the equation of this Bézier curve can be written as

$$P = m_0P_0 + m_1P_1 + m_2P_2 + m_3P_3.$$

Note that when $t=0$, $m_0=1$ and $m_1=m_2=m_3=0$ The curve is forced to pass through P_0 . Also, when $t=1$, $m_3=1$ and $m_0=m_1=m_2=0$, the curve passes through point P_3 . Furthermore, the curve is tangent to P_0-P_1 and P_3-P_2 . These properties make Bézier

curve an intuitively meaningful way to describe free-form shapes.

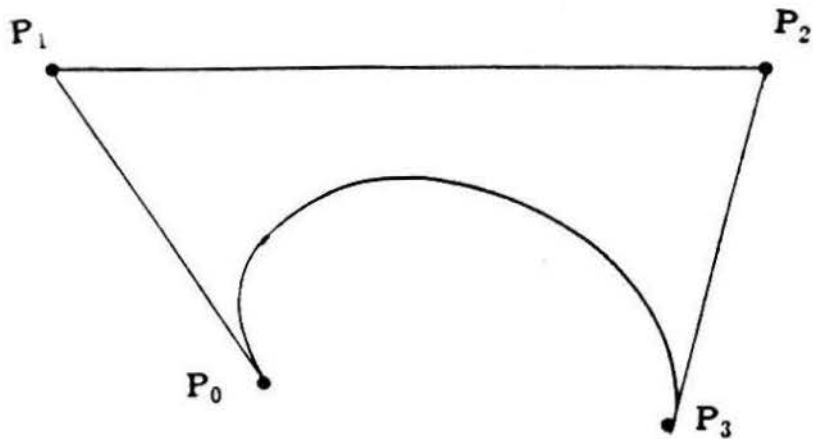


Figure D.3 Cubic Bézier curve

Some examples of cubic Bézier curves which illustrate these properties are given in figure D.4.

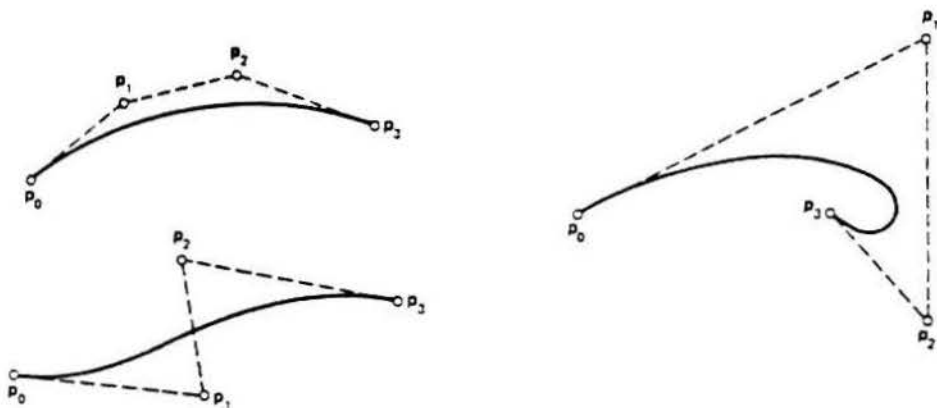


Figure D.4 Example of Bézier curves.

These variable masses m_i are normally called blending

functions and their locations P_i are known as control points or Bézier points. These blending functions, in the case of Bézier curves, are known as Bernstein polynomials.

Bézier curves of any degree can be defined. A Bézier curve of degree n will have $n+1$ control points whose blending functions are given as $B_i^n(t)$

$$B_i^n(t) = \frac{n!}{i!(n-i)!} (1-t)^{n-i} t^i, \quad i = 0, 1, \dots, n.$$

For example, let $n=3$ then $m_0=B_0^3=(1-t)^3$, $m_1=B_1^3=3t(1-t)^2$, $m_2=B_2^3=3t^2(1-t)$ and $m_3=B_3^3=t^3$. $B_i^n(t)$ is also referred to as the i th Bernstein polynomial of degree n . Many important properties of Bézier curves can be found in the article by Bohm, Farin and Kahmann³¹.

SUBDIVISION METHOD

The most important and fundamental algorithm for dealing with Bézier curves is the subdivision algorithm. This was devised in 1950 by F. de Casteljau and is sometimes referred to as the de Casteljau algorithm. It is also known as the geometric construction algorithm.

As a Bézier curve is defined for parameter values between 0 and 1. It is often useful to be able to subdivide this curve into two new Bézier curves, one of them in the range $0 \leq t \leq \alpha$, and the other in the range $\alpha \leq t \leq 1$. These two curves are considered to be together and equivalent to the single original curve from which they were derived. The control points for these two new curves are found using a

simple geometric rule. We illustrate with a cubic Bézier curve in figure B.5.

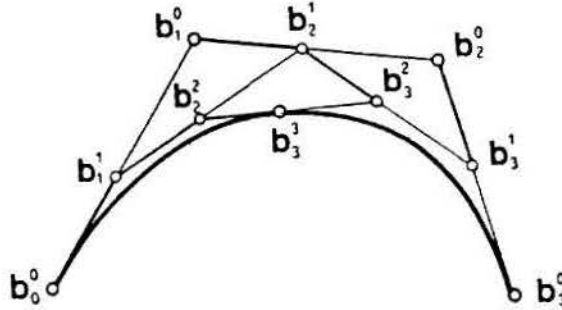


Figure D.5 Subdividing a cubic Bézier curve.

Let the original four control points be denoted b_0^0 , b_1^0 , b_2^0 , b_3^0 . The control polygon consists of three lines connected between four control points. α is the parameter value at which we wish to split the curve, then more auxiliary points have been created between the original control points. The auxiliary points are located at the distance α of the previous control points on the control lines. The recurrence relation of the control points can be expressed by the formula:

$$b_i^k(\alpha) = (1-\alpha)b_{i-1}^{k-1} + \alpha b_{i+1}^{k-1}$$

Now, two new Bézier curves are defined from these auxiliary points. For the left portion where $0 \leq t \leq \alpha$, there exist control points b_0^0 , b_1^1 , b_2^1 , and b_3^1 . The right portion where $\alpha \leq t \leq 1$ has control points b_3^1 , b_2^2 , b_1^2 , and b_0^0 . These two

Bézier curves can be repeated subdivided until small enough, then the Bézier curve points are defined.

The Bézier curves which have been defined above are known as integral Bézier curves. Since it is essential for the blending function to sum to one (otherwise the curve will change with the coordinate system) the blending function must be normalized by dividing through their total value.

$$X(t) = \frac{\sum_{i=0}^n w_i B_i^n(t) P_i}{\sum_{i=0}^n w_i B_i^n(t)}$$

The curve is called a rational Bézier curve, because the blending functions are rational polynomials. The important reason for rational Bézier curves is that they can define a curves without using higher order control points. Figure D.6 shows the change of weights in the control points to influence the shape of the curve.

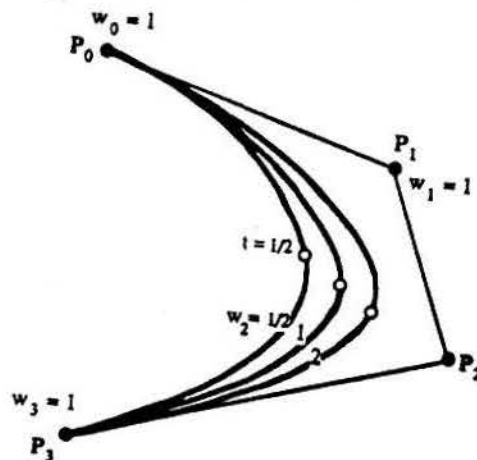
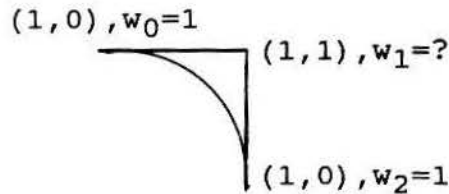


Figure D.6 Change the curve shape by the weights

The following example shows the way to create a circle by two method using the Bézier curve. To create a quarter circle from a line by degree two control points, the weight control method²⁸ needs to be used. In the following figure let a rational quadratic degree two Bézier curve was defined as:



$$x(t) = \frac{w_0 x_0 (1-t)^2 + 2w_1 x_1 t(1-t) + w_2 x_2 t^2}{w_0 (1-t)^2 + 2w_1 t(1-t) + w_2 t^2}$$

$$y(t) = \frac{w_0 y_0 (1-t)^2 + 2w_1 y_1 t(1-t) + w_2 y_2 t^2}{w_0 (1-t)^2 + 2w_1 t(1-t) + w_2 t^2}$$

In order to satisfied a circular equation, $x(t)^2 + y(t)^2 = 1$, we substitute x and y into the equation to solve for the weight at w_1 , and then the weight for the middle point $w_1 = \sqrt{2}/2$. Then this value will be used as a control point weight at that middle point. After the free-form deformation this line will become a quarter circle as defined.

Additionally, there is a fixed control points weight method can be used. By increasing the degree of the Bézier curve and adjusting the location of some control points, one can get the same quarter circle. Let a fourth degree Bézier curve have the same weight all around the control points,

$w_0=w_1=w_2=w_3=w_4=1$ and set $P_0=(0,1)$, $P_1=(0.4,1)$, $P_2=(0.8,0.8)$, $P_3=(1,0.4)$, $P_4=(1,0)$ the resulting curve will be exactly the same as the second degree curve has different weights on each control point.

BÉZIER SURFACE

After the Bézier curve was defined, this curve can be expand into a surface. Points on a Bézier surface are given by an extension of the general equation for points on a Bézier curve.

$$P(s,t) = \sum_{i=0}^m \sum_{j=0}^n P_{ij} B_i^m(s) B_j^n(t) \quad s,t \in [0,1]$$

Where P_{ij} are control points of the Bézier surface that form an $(m+1) \times (n+1)$ rectangular array of points. By using the binomial representation of the cubic tensor product, the Bézier surface patch takes the form:

$$P(s,t) = [s^3 \quad 3s^2(1-s) \quad 3s(1-s)^2 \quad (1-s)^3] P \begin{bmatrix} t^3 \\ 3t^2(1-t) \\ 3t(1-t)^2 \\ (1-t)^3 \end{bmatrix}$$

where

$$P = \begin{bmatrix} P_{33} & P_{23} & P_{13} & P_{03} \\ P_{32} & P_{22} & P_{12} & P_{02} \\ P_{31} & P_{21} & P_{11} & P_{01} \\ P_{30} & P_{20} & P_{10} & P_{00} \end{bmatrix}$$

The matrix P contains the position of the control points for Bézier surface. In the Bézier surface only the four corner points where $s=0$, $s=1$, $t=0$, $t=1$ are lie on the patch. Figure D.7 shows the bicubic Bézier surface.

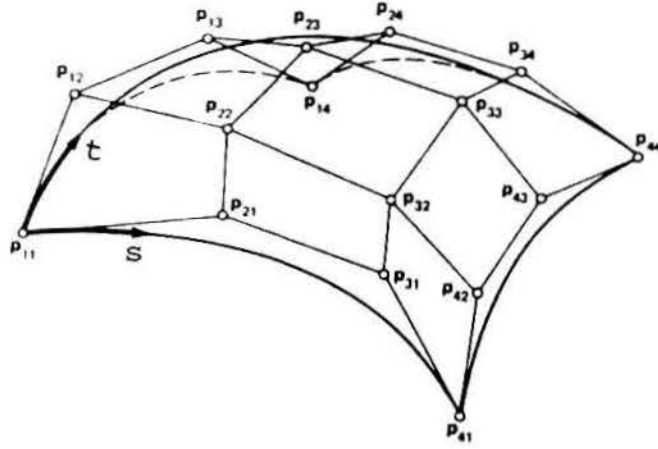


Figure D.7 Bicubic Bézier surface

APPENDIX E

SOURCE CODE OF FREE-FORM DEFORMATION MESH GENERATOR

```

C-----
C      subroutine gtsoli(fname)
C-----
C
C PROGRAM: GTSOLI.FOR
C
C PURPOSE:  get the required data for volume and prepared them for
C           volume meshing.
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 8, 1988
C MODIFIED BY:           CHECKED BY:           DATE:
C-----
C
C ALGORITHM EXPLANATION: READIN THE GEOMETRY FILE AND DEFORMATION FILE
C THEN DO THE LINE DEFORMATION AND STORE THE DATA FOR MESHING
C-----
C
C called:
C   readso.for -- read in the 8 corner point in a volume
C   orimsh.for -- store the premesh deformed data
C   frform.for -- readin the free formed deformation control points
C   linnod.for -- find the nodes in each line
C   defor1.for -- do the deformation for the lines
C   mshlin.for -- mesh the boundary lines first
C   mpback.for -- mapping the deformed nodes back to original hexahedron
C   drstor.for -- store the nodes for draw picture
C   prepar.for -- store the nodes information for volume meshing
C-----
C
C called by: threed.for -- the user's interface
C-----
C
C   include 'para.inc'
C   dimension xg(8),yg(8),zg(8),ck(8),nnum(12),dist(12)
C   dimension x(maxsid,12),y(maxsid,12),z(maxsid,12),w(maxsid,12)
C   dimension xo(maxsid,12),yo(maxsid,12),zo(maxsid,12)
C   character fname=(*)
C
C   store deformation information
C
C   call frform(fname,ierr)
C
C   read in the prime volume 8 points
C
C   call readso(fname,nx,ny,nz,unit,xg,yg,zg,ierr)
C   if(ierr .ne. 0)then
C     write(*,*) 'None exist file name'
C     return
C   end if
C
C   check the dimension
C
C   maxd=max0(nx,ny,nz)+1
C   if(maxd .gt. maxsid)then
C     write(*,*) ' The <maxsid> need to set to ',maxd
C     stop
C   end if
C
C   find the nodes for twelve lines
C
C   call linnod(nx,ny,nz,xg,yg,zg,x,y,z,maxsid,nnum)
C
C   store the node for plot the original mesh only, can be deleted
C
C   call nodsur(nx,ny,nz,x,y,z,maxsid)
C   call nodsto(fname,nx,ny,nz)
C
C   get the deform shape for the prime volume lines
C
C   call defor1(xo,yo,zo,x,y,z,nnum)
C
C   mesh the twelve boundary line nodes
C
C   call mshlin(nx,ny,nz,x,y,z,nnum,unit,dist)
C
C   map the deformed nodes back to the normal cube

```



```

c      call mback(nx,ny,nz,xg,yg,zg,dist,xo,yo,zo,x,y,z,w,nnum)
c
c store data for drafting purpose, no need at this time
c do not open it while maxdeg<30 or nodsto is in used
c      call drstor(x,y,z,nnum,unit)
c
c reset the storage address
c
c      call prepar(x,y,z,w,nnum,unit)
c      return
c      end

CNAME VMESH.FOR
C-----
c      subroutine vmesh(cname)
C-----
C
C PROGRAM: VMESH.FOR
C
C PURPOSE: CHOOSE THE REQUIRE MESH PROCESSER TO DO MESHING
C
C WRITTEN BY: SSUTA HSU      CHECKED BY:      DATE: OCT. 8, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C-----
C ALGORITHM EXPLANATION: check the number of nodes in each line, if the
c      number equal two then that will be the final state for that element
c      or the element need todo the transition on it, otherwise continue
c      the subdivision algorithm.
C-----
c      called:
c      push.for    -- stack one step in the loop
c      pop.for     -- pop up one step in the loop
c      vstore.for  -- store nodes for plot
c      xmesh.for   -- mesh in x-direction
c      ymesh.for   -- mesh in y-direction
c      zmesh.for   -- mesh in z-direction
c      lnflag.for  -- set flag to choose the required transition method
c      transi.for  -- set the required transition method
c      ffdefo.for  -- process the free formed deformation
c      wrtmov.for  -- write out the movie and navgraph files for futher
c                  processing.
C-----
c      called by: threed.for -- user's interface
c
c-----
c      character cname*(*)
c
c      write(*,*)' Automatic mesh generating now ...'
c
c      iter = 0
c      loop = 1
c
c      call push(loop,iter)
c
c...start the looping method
c
10      continue
c      call pop(loop,iter)
c
c check the number of nodes in all the lines, then do the proper method
c
c      call lnflag(loop,iflagx,iflagy,iflagz)
c      if(iflagx.eq.0 .and. iflagy.eq.0 .and. iflagz.eq.0)then
c          call vstore(loop)
c
c check the nodes in z,y,and x direction and mesh it
c
c      else if(iflagy.eq.6)then
c          call ymesh(loop,iter)
c      else if(iflagz.eq.6)then
c          call zmesh(loop,iter)
c      else if(iflagx.eq.6)then
c          call xmesh(loop,iter)
c      else

```

```

c
c do the transition for all the possible ways (for 0 <iflag < 8)
c
c     call transi(loop,iter,iflagx,iflagy,iflagz)
c     end if
c     if(iter.gt.0)go to 10
c
c     call surbnd
c     call smooth
c     call ffdefo
c     call wrtmov(cname)
c     call wrtnav(cname)
c
c     return
c     end

CNAME XMESH.FOR
c-----
c     subroutine xmesh(loop,iter)
c-----
c
c C PROGRAM: XMESH.FOR
c C PURPOSE: do subdivision in x direction
c C WRITTEN BY: SSUTA HSU      CHECKED BY:      DATE: OCT.8, 1988
c C MODIFIED BY:              CHECKED BY:      DATE:
c C-----
c C ALGORITHM EXPLANATION: CHECK THE BEST NODE FOR SUBDIVISION AND ADD THE
c C NEW NODES TO THE SUBDIVISION LINES
c C-----
c
c called:
c     vgtnod.for -- choose the required node to split with
c     vnewnd.for -- find the new node
c     vnewln.for -- find the new line
c     push.for   -- stack for the new loop
c
c-----
c
c called by: vmesh.for -- volume mesh controler
c-----
c
c     include 'para.inc'
c     dimension coor(4,maxnod),node(maxsid,12,maxloop)
c     common/auto/coor,node,last,unit
c     logical found
c
c store extra data for remesh check
c
c     do 20 j = 1,12
c         do 20 i = 1,node(1,j,loop)+1
c             node(i,j,maxloop-1) = node(i,j,loop)
20     continue
c     klast = last
c
c find the split line ending node for surface 1-3, then connect the
c third and fourth nodes to create the split line for surface 4,
c
c     mini = 10000
c     do 10 i = 1,3
c         if(node(1,i,loop).lt.mini)then
c             mini = node(1,i,loop)
c             line = i
c         end if
10     continue
c     if(line.eq.1)then
c         line1 = 1
c         num1 = 2
c         call vgtnod(loop,1,line1,line2,num1,num2)
c         call vgtnod(loop,1,line2,line4,num2,num4)
c         call vgtnod(loop,2,line1,line3,num1,num3)
c     else if(line.eq.2)then
c         line2 = 2
c         num2 = 2
c         call vgtnod(loop,1,line2,line4,num2,num4)
c         call vgtnod(loop,2,line2,line1,num2,num1)

```

```

        call vgtnod(loop,2,line1,line3,num1,num3)
    else
        line3 = 3
        num3 = 2
        call vgtnod(loop,1,line3,line1,num3,num1)
        call vgtnod(loop,1,line1,line2,num1,num2)
        call vgtnod(loop,1,line2,line4,num2,num4)
    end if
c
c find the nodes for each new split lines.
c
    call vnewnd(line1,line2,num1,num2,loop)
    call vnewnd(line1,line3,num1,num3,loop)
    call vnewnd(line2,line4,num2,num4,loop)
    call vnewnd(line3,line4,num3,num4,loop)
c
c split the old boundary line into two new lines
c
    call vnewln(line1,num1,loop)
    call vnewln(line2,num2,loop)
    call vnewln(line3,num3,loop)
    call vnewln(line4,num4,loop)
c
c check for the impossible transition case
c
    call push(loop,iter)
    loop = loop+1
    call push(loop,iter)
c
    return
    end

CNAME YMESH.FOR
c-----
    subroutine ymesh(loop,iter)
c-----
C
C PROGRAM: YMESH.FOR
C
C PURPOSE: do subdivision in y direction
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:                  CHECKED BY:          DATE:
c-----
C
C ALGORITHM EXPLANATION: FIND THE BEST SUBDIVISION LOCATION AND ADD THE
C NOCES TO THOSE SUBDIVISION LINES
c-----
c
c called:
c vgtnod.for -- choose the required node to split with
c vnewnd.for -- find the new node
c vnewln.for -- find the new line
c push.for   -- stack for the new loop
c-----
c
c called by: vmesh.for -- volume mesh controler
c-----
    include 'para.inc'
    dimension coor(4,maxnod),node(maxsid,12,maxloop)
    common/auto/coor,node,last,unit
    logical found
c
c store extra data for remesh check
c
    do 20 j = 1,12
        do 20 i = 1,node(1,j,loop)+1
            node(i,j,maxloop-1) = node(i,j,loop)
20    continue
        klast = last
c
c find the split line ending node for surface 5-7, then connect the
c third and fourth nodes to create the split line for surface 8.
c
    mini = 1000
    do 10 i = 5,7

```

```

        if(node(1,i,loop).lt.mini)then
            mini = node(1,i,loop)
            line = i
        end if
10    continue
    if(line.eq.6)then
        line1 = 5
        num1 = 2
        call vgtnod(loop,1,line1,line2,num1,num2)
        call vgtnod(loop,1,line2,line4,num2,num4)
        call vgtnod(loop,2,line1,line3,num1,num3)
    else if(line.eq.8)then
        line2 = 8
        num2 = 2
        call vgtnod(loop,1,line2,line4,num2,num4)
        call vgtnod(loop,2,line2,line1,num2,num1)
        call vgtnod(loop,2,line1,line3,num1,num3)
    else
        line3 = 7
        num3 = 2
        call vgtnod(loop,1,line3,line1,num3,num1)
        call vgtnod(loop,1,line1,line2,num1,num2)
        call vgtnod(loop,1,line2,line4,num2,num4)
    end if

c
c    find the nodes for each new split lines.
c
c        call vnewnd(line1,line2,num1,num2,loop)
c        call vnewnd(line1,line3,num1,num3,loop)
c        call vnewnd(line2,line4,num2,num4,loop)
c        call vnewnd(line3,line4,num3,num4,loop)
c
c    split the old boundary line into two new lines
c
c        call vnewln(line1,num1,loop)
c        call vnewln(line2,num2,loop)
c        call vnewln(line3,num3,loop)
c        call vnewln(line4,num4,loop)
c
c        call push(loop,iter)
c        loop = loop+1
c        call push(loop,iter)
c
c    return
c    end

CNAME ZMESH.FOR
c-----
c        subroutine zmesh(loop,iter)
c-----
c
c    PROGRAM: ZMESH.FOR
c
c    PURPOSE: do subdivision in z direction
c
c    WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
c    MODIFIED BY:                CHECKED BY:           DATE:
c-----
c
c    ALGORITHM EXPLANATION:  FIND THE BEST SUBDIVISION LOCATION AND ADD THE
c                          NOCES TO THOSE SUBDIVISION LINES
c-----
c
c    called:
c        vgtnod.for -- choose the required node to split with
c        vnewnd.for -- find the new node
c        vnewln.for -- find the new line
c        push.for  -- stack for the new loop
c-----
c
c    called by: vmesh.for -- volume mesh controler
c-----
c
c        include 'para.inc'
c        dimension coor(4,maxnod),node(maxsid,12,maxloop)
c        common/auto/coor,node,last,unit
c        logical found

```

```

c
c store extra data for remesh check
c
  do 20 j = 1,12
    do 20 i = 1,node(1,j,loop)+1
      node(i,j,maxloop-1) = node(i,j,loop)
20  continue
    klast = last
c
c find the split line ending node for surface 9-11, then connect the
c third and fourth nodes to create the split line for surface 12.
c
  mini = 10000
  do 10 i = 9,11
    if(node(1,i,loop).lt.mini)then
      mini = node(1,i,loop)
      line = i
    end if
10  continue
  if(line.eq.9)then
    line1 = 9
    num1 = 2
    call vgtnod(loop,1,line1,line2,num1,num2)
    call vgtnod(loop,1,line2,line4,num2,num4)
    call vgtnod(loop,2,line1,line3,num1,num3)
  else if(line.eq.10)then
    line2 = 10
    num2 = 2
    call vgtnod(loop,1,line2,line4,num2,num4)
    call vgtnod(loop,2,line2,line1,num2,num1)
    call vgtnod(loop,2,line1,line3,num1,num3)
  else
    line3 = 11
    num3 = 2
    call vgtnod(loop,1,line3,line1,num3,num1)
    call vgtnod(loop,1,line1,line2,num1,num2)
    call vgtnod(loop,1,line2,line4,num2,num4)
  end if
c
c find the nodes for each new split lines.
c
  call vnewnd(line1,line2,num1,num2,loop)
  call vnewnd(line1,line3,num1,num3,loop)
  call vnewnd(line2,line4,num2,num4,loop)
  call vnewnd(line3,line4,num3,num4,loop)
c
c split the old boundary line into two new lines
c
  call vnewln(line1,num1,loop)
  call vnewln(line2,num2,loop)
  call vnewln(line3,num3,loop)
  call vnewln(line4,num4,loop)
c
  call push(loop,iter)
  loop = loop+1
  call push(loop,iter)
c
  return
end
CNAME CHOLIN.FOR
c-----
  subroutine choln(nod1,nod2,oline,nold)
c-----
C
C PROGRAM: CHOLIN.FOR
C
C PURPOSE: CHECK TO FIND THE POSSIBLE OLD LINE
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:                  CHECKED BY:          DATE:
C-----
C
C ALGORITHM EXPLANATION:
c-----
c
c called by: vnext1.for -- find the next possible line
c-----

```

```

      include 'para.inc'
      dimension ick(maxsid),itemp(20)
      logical oline
      common/oldln/nnum(20,maxsid*6),lold(maxsid*6),nall
c
      oline=.false.
c
c...check for the discrete line case or partial line case
c
      ntemp = nod1
      icnt = 0
10      continue
      do 40 i = 1,nall
c...start normal checking
      if(ntemp.eq.nnum(1,i))then
        icnt = icnt + 1
        ick(icnt) = i
c...regular coincident case and final condition of small lines case
      if(ntemp.eq.nod2)then
        oline = .true.
        go to 60
      end if
c...partial line case
      if(lold(i).gt.2)then
        do 20 k = 2,lold(i)
          if(nod2.eq.nnum(k,i))then
            oline = .true.
            nall = nall + 1
            mc = 1
            nnum(1,nall) = nod1
            do 21 m = 1,icnt-1
              do 21 n = 2,lold(ick(m))
                mc = mc + 1
                nnum(mc,nall) = nnum(n,ick(m))
21              continue
              do 30 j = 2,k
                mc = mc + 1
                nnum(mc,nall) = nnum(j,i)
30              continue
              lold(nall) = mc
              nold = nall
c...rearrange the original line
              do 31 j = 0,lold(i)-k
                nnum(j+1,i) = nnum(j+k,i)
31              continue
              lold(i) = lold(i)-k+1
              return
            end if
          end if
        do 20 k = 2,lold(i)
          nnum(mc,nall) = nnum(j,i)
20          continue
        end if
c...many small line segments case
      go to 10
    end if
40      continue
60      continue
c
      return
    end
CNAME CKOLIN.FOR
c-----
      subroutine ckolin(nod1,nod2,oline,nold)
c-----
c
c PROGRAM: CKOLIN.FOR
c
c PURPOSE: CHECK NODE IN ALL THE OLD LINE
c
c WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
c MODIFIED BY:                  CHECKED BY:          DATE:
c
c-----
c
c ALGORITHM EXPLANATION: CHECK THROUGH THE ALL OLD LINES TO SEE IF THE
c GIVEN LINE WAS AN EXIST OLD LINE
c
c-----
c
c called by: fnewnd.for -- find the new node
c

```

```

c.....
  include 'para.inc'
  dimension ick(maxsid),itemp(20)
  logical oline
  common/oldln/nnum(20,maxsid*8),lold(maxsid*8),nall
c
  oline=.false.
c
c...check for the discrete line case or partial line case
c
  ntemp = nod1
  icnt = 0
10  continue
  do 40 i = 1,nall
c...check for the inverse special case
  if(nod1.eq.nnum(lold(i),i) .and. nod2.eq.nnum(1,i))then
  do 60 km = 1,lold(i)
  itemp(km) = nnum(km,i)
50  continue
  do 61 km = 1,lold(i)
  nnum(km,i) = itemp(lold(i)-km+1)
61  continue
  end if
c...start normal checking
  if(ntemp.eq.nnum(1,i))then
  icnt = icnt + 1
  ick(icnt) = i
  ntemp = nnum(lold(i),i)
c...regular coincident case and final condition of small lines case
  if(ntemp.eq.nod2)then
  oline = .true.
  go to 60
  end if
c...partial line case
  if(lold(i).gt.2)then
  do 20 k = 2,lold(i)
  if(nod2.eq.nnum(k,i))then
  oline = .true.
  nall = nall + 1
  mc = 1
  nnum(1,nall) = nod1
  do 21 m = 1,icnt-1
  do 21 n = 2,lold(ick(m))
  mc = mc + 1
  nnum(mc,nall) = nnum(n,ick(m))
21  continue
  do 30 j = 2,k
  mc = mc + 1
  nnum(mc,nall) = nnum(j,i)
30  continue
  lold(nall) = mc
  nold = nall
c...restore the original line
  do 31 j = 0,lold(i)-k
  nnum(j+1,i) = nnum(j+k,i)
31  continue
  lold(i) = lold(i)-k+1
  return
  end if
20  continue
  end if
c...many small line segments case
  go to 10
  end if
40  continue
60  continue
  if(oline)then
  if(icnt.gt.1)then
c...store the the lines found into one line
  nall = nall + 1
  nnum(1,nall) = nod1
  jc = 1
  do 90 i = 1,icnt
  do 90 j = 2,lold(ick(i))
  jc = jc + 1
  nnum(jc,nall) = nnum(j,ick(i))
90  continue
  lold(nall) = jc
c...delete those small lines
  do 71 ik = icnt,1,-1

```

```

      do 70 i = ick(ik),nall-1
        lold(i) = lold(i+1)
        do 80 jj = 1,lold(i)
          nnum(jj,i) = nnum(jj,i+1)
80          continue
70          continue
          nall = nall - 1
71          continue
          nold = nall
          else
            nold = ick(icnt)
          end if
        end if
      end if
c
      return
      end
CNAME CKTRAN.FOR
c-----
      subroutine cktran(temp,kt)
c-----
C
C PROGRAM: CKTRAN.FOR
C
C PURPOSE: CHECK THE EXIST TRANSITION NODES
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C
c-----
C
C ALGORITHM EXPLANATION: check, if the node already exist then use it
c for the new surface and delete it from the check list (to save time
c and space), otherwise create a new node.
c
c-----
c
c called by:
C tran1.for -- transition method 1
C tran2.for -- transition method 2
c
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),temp(4)
      common/auto/coor,node,last,unit
      common/tran/ntran(1000),kall
      logical found
c
      found = .false.
      do 10 i = 1,kall
        if(abs(temp(1)-coor(1,ntran(i))).le.0.005 .and.
          $ abs(temp(2)-coor(2,ntran(i))).le.0.005 .and.
          $ abs(temp(3)-coor(3,ntran(i))).le.0.005) then
          kt = ntran(i)
          nod = i
          found = .true.
          go to 20
        end if
10      continue
20      continue
        if(found)then
          do 30 j = nod,kall-1
            ntran(j) = ntran(j+1)
30          continue
            kall = kall - 1
          else
            kall = kall + 1
            ntran(kall) = kt
            last = kt
          end if
        end if
      end if
c
      return
      end
CNAME DEFOR1.FOR
c-----
      subroutine defor1(xo,yo,zo,x,y,z,nodnum)
c-----
C
C PROGRAM: DEFOR1.FOR
C

```



```

C  PURPOSE: DO THE DEFORMATION FOR BOUNDARY LINE
C
C  WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C  MODIFIED BY:          CHECKED BY:          DATE:
C
C -----
C  ALGORITHM EXPLANATION: deform the boundary lines first, then store
c    xo,yo,zo as reference points to map the real nodes back to original
c    volume for auto meshing
c -----
c
c  called: deform.for -- free formed deformation
c -----
c
c  called by: gtsoli.for -- get volume data
c -----
c
c    include 'para.inc'
c    dimension x(maxsid,12),y(maxsid,12),z(maxsid,12),nodnum(12)
c    dimension xo(maxsid,12),yo(maxsid,12),zo(maxsid,12)
c
c  deformed after meshing the cube surfaces
c
c    do 10 j=1,12
c    do 10 i=1,nodnum(j)
c      xi=x(i,j)
c      yi=y(i,j)
c      zi=z(i,j)
c      call deform(xi,yi,zi,xd,yd,zd)
c      x(i,j)=xd
c      y(i,j)=yd
c      z(i,j)=zd
c      xo(i,j)=xd
c      yo(i,j)=yd
c      zo(i,j)=zd
10  continue
c
c    return
c    end
CNAME DRSTOR.FOR
C -----
c    subroutine drstor(xx,yy,zz,nodnum,unit)
C -----
C
C  PROGRAM: DRSTOR.FOR
C
C  PURPOSE:  store the nodal information for plot
C
C  WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C  MODIFIED BY:          CHECKED BY:          DATE:
C
C -----
C  ALGORITHM EXPLANATION: STORE DATA INTO GRAPHIC BUCKET
c -----
c
c  called: deform.for -- do the free formed deformation
c -----
c
c  called by:
c    gtsoli.for -- get the original volume data
c    vmesh.for  -- volume mesh controller
c -----
c
c    include 'para.inc'
c    dimension xx(maxsid,12),yy(maxsid,12),zz(maxsid,12),nodnum(12)
c    dimension x(0:maxdeg,maxid),y(0:maxdeg,maxid),
c    $          z(0:maxdeg,maxid),ideg(maxid),isurf(4,6)
c    common/data/x,y,z,ideg
c    data isurf/1,6,2,5, 1,10,3,9, 10,8,12,6, 2,12,4,11,
c    $          9,7,11,5, 3,8,4,7/
c
c  from the boundary lines segments to get surface boundary nodes
c
c    do 10 k = 1,6

```

```

num=-1
do 20 i=1,4
  line=isurf(i,k)
  ln=nodnum(line)-1
  if(i.eq.1 .or. i.eq.2)then
    do 30 j=1,ln
      num=num+1
      xi=xx(j,line)
      yi=yy(j,line)
      zi=zz(j,line)
      call deform(xi,yi,zi,xd,yd,zd)
      x(num,k)=xd
      y(num,k)=yd
      z(num,k)=zd
30    continue
    else
      do 40 j=ln+1,2,-1
        num=num+1
        xi=xx(j,line)
        yi=yy(j,line)
        zi=zz(j,line)
        call deform(xi,yi,zi,xd,yd,zd)
        x(num,k)=xd
        y(num,k)=yd
        z(num,k)=zd
40    continue
      end if
    continue
    ideg(k) = num
10  continue
c
  return
end

```

CNAME EVENUM.FOR

```

c-----
c      subroutine evenum(dist,nodnum,iflag)
c-----
C
C PROGRAM: EVENUM.FOR
C
C PURPOSE: set the surface close loop to be an even number
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C
C ALGORITHM EXPLANATION: CHECK THE CLOSE LOOP OF A SURFACE IF THE TOTAL
C NUMBER WAS NOT AN EVEN NUMBER THEN ADD 1 TO MAKE IT EVEN
C-----
c
c called by: mshlin.for -- mesh the boundary lines first
c-----
c      dimension nodnum(12),isurf(4,8),line(12),iflag(12),dist(12)
c      logical found
c      data line /1,2,3,4,5,6,7,8,9,10,11,12/
c      data isurf /1,6,2,5, 1,10,3,9, 10,8,12,6, 2,12,4,11,
c      $          5,9,7,11, 3,8,4,7/
c
c increase the largest line segment by one to change the opposit
c odd number to even number.
c
  do 10 i=1,8
    do 20 j = 1,2
      if(mod(nodnum(isurf(j,i)),2) .ne.
c      $          mod(nodnum(isurf(j+2,i)),2))then
c      $          if(line(isurf(j,i)).gt.0 .and.
c      $          line(isurf(j+2,i)).gt.0)then
c      $          if(nodnum(isurf(j,i)).gt.nodnum(isurf(j+2,i)))then
c      $          ic = isurf(j,i)
c      $          else
c      $          ic = isurf(j+2,i)
c      $          end if
c      $          else if(line(isurf(j,i)).gt.0)then
c      $          ic = isurf(j,i)
c      $          else if(line(isurf(j+2,i)).gt.0)then
c      $          ic = isurf(j+2,i)

```

```

                end if
                nodnum(ic) = nodnum(ic)+1
                iflag(ic) = 1
            end if
            line(isurf(j,i)) = -isurf(j,i)
            line(isurf(j+2,i)) = -isurf(j+2,i)
20          continue
10          continue
c
            return
            end

CNAME FNEWND.FOR
c-----
$      subroutine fnewnd(line1,line2,num1,num2,left,loop,nseg,
c-----
c          oline,nold)
c-----
c
c PROGRAM: FNEWND.FOR
c
c PURPOSE: FIND THE NEW NODES
c
c WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
c MODIFIED BY:          CHECKED BY:          DATE:
c-----
c
c ALGORITHM EXPLANATION: FIND THE OLD NODES, OR CREATE NEW NODES
c-----
c
c called: vnewds.for -- find the distance for the deformed line
c          ckolin.for -- check the existence of the old line
c-----
c
c called by: vnewnd.for -- find the new node
c-----
c
c          include 'para.inc'
c          dimension coor(4,maxnod),node(maxsid,12,maxloop),xo(maxsid),
c          $              yo(maxsid),zo(maxsid)
c          real new,newsum
c          logical oline
c          common/auto/coor,node,last,unit
c          common/oldln/nnum(20,maxsid*6),lold(maxsid*6),nall
c
c          oline=.false.
c          num = node(1,left,loop)
c          nod1 = node(num1+1,line1,loop)
c          nod2 = node(num2+1,line2,loop)
c if this line was an already exist old line then return require parameter
c for copy its line data
c          do 30 i = 1,nall
c              if(nod2.eq.nnum(lold(i),i) .and. nod1.eq.nnum(1,i))then
c                  oline = .true.
c                  nold = i
c                  return
c              end if
30          continue
c...check for the discontinue line case or partial line case
c          call ckolin(nod1,nod2,oline,nold)
c          if(oline) return
c
c find the distance for the deformed line
c
c          call vnewds(nod1,nod2,num,xo,yo,zo,dist)
c
c find the require segments (even number in a close polygon) for
c the deformed line
c
c          nseg = nint(dist/unit) + 1
c          if(nseg.lt.2) nseg = 2
c          isum = num1 + num2 + num + nseg
c          if(mod(isum,2).ne.0)then
c              nseg = nseg + 1
c          end if
c
c map this line back to the reference hexahedron
c

```

```

new=dist/(nseg-1)
old=1.0e30
k = 0
newsum=new
i = 0
oldsum=0.0
ratio =0.0
10 continue
  if(newsum.le.oldsum+0.001)then
    diff = newsum - (oldsum-old)
    if(diff.gt.new)diff=new
    ratio = ratio + diff/old/(num-1)
    k = k+1
    $   coor(1,last+k)=coor(1,nod1) + (coor(1,nod2)-coor(1,nod1))
    $   *ratio
    $   coor(2,last+k)=coor(2,nod1) + (coor(2,nod2)-coor(2,nod1))
    $   *ratio
    $   coor(3,last+k)=coor(3,nod1) + (coor(3,nod2)-coor(3,nod1))
    $   *ratio
    coor(4,last+k)=ratio
    newsum = new*(k+1)
  else
    diff = oldsum - new*k
    if(diff.gt.old)diff=old
    ratio = ratio + diff/old/(num-1)
    i = i+1
    $   old = sqrt((xo(i+1)-xo(i))**2 + (yo(i+1) -
    $   yo(i))**2 + (zo(i+1)-zo(i))**2)
    oldsum = oldsum + old
  end if
  if(k.ge.nseg-2) go to 20
  go to 10
20 continue
c
  return
end
CNAME GTSIDE.FOR
-----
C   subroutine gtside(ibeg,next,left,nrit)
-----
C
C PROGRAM: GTSIDE.FOR
C
C PURPOSE: FIND THE OTHER OPPOSITE LINES OF A GIVEN SURFACE
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C MODIFIED BY:                CHECKED BY:           DATE:
C
-----
C ALGORITHM EXPLANATION: from the given opposite lines to find the othe
c opposite lines of that surface
c
-----
c called by:
c   vgtnod.for -- find the best node for split line
c   vnewnd.for -- find the new distance
c
-----
c for x direction
  if(ibeg.eq.1 .and. next.eq.2)then
    left = 5
    nrit = 6
  else if(ibeg.eq.1 .and. next.eq.3)then
    left = 9
    nrit = 10
  else if(ibeg.eq.2 .and. next.eq.4)then
    left = 11
    nrit = 12
  else if(ibeg.eq.3 .and. next.eq.4)then
    left = 7
    nrit = 8
c for y direction
  else if(ibeg.eq.5 .and. next.eq.6)then
    left = 1
    nrit = 2
  else if(ibeg.eq.5 .and. next.eq.7)then
    left = 9
    nrit = 11

```

```

        else if(ibeg.eq.6 .and. next.eq.8)then
            left = 10
            nrit = 12
        else if(ibeg.eq.7 .and. next.eq.8)then
            left = 3
            nrit = 4
c   for z direction
        else if(ibeg.eq.9 .and. next.eq.10)then
            left = 1
            nrit = 3
        else if(ibeg.eq.9 .and. next.eq.11)then
            left = 5
            nrit = 7
        else if(ibeg.eq.11 .and. next.eq.12)then
            left = 2
            nrit = 4
        else if(ibeg.eq.10 .and. next.eq.12)then
            left = 6
            nrit = 8
        end if
c
        return
    end
CNAME LFLAG.FOR
c-----
      subroutine lflag(loop,line1,line2,line3,line4,iflag)
c-----
C
C   PROGRAM: LFLAG.FOR
C
C   PURPOSE: SET FLAG FOR THE LINES
C
C   WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C   MODIFIED BY:                CHECKED BY:           DATE:
C
c-----
C   ALGORITHM EXPLANATION: check the number of nodes on each line at x,y,
c   or z direction and set flag for them
c
c-----
c   called by: lnflag.for -- get flag number for lines
c
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop)
      common/auto/coor,node,last,unit
c
      num1 = node(1,line1,loop)
      num2 = node(1,line2,loop)
      num3 = node(1,line3,loop)
      num4 = node(1,line4,loop)
      if(num1.eq.2 .and. num2.eq.2 .and.
$         num3.eq.2 .and. num4.eq.2)then
          iflag = 0
c
c   for transition case 1.
$   else if(num1.eq.4 .and. num2.eq.2 .and. num3.eq.2
$         .and. num4.eq.2)then
          iflag = 11
$   else if(num2.eq.4 .and. num3.eq.2 .and. num4.eq.2
$         .and. num1.eq.2)then
          iflag = 12
$   else if(num3.eq.4 .and. num4.eq.2 .and. num1.eq.2
$         .and. num2.eq.2)then
          iflag = 13
$   else if(num4.eq.4 .and. num1.eq.2 .and. num2.eq.2
$         .and. num3.eq.2)then
          iflag = 14
c
c   for transition case 2.
$   else if((num1.eq.2 .and. num3.eq.2 .and.
$         mod(num2,2).eq.0 .and. mod(num4,2).eq.0) .or.
$         (num2.eq.2 .and. num4.eq.2 .and.
$         mod(num1,2).eq.0 .and. mod(num3,2).eq.0))then
          iflag = 21
$   else if((num1.eq.2 .and. num2.eq.2 .and.
$         mod(num3,2).eq.0 .and. mod(num4,2).eq.0) .or.
$         (num3.eq.2 .and. num4.eq.2 .and.

```

```

$      mod(num1,2).eq.0 .and. mod(num2,2).eq.0))then
  iflag = 22
c
c for transition case 3.
c   else if(num1.eq.3 .and. num2.eq.3 .and. ((num3.eq.2 .and.
$     mod(num4,2).eq.0).or.(mod(num3,2).eq.0 .and. num4.eq.2)))then
c     iflag = 31
c   else if(num3.eq.3 .and. num4.eq.3 .and. ((num1.eq.2 .and.
$     mod(num2,2).eq.0).or.(mod(num1,2).eq.0 .and. num2.eq.2)))then
c     iflag = 32
c   else if(num1.eq.3 .and. num3.eq.3 .and. ((num2.eq.2 .and.
$     mod(num4,2).eq.0).or.(mod(num2,2).eq.0 .and. num4.eq.2)))then
c     iflag = 33
c   else if(num2.eq.3 .and. num4.eq.3 .and. ((num1.eq.2 .and.
$     mod(num3,2).eq.0).or.(mod(num1,2).eq.0 .and. num3.eq.2)))then
c     iflag = 34
c
c for transition case 4
c   else if((num1.eq.3 .and. mod(num4,2).eq.0 .and. mod(num3,2).eq.0
$     .and. mod(num2,2).eq.0) .or.
c   $$ (num1.eq.3 .and. mod(num2,2).eq.0 .and. mod(num3,2).eq.0
$     .and. mod(num4,2).eq.0) .or.
c   $ (num1.eq.3 .and. mod(num3,2).eq.0 .and. mod(num4,2).eq.0
$     .and. mod(num2,2).eq.0)) then
c     iflag = 41
c   else if((num2.eq.3 .and. mod(num3,2).eq.0 .and. mod(num4,2).eq.0
$     .and. mod(num1,2).eq.0) .or.
c   $ (num2.eq.3 .and. mod(num1,2).eq.0 .and. mod(num4,2).eq.0
$     .and. mod(num3,2).eq.0) .or.
c   $ (num2.eq.3 .and. mod(num4,2).eq.0 .and. mod(num3,2).eq.0
$     .and. mod(num1,2).eq.0)) then
c     iflag = 42
c   else if((num3.eq.3 .and. mod(num2,2).eq.0 .and. mod(num1,2).eq.0
$     .and. mod(num4,2).eq.0) .or.
c   $ (num3.eq.3 .and. mod(num1,2).eq.0 .and. mod(num2,2).eq.0
$     .and. mod(num4,2).eq.0) .or.
c   $ (num3.eq.3 .and. mod(num4,2).eq.0 .and. mod(num1,2).eq.0
$     .and. mod(num2,2).eq.0)) then
c     iflag = 43
c   else if((num4.eq.3 .and. mod(num1,2).eq.0 .and. mod(num2,2).eq.0
$     .and. mod(num3,2).eq.0) .or.
c   $ (num4.eq.3 .and. mod(num2,2).eq.0 .and. mod(num1,2).eq.0
$     .and. mod(num3,2).eq.0) .or.
c   $ (num4.eq.3 .and. mod(num3,2).eq.0 .and. mod(num2,2).eq.0
$     .and. mod(num1,2).eq.0)) then
c     iflag = 44
c
c for regular subdivision
$   else if(num1.ne.2 .and. num3.ne.2 .and.
$     num2.ne.2 .and. num4.ne.2)then
  iflag = 6
c
c unknown situation --- transition case 5
c
  else
    iflag = 5
  end if
c
  return
  end

CNAME LINDIS.FOR
c-----
  subroutine lindis(num,j,x,y,z,maxd,dist)
c-----
C
C PROGRAM: LINDIS.FOR
C
C PURPOSE: find the distance for 12 boundary lines
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:                  CHECKED BY:          DATE:
C
c-----
C ALGORITHM EXPLANATION: USE LINEAR METHOD TO FIND THE DISTANCE
c-----
c
c called by: mshlin.for -- mesh the boundary lines first

```

```

c
c-----
      dimension x(maxd,12),y(maxd,12),z(maxd,12)
c
      sum=0.0
      do 10 i=1,num
        sum = sum+sqrt((x(i+1,j)-x(i,j))**2 +
10      $      (y(i+1,j)-y(i,j))**2 + (z(i+1,j)-z(i,j))**2)
      continue
      dist = sum
c
      return
      end

CNAME LINNOD.FOR
c-----
      subroutine linnod(nx,ny,nz,xg,yg,zg,x,y,z,maxd,nodnum)
c-----
C
C PROGRAM: LINNOD.FOR
C
C PURPOSE: find the nodal data for twelve lines
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 8, 1988
C MODIFIED BY:           CHECKED BY:           DATE:
C
c-----
C ALGORITHM EXPLANATION:
c
c called by: gtsoli.for -- prepare the information for meshing
c
c-----
      dimension x(maxd,12),y(maxd,12),z(maxd,12),nodnum(12)
      dimension xg(8),yg(8),zg(8),ick(24)
      data ick/1,4,5,8,2,3,6,7, 1,2,5,6,4,3,8,7, 1,2,4,3,5,8,8,7/
c
c store nodes on x direction line 1,2,3,4
c
      do 20 m=1,4
        do 10 i=1,nx+1
          ratio=float(i-1)/float(nx)
          x(i,m)=xg(ick(m))+(xg(ick(m+4))-xg(ick(m)))*ratio
          y(i,m)=yg(ick(m))+(yg(ick(m+4))-yg(ick(m)))*ratio
10          z(i,m)=zg(ick(m))+(zg(ick(m+4))-zg(ick(m)))*ratio
          continue
          nodnum(m)=nx+1
20          continue
c
c store nodes on y direction line 5,6,7,8
c
      do 40 m=9,12
        do 30 i=1,ny+1
          ratio=float(i-1)/float(ny)
          x(i,m-4)=xg(ick(m))+(xg(ick(m+4))-xg(ick(m)))*ratio
          y(i,m-4)=yg(ick(m))+(yg(ick(m+4))-yg(ick(m)))*ratio
30          z(i,m-4)=zg(ick(m))+(zg(ick(m+4))-zg(ick(m)))*ratio
          continue
          nodnum(m-4)=ny+1
40          continue
c
c store nodes on z direction line 9,10,11,12
c
      do 60 m=17,20
        do 50 i=1,nz+1
          ratio=float(i-1)/float(nz)
          x(i,m-8)=xg(ick(m))+(xg(ick(m+4))-xg(ick(m)))*ratio
          y(i,m-8)=yg(ick(m))+(yg(ick(m+4))-yg(ick(m)))*ratio
50          z(i,m-8)=zg(ick(m))+(zg(ick(m+4))-zg(ick(m)))*ratio
          continue
          nodnum(m-8)=nz+1
60          continue
c
      return
      end
CNAME LNFLAG.FOR
c-----
      subroutine lnflag(loop,iflagx,iflagy,iflagz)

```

```

C-----
C
C PROGRAM: LNFLAG.FOR
C
C PURPOSE: SET NODAL FLAG FOR X,Y,Z DIRECTION
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C
C ALGORITHM EXPLANATION: check the number of nodes on each line to
c      determind the require flag number for proper processing.
c-----
c
c      called : lflag.for -- set flag for the require direction
c-----
c
c      called by: vmesh.for -- volume mesh controler
c-----
c
c set flag for x direction
c
c      call lflag(loop,1,2,3,4,iflagx)
c
c set flag for y direction
c
c      call lflag(loop,5,6,7,8,iflagy)
c
c set flag for z direction
c
c      call lflag(loop,9,10,11,12,iflagz)
c
c      return
c      end

CNAME MPBACK.FOR
c-----
c      subroutine mpback(nx,ny,nz,xg,yg,zg,dist,xo,yo,zo,x,y,z,w
c      $                      ,nodnum)
c-----
C
C PROGRAM: MPBACK.FOR
C
C PURPOSE: mapped the nodes back to original volume for meshing
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C
C ALGORITHM EXPLANATION: FIND THE PROPER LOCATION OF EACH DEFORMED NODE
C      IN IT'S ORIGINAL CUBE FOR FURTHER PROCESSING
c-----
c
c      called by: gtsoli.for -- prepare the volume information
c-----
c
c      include 'para.inc'
c      real new,newsum
c      dimension x(maxsid,12),y(maxsid,12),z(maxsid,12),w(maxsid,12)
c      dimension xo(maxsid,12),yo(maxsid,12),zo(maxsid,12),
c      $          nodnum(12),dist(12)
c      dimension xg(8),yg(8),zg(8),ick(24)
c      data ick/1,4,5,8,1,2,5,6,1,2,4,3, 2,3,6,7,4,3,8,7,5,6,8,7/
c
c      do 20 j=1,12
c        if(j.le.4)then
c          nm=nx
c        else if(j.gt.8)then
c          nm=nz
c        else
c          nm=ny
c        end if
c        new=dist(j)/(nodnum(j)-1)
c        old=1.0e30

```



```

      k=0
      newsum=0.0
      i=0
      oldsum=0.0
      ratio =0.0
10  continue
      if(newsum.le.oldsum+0.001)then
        diff = newsum - (oldsum-old)
        if(diff.gt.new)diff=new
        ratio = ratio + diff/old/nm
        if(ratio.gt.1.0)ratio=1.0
        k = k+1
        x(k,j)=xg(ick(j))+(xg(ick(j+12))-xg(ick(j)))*ratio
        y(k,j)=yg(ick(j))+(yg(ick(j+12))-yg(ick(j)))*ratio
        z(k,j)=zg(ick(j))+(zg(ick(j+12))-zg(ick(j)))*ratio
        w(k,j) = ratio
        newsum = new*k
      else
        diff = oldsum - new*(k-1)
        if(diff.gt.old)diff=old
        ratio = ratio + diff/old/nm
        i = i+1
        $      old = sqrt((xo(i+1,j)-xo(i,j))**2 + (yo(i+1,j) -
          yo(i,j))**2 + (zo(i+1,j)-zo(i,j))**2)
        oldsum = oldsum + old
      end if
      if(k.ge.nodnum(j)) go to 20
      go to 10
20  continue
c
      return
      end

CNAME MSHLIN.FOR
c-----
      subroutine mshlin(nx,ny,nz,x,y,z,nodnum,unit,dist)
c-----
c
c PROGRAM: MSHLIN.FOR
c
c PURPOSE: MESH THE TWELVE BOUNDARY LINE
c
c WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 8, 1988
c MODIFIED BY:           CHECKED BY:           DATE:
c-----
c
c ALGORITHM EXPLANATION: FIND THE NODES ON BOUNDARY LINE BY THE REQUIRE
c ELEMENT SIZE
c
c-----
c
c called:
c   lindic.for -- find the total distance of a deformed line
c   evenum.for -- set the total node number to be even number
c   renode.for -- if the number is not even then reset the number
c-----
c
c called by: gtsoli.for -- prepare the information for meshing
c-----
c
c   include 'para.inc'
c   dimension x(maxsid,12),y(maxsid,12),z(maxsid,12),nodnum(12)
c   dimension xb(maxsid),yb(maxsid),zb(maxsid),dist(12),iflag(12)
c
c   do 20 j=1,12
c     sum=0.0
c     nnod=0
c     gdis=0.0
c     if(j.le.4)then
c       num=nx
c     else if(j.gt.4 .and. j.le.8)then
c       num=ny
c     else if(j.gt.8)then
c       num=nz
c     end if
c
c find the boundary distances for each line first
c

```

```

      call lindis(num,j,x,y,z,maxsid,dis)
c
c reset the segment of each line after deformed
c
      nseg = nint(dis/unit)
      if(nseg.eq.0) nseg=1
      fdis=dis/nseg
      do 10 i=1,num+1
        if(i.gt.1)then
          delt = sqrt((x(i,j)-x(i-1,j))**2 + (y(i,j)-
$           y(i-1,j))**2 + (z(i,j)-z(i-1,j))**2)
          sum = sum+delt
        else
          delt = 1.e10
        end if
c
c if the final length between two node (delt) is greater than deformed length
c then increase the deformed node number otherwise increase the final node
c number.
11      if(gdis.le.sum+.001)then
          nnod=nnod+1
          if(nnod.gt.maxsid)then
            write(*,*)'ERROR : increase maxsid to',nnod
            return
          end if
          diff=sum-gdis
          xb(nnod) = x(i,j) - (x(i,j)-x(i-1,j))/delt*diff
          yb(nnod) = y(i,j) - (y(i,j)-y(i-1,j))/delt*diff
          zb(nnod) = z(i,j) - (z(i,j)-z(i-1,j))/delt*diff
          gdis=gdis+fdis
          go to 11
        end if
10      continue
        nodnum(j)=nnod
c
c restore the line node data
      do 30 i=1,nnod
        x(i,j)=xb(i)
        y(i,j)=yb(i)
        z(i,j)=zb(i)
30      continue
        dist(j)=dis
20      continue
c
c set the surface close loop is even number to perform hexahedron mesh
c
      call evenum(dist,nodnum,iflag)
      call renode(nx,ny,nz,x,y,z,nodnum,dist,iflag)
c
      return
      end

CNAME NODLIN.FOR
c-----
      subroutine nodlin(xg,yg,zg,nx,ny,nz,x,y,z)
c-----
C
C PROGRAM: NODLIN.FOR
C
C PURPOSE: store the nodes data of twelve lines
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:                  CHECKED BY:          DATE:
c-----
C ALGORITHM EXPLANATION:
c-----
c
c called:
c-----
c
c called by: orimsh.for -- find the original mesh
c-----
      dimension x(nx+1,ny+1,nz+1),y(nx+1,ny+1,nz+1),z(nx+1,ny+1,nz+1)
      dimension xg(8),yg(8),zg(8),ick(8)
      data ick/1,2,4,3,5,6,8,7/

```

```

c
c store the corner point first
c
      m = 0
      do 10 k = 1,nz+1,nz
        do 10 j = 1,ny+1,ny
          do 10 i = 1,nx+1,nx
            m = m+1
            x(i,j,k) = xg(ick(m))
            y(i,j,k) = yg(ick(m))
            z(i,j,k) = zg(ick(m))
10      continue
c
c store nodes into x-direction boundary lines
c
      do 20 k=1,nz+1,nz
        do 20 j=1,ny+1,ny
          do 20 i=2,nx
            ratio=float(i-1)/float(nx)
            x(i,j,k) = x(1,j,k) + (x(nx+1,j,k) - x(1,j,k))*ratio
            y(i,j,k) = y(1,j,k) + (y(nx+1,j,k) - y(1,j,k))*ratio
            z(i,j,k) = z(1,j,k) + (z(nx+1,j,k) - z(1,j,k))*ratio
20      continue
c
c store nodes into y-direction boundary lines
c
      do 30 k=1,nz+1,nz
        do 30 i=1,nx+1,nx
          do 30 j=2,ny
            ratio=float(j-1)/float(ny)
            x(i,j,k) = x(i,1,k) + (x(i,ny+1,k) - x(i,1,k))*ratio
            y(i,j,k) = y(i,1,k) + (y(i,ny+1,k) - y(i,1,k))*ratio
            z(i,j,k) = z(i,1,k) + (z(i,ny+1,k) - z(i,1,k))*ratio
30      continue
c
c store nodes into z-direction boundary lines
c
      do 40 i=1,nx+1,nx
        do 40 j=1,ny+1,ny
          do 40 k=2,nz
            ratio=float(k-1)/float(nz)
            x(i,j,k) = x(i,j,1) + (x(i,j,nz+1) - x(i,j,1))*ratio
            y(i,j,k) = y(i,j,1) + (y(i,j,nz+1) - y(i,j,1))*ratio
            z(i,j,k) = z(i,j,1) + (z(i,j,nz+1) - z(i,j,1))*ratio
40      continue
c
      return
      end
CNAME NODSTO.FOR
c
-----
c
      subroutine nodsto(cname,nx,ny,nz)
c
-----
C
C PROGRAM: NODSTO.FOR
C
C PURPOSE: store all the nodes for movie file display
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 8, 1988
C MODIFIED BY:                CHECKED BY:           DATE:
C
-----
C
C ALGORITHM EXPLANATION:
c
-----
c
c called:
c
-----
c
c called by: orimsh.for -- find the original mesh
c
-----
      include 'para.inc'
      dimension x(0:maxdeg,maxid),y(0:maxdeg,maxid),
      z(0:maxdeg,maxid),ideg(maxid)
      common/data/x,y,z,ideg
      dimension coor(4,maxnod),node(maxsid,12,maxloop),jp(4*maxid)
      common/auto/coor,node,last,unit
      character cname=(*)

```

```

c
  imov = 13
  open(imov,file=cname//'.def',status='new')
c
  np = 1
  ic = 0
  npt = 0
c
c store mesh for the front,back face
c
  ic = 0
  m = (nx+1)*(ny+1)
  do 10 k = 0,m,m
    do 10 j = 1,ny
      do 10 i = 1,nx
        jp(ic+1) = k+(nx+1)-(j-1)+i
        jp(ic+2) = k+(nx+1)-(j-1)+i+1
        jp(ic+3) = k+(nx+1)+j+i+1
        jp(ic+4) = -(k+(nx+1)*j+i)
        ic = ic + 4
10    continue
c
c store mesh for the left,right face
c
  n = (ny+1)*(nz+1)
  do 20 k = 2*m,(2*m+n),n
    do 20 j = 1,ny
      do 20 i = 1,nz
        jp(ic+1) = k+(nz+1)-(j-1)+i
        jp(ic+2) = k+(nz+1)-(j-1)+i+1
        jp(ic+3) = k+(nz+1)+j+i+1
        jp(ic+4) = -(k+(nz+1)*j+i)
        ic = ic + 4
20    continue
c
c store mesh for the top,bottom face
c
  l = (nx+1)*(nz+1)
  do 30 k = 2*(m+n),(2*(m+n)-1),l
    do 30 j = 1,nz
      do 30 i = 1,nx
        jp(ic+1) = k+(nx+1)-(j-1)+i
        jp(ic+2) = k+(nx+1)-(j-1)+i+1
        jp(ic+3) = k+(nx+1)+j+i+1
        jp(ic+4) = -(k+(nx+1)*j+i)
        ic = ic + 4
30    continue
  nedge = ic
  nj = last
  npt = ic/4
c
c store mesh for plot
c
  do 40 id = 1,npt
    do 50 mc = 1,4
      if(mc.eq.4)then
        x(mc-1,id) = coor(1,-jp((id-1)*4+mc))
        y(mc-1,id) = coor(2,-jp((id-1)*4+mc))
        z(mc-1,id) = coor(3,-jp((id-1)*4+mc))
      else
        x(mc-1,id) = coor(1,jp((id-1)*4+mc))
        y(mc-1,id) = coor(2,jp((id-1)*4+mc))
        z(mc-1,id) = coor(3,jp((id-1)*4+mc))
      end if
50    continue
  ideg(id) = 3
40  continue
c
c deform the mesh for first view of the shape
c
  do 60 i = 1,last
    call deform(coor(1,i),coor(2,i),coor(3,i),xo,yo,zo)
    coor(1,i) = xo
    coor(2,i) = yo
    coor(3,i) = zo
60  continue
c
c write movie polygon file
c
  write(imov,150) np,nj,npt,nedge

```

```

        write(imov,150) np,npt
        write(imov,180) ((coor(i,j),i=1,3),j=1,nj)
        write(imov,150) (jp(i),i=1,nedge)
        close(imov)
c
150      format(10i8)
180      format(1p8e12.5)
c
        return
        end
CNAME NODSUR.FOR
c-----
      subroutine nodsur(nx,ny,nz,x,y,z,maxd)
c-----
C
C PROGRAM: NODSUR.FOR
C
C PURPOSE: store the nodes data of six surfaces
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C ALGORITHM EXPLANATION:
c-----
c called:
c-----
c called by: orimsh.for -- find the original mesh
c-----
      include 'para.inc'
      dimension x(maxd,12),y(maxd,12),z(maxd,12)
      dimension coor(4,maxnod),node(maxsid,1,maxloop)
      common/auto/coor,node,last,unit
c
      nod = 0
c
c store nodes on two z direction surfaces
c
      do 10 k = 5,7,2
      do 10 j = 1,ny+1
        do 10 i = 1,nx+1
          nod = nod+1
          ratio = float(i-1)/float(nx)
          coor(1,nod) = x(j,k)+(x(j,k-1)-x(j,k))*ratio
          coor(2,nod) = y(j,k)+(y(j,k-1)-y(j,k))*ratio
          coor(3,nod) = z(j,k)+(z(j,k-1)-z(j,k))*ratio
10      continue
c
c store nodes on two x direction surfaces
c
      do 20 k = 5,8
      do 20 j = 1,ny+1
        do 20 i = 1,nz+1
          nod = nod+1
          ratio = float(i-1)/float(nz)
          coor(1,nod) = x(j,k)+(x(j,k+2)-x(j,k))*ratio
          coor(2,nod) = y(j,k)+(y(j,k+2)-y(j,k))*ratio
          coor(3,nod) = z(j,k)+(z(j,k+2)-z(j,k))*ratio
20      continue
c
c store nodes on two y direction surfaces
c
      do 30 k = 9,11,2
      do 30 j = 1,nz+1
        do 30 i = 1,nx+1
          nod = nod+1
          ratio = float(i-1)/float(nx)
          coor(1,nod) = x(j,k)+(x(j,k+1)-x(j,k))*ratio
          coor(2,nod) = y(j,k)+(y(j,k+1)-y(j,k))*ratio
          coor(3,nod) = z(j,k)+(z(j,k+1)-z(j,k))*ratio
30      continue
      last = nod
c
      return

```

```

      end
CNAME ORIMSH.FOR
C*****
      subroutine orimsh(xg,yg,zg,nx,ny,nz)
C*****
C
C PROGRAM: ORIMSH.FOR
C
C PURPOSE: store the nodes data for original mesh
C
C WRITTEN BY: SSUTA HSU      CHECKED BY:      DATE: OCT. 8, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C
C*****
C ALGORITHM EXPLANATION: store the original mesh data for check the
c difference with automatic meshing
c
C*****
c called:
c   nodlin.for -- get the nodes for lines
c   nodsur.for -- get the nodes for surface
c   nodsto.for -- store all the nodes for graphic bucket
C*****
c called by: gtsoli.for -- prepare the information for meshing
C*****
      dimension x(20,20,20),y(20,20,20),z(20,20,20)
      dimension xg(8),yg(8),zg(8),ick(8)
c
c store the line nodes
c
c   call nodlin(xg,yg,zg,nx,ny,nz,x,y,z)
c
c store the surface nodes
c
c   call nodsur(nx,ny,nz,x,y,z)
c
c store into the mesh for original mesh plot
c
c   call nodsto(nx,ny,nz,x,y,z)
c
c   return
c   end
CNAME PREPAR.FOR
C*****
      subroutine prepar(x,y,z,w,nodnum,unitt)
C*****
C
C PROGRAM: PREPAR.FOR
C
C PURPOSE: store the nodal data for volume meshing
C
C WRITTEN BY: SSUTA HSU      CHECKED BY:      DATE: OCT. 8, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C
C*****
C ALGORITHM EXPLANATION:
c
C*****
c called by: gtsoli.for -- prepare the information for meshing
C*****
      include 'para.inc'
      dimension x(maxsid,12),y(maxsid,12),z(maxsid,12),
      §          w(maxsid,12),nodnum(12)
      dimension coor(4,maxnod),node(maxsid,12,maxloop)
      common/auto/coor,node,last,unit
c
c store the coordinate and weight for each nodes
c
c   k = 0
c   do 10 j = 1,12
c     if(j.le.4)then
c       ibeg = 1

```

```

        iend = nodnum(j)
    else
        ibeg = 2
        iend = nodnum(j)-1
    end if
c
c  get the line nodal information for the first volume
c
        node(1,j,1) = nodnum(j)
        do 20 i = ibeg,iend
            k = k+1
            coor(1,k)=x(i,j)
            coor(2,k)=y(i,j)
            coor(3,k)=z(i,j)
            coor(4,k)=w(i,j)
            node(i+1,j,1) = k
20        continue
c
c  find the first and last node for rest of the lines
c
        if(j.gt.4)then
            do 30 i = 1,nodnum(j),nodnum(j)-1
                do 40 m = 1,4
                    do 40 n = 1,nodnum(m),nodnum(m)-1
                        delx = abs(x(i,j)-x(n,m))
                        dely = abs(y(i,j)-y(n,m))
                        delz = abs(z(i,j)-z(n,m))
                        if(delx.lt.0.001 .and. dely.lt.0.001 .and.
$                            delz.lt.0.001)then
                            node(i+1,j,1) = node(n+1,m,1)
                        end if
40                    continue
30                continue
            end if
10        continue
            last = k
            unit = unitt
c
c  store extra data for smoothing purpose
c
        do 50 j = 1,12
            do 60 i = 1,node(1,j,1)+1
                node(i,j,maxloop) = node(i,j,1)
60            continue
50        continue
c
        return
        end

```

CNAME READSO.FOR

```

c-----
      subroutine readso(fname,nx,ny,nz,unit,xg,yg,zg,ierr)
c-----
C
C  PROGRAM: READSO.FOR
C
C  PURPOSE:  read in the volume data by counter clockwise order
C
C  WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C  MODIFIED BY:          CHECKED BY:          DATE:
C-----
C
C  ALGORITHM EXPLANATION:
c-----
c
c  called by: gtsoli.for -- prepare volume information
c-----
      dimension xg(8),yg(8),zg(8)
      character fname=(*)
c
c  read in the prime volume 8 points
c
      open(11,file=fname//'.sol',status='old',err=10)
      read(11,*,end=99)unit
      read(11,*)(xg(i),yg(i),zg(i),i=1,8)
c
c  find the increments on each direction

```

```

c
  nx = nint(sqrt((xg(1)-xg(2))**2 + (yg(1)-yg(2))**2 +
  $           (zg(1)-zg(2))**2)/unit)
  if(nx.le.1) nx=1
  ny = nint(sqrt((xg(1)-xg(4))**2 + (yg(1)-yg(4))**2 +
  $           (zg(1)-zg(4))**2)/unit)
  if(ny.le.1) ny=1
  nz = nint(sqrt((xg(1)-xg(5))**2 + (yg(1)-yg(5))**2 +
  $           (zg(1)-zg(5))**2)/unit)
  if(nz.le.1) nz=1
  go to 99

c
10  ierr=1
99  close(11)
    return
    end
CNAME RENODE.FOR
c-----
c      subroutine renode(nx,ny,nz,x,y,z,nodnum,dist,iflag)
C-----
C
C PROGRAM: RENODE.FOR
C
C PURPOSE: re-mesh line for newly even number
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C MODIFIED BY:           CHECKED BY:           DATE:
C
C-----
C ALGORITHM EXPLANATION: IF THE CLOSE LOOP WAS AN ADD ONE TO MAKE IT EVEN
C LOOP THEN RE-ARRANGE THE NODAL LOCATION
C
c-----
c
c CALLED BY:MSHLIN.FOR
c-----
c
  include 'para.inc'
  dimension x(maxsid,12),y(maxsid,12),z(maxsid,12)
  dimension xb(maxsid),yb(maxsid),zb(maxsid),
  $          nodnum(12),dist(12),iflag(12)
c
  do 20 j=1,12
  if(iflag(j).eq.1)then
    sum=0.0
    nnod=0
    gdis=0.0
    fdis=dist(j)/(nodnum(j)-1)
    do 10 i=1,nodnum(j)-1
      if(i.gt.1)then
        delt=dist(j)/(nodnum(j)-2)
        sum=delt*(i-1)
      else
        delt = 1.e10
      end if
11    if(gdis.le.sum+.001)then
      nnod=nnod+1
      diff=sum-gdis
      xb(nnod) = x(i,j) - (x(i,j)-x(i-1,j))/delt-diff
      yb(nnod) = y(i,j) - (y(i,j)-y(i-1,j))/delt-diff
      zb(nnod) = z(i,j) - (z(i,j)-z(i-1,j))/delt-diff
      gdis=gdis+fdis
      go to 11
    end if
10    continue
    nodnum(j)=nnod
c
c restore the line node data
  do 30 i=1,nnod
    x(i,j)=xb(i)
    y(i,j)=yb(i)
    z(i,j)=zb(i)
30  continue
  end if
20  continue
c
  return
  end
CNAME SMAREA.FOR

```



```

c-----
c      subroutine smarea(id,coox,cooy,cooz,area)
c-----
C
C PROGRAM: SMAREA.FOR
C
C PURPOSE: FIND THE AREA FOR THE ELEMENT
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C
c-----
C ALGORITHM EXPLANATION: FIND THE AREA BY TRIANGULAR METHOD
c
c called by:
c      smooth.for -- mesh smoother
c
c-----
c      include 'para.inc'
c      common/auto/coor,node,last,unit
c      common/fin/ncon(0:10,maxid)
c      dimension node(maxsid,12,maxloop),coor(4,maxnod),cood(3,20)
c
c...find the center vector
c      coox = (coor(1,ncon(0,id)) + coor(1,ncon(1,id)) +
$           coor(1,ncon(2,id)) + coor(1,ncon(3,id)))/4.0
c      cooy = (coor(2,ncon(0,id)) + coor(2,ncon(1,id)) +
$           coor(2,ncon(2,id)) + coor(2,ncon(3,id)))/4.0
c      cooz = (coor(3,ncon(0,id)) + coor(3,ncon(1,id)) +
$           coor(3,ncon(2,id)) + coor(3,ncon(3,id)))/4.0
c...find the area for that surface
c      a = sqrt((coor(1,ncon(0,id))-coor(1,ncon(1,id)))**2 +
$           (coor(2,ncon(0,id))-coor(2,ncon(1,id)))**2 +
$           (coor(3,ncon(0,id))-coor(3,ncon(1,id)))**2)
c      b = sqrt((coor(1,ncon(0,id))-coor(1,ncon(3,id)))**2 +
$           (coor(2,ncon(0,id))-coor(2,ncon(3,id)))**2 +
$           (coor(3,ncon(0,id))-coor(3,ncon(3,id)))**2)
c      c = sqrt((coor(1,ncon(3,id))-coor(1,ncon(1,id)))**2 +
$           (coor(2,ncon(3,id))-coor(2,ncon(1,id)))**2 +
$           (coor(3,ncon(3,id))-coor(3,ncon(1,id)))**2)
c      d = sqrt((coor(1,ncon(2,id))-coor(1,ncon(1,id)))**2 +
$           (coor(2,ncon(2,id))-coor(2,ncon(1,id)))**2 +
$           (coor(3,ncon(2,id))-coor(3,ncon(1,id)))**2)
c      e = sqrt((coor(1,ncon(3,id))-coor(1,ncon(2,id)))**2 +
$           (coor(2,ncon(3,id))-coor(2,ncon(2,id)))**2 +
$           (coor(3,ncon(3,id))-coor(3,ncon(2,id)))**2)
c      s1 = (a + b + c)/2.0
c      s2 = (c + d + e)/2.0
c      areal = sqrt(s1*(s1-a)*(s1-b)*(s1-c))
c      area2 = sqrt(s2*(s2-d)*(s2-e)*(s2-c))
c      area = areal + area2
c
c      return
c      end
CNAME SMOOTH.FOR
c-----
c      subroutine smooth
c-----
C
C PROGRAM: SMOOTH.FOR
C
C PURPOSE: SMOOTHING THE ELEMENTS FOR BETTER VIEW
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C
c-----
C ALGORITHM EXPLANATION: optimize the node coordinate for best hexahedron
c      element by area pull method( modifed Laplacian method)
c
c called:
c      smarea.for -- find the area for the element
c
c-----
c called by:
c      vmesh.for -- volume mesh controler

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```

c
c*****
c      include 'para.inc'
c      dimension x(0:maxdeg,maxid),y(0:maxdeg,maxid),
$      z(0:maxdeg,maxid),ideg(maxid)
c      common/data/x,y,z,ideg
c      dimension node(maxsid,12,maxloop),coor(4,maxnod),isurf(maxid/4)
c      common/auto/coor,ncon,last,unit
c      common/fin/ncon(0:10,maxid)
c      common/oldln/nnum(20,maxsid*6),lold(maxsid*6),nall
c      common/bsur/isurf,icont
c      logical found

c
c      write(*,*)' Smoothing now please wait...'
c      do 9 itr = 1,2
c      do 10 nod = nnum(1,1),last
c          sumx = 0.0
c          sumy = 0.0
c          sumz = 0.0
c          suma = 0.0
c          found = .false.

c
c      check to see if the node was in boundary(and transition) surface or not
c
c          do 20 i = 1,icont
c          do 30 j = 0,3
c              if(nod.eq.ncon(j,isurf(i)))then
c                  found = .true.
c                  go to 50
c              end if
30          continue
20          continue
c
c      if the node was in boundary surfaces, but not in boundary lines then do
c      the surface smoothing, otherwise do the solid smoothing
c
c      50          continue
c          if(found)then
c              do 60 j = 1,12
c              do 60 i = 1,node(1,j,maxloop)
c                  if(nod.eq.node(i+1,j,maxloop))go to 10
60          continue
c              do 70 id = 1,maxid
c                  if(ncon(0,id).le.0)go to 120
c...check if the element is on the boundary surface or not
c                  if(nod.eq.ncon(0,id) .or. nod.eq.ncon(1,id) .or.
$                  nod.eq.ncon(2,id) .or. nod.eq.ncon(3,id))then
c                      do 90 im = 1,icont
c                          if(id .eq. isurf(im))then
c                              call smarea(id,coox,cooy,cooz,area)
c                              sumx = sumx + area*coox
c                              sumy = sumy + area*cooy
c                              sumz = sumz + area*cooz
c                              suma = suma + area
c                          end if
90                      continue
c                      end if
70                      continue
c                  else
c.... do the volume smoothing
c                  do 100 id = 1,maxid
c                  if(ncon(0,id).le.0)go to 120
c                  do 110 nc = 0,3
c                  if(nod.eq.ncon(nc,id))then
c                      call smarea(id,coox,cooy,cooz,area)
c                      sumx = sumx + area*coox
c                      sumy = sumy + area*cooy
c                      sumz = sumz + area*cooz
c                      suma = suma + area
c                  end if
110                  continue
100                  continue
c                  end if
120                  continue
c                  if(suma.le.0)go to 10
c                  coor(1,nod) = sumx/suma
c                  coor(2,nod) = sumy/suma
c                  coor(3,nod) = sumz/suma
10          continue
9          continue

```

```

c
c store the coordinate into common block for plot
c
      iele = 1
      do while (ideg(iele).ne.0)
        do 130 i = 0,ideg(iele)
          x(i,iele)=coor(1,ncon(i,iele))
          y(i,iele)=coor(2,ncon(i,iele))
          z(i,iele)=coor(3,ncon(i,iele))
130      continue
          iele = iele+1
        end do
c
      return
      end
CNAME SURBND.FOR
c-----
      subroutine surbnd
c-----
      include 'para.inc'
      common/auto/coor,node,last,unit
      common/fin/ncon(0:10,maxid)
      common/oldln/nnum(20,maxsid*6),lold(maxsid*6),nall
      common/tran/ntran(1000),kall
      common/bsur/isurf,icont
      dimension coor(4,maxnod),node(maxsid,12,maxloop),isurf(maxid/4)
      logical lck(0:3),coin(0:3)
c
c check to see if the node was in boundary(and transition) surface or not
c
      icont = 0
      do 10 id = 1,maxid
        if(ncon(0,id).le.0)go to 100
        do 20 i = 0,3
          lck(i) = .false.
20      continue
          do 30 k = 0,3
c...check boundary line nodes
          do 32 i = 1,12
            do 32 j = 2,node(1,i,maxloop)+1
              if(ncon(k,id).eq.node(j,i,maxloop))then
                lck(k) = .true.
                go to 31
              end if
32      continue
c...check boundary surface nodes
          do 40 i = 1,nall
            do 40 j = 1,lold(i)
              if(ncon(k,id).eq.nnum(j,i))then
                lck(k) = .true.
                go to 31
              end if
40      continue
c...check boundary transition nodes
          do 50 i = 1,kall
            if(ncon(k,id).eq.ntran(i))then
              lck(k) = .true.
              go to 31
            end if
50      continue
31      continue
          if(.not.lck(k))go to 10
30      continue
c
c if the all the nodes were on boundary surfaces then store that surface id
c but check for the coincident surface which is a internal surface between
c two boundary shell
c
      if(lck(0) .and. lck(1) .and. lck(2) .and. lck(3))then
        do 80 i = 1,icont
          do 80 im = 0,3
            coin(im) = .false.
            do 90 in = 0,3
              if(ncon(im,id).eq.ncon(in,isurf(i)))then
                coin(im) = .true.
              end if
90      continue
          if(.not.coin(im))go to 80
80      continue
        if(coin(0) .and. coin(1) .and. coin(2) .and. coin(3))then

```

```

          do 70 k = i,icont
             isurf(k) = isurf(k+1)
          continue
          icont = icont -1
          go to 10
        end if
      continue
      icont = icont + 1
      isurf(icont) = id
    end if
  continue
100  return
    end

```

CNAME TRAN1.FOR

```

-----
c      subroutine tran1(loop,iter,jt,iflag)
-----
C
C PROGRAM: TRAN1.FOR
C
C PURPOSE: DO TRANSITION BY #1 METHOD
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C
-----
C ALGORITHM EXPLANATION: do transition on only one line
c
-----
c
c called:
c   cktran  -- check the nodes on the line to see if the node have
c             been used or not
c   tran11  -- do the transition on line 1,5,9.
c   tran12  -- do the transition on line 2,8,10.
c   tran13  -- do the transition on line 3,7,11.
c   tran14  -- do the transition on line 4,3,12.
c   push    -- store the data into new loop
c
-----
c
c called by: transi.for -- transition flag checker
c
-----
c
c   include 'para.inc'
c   dimension coor(4,maxnod),node(maxsid,12,maxloop),kt(6)
c   common/auto/coor,node,last,unit
c
c   if(iflag.eq.11)then
c     jn = jt + 1
c     jm = jt + 2
c     jo = jt + 3
c   else if(iflag.eq.12)then
c     jn = jt + 2
c     jm = jt - 1
c     jo = jt + 1
c   else if(iflag.eq.13)then
c     jn = jt + 1
c     jm = jt - 2
c     jo = jt - 1
c   else if(iflag.eq.14)then
c     jn = jt - 2
c     jm = jt - 1
c     jo = jt - 3
c   end if
c   kt(1) = last + 1
c   njt = node(1,jt,loop)
c   do 10 i = 1,4
c     coor(i,kt(1)) = ((coor(i,node(2,jn,loop)) +
$             coor(i,node(3,jn,loop)))/2.0 +
$             coor(i,node(3,jt,loop)))/2.0
10  continue
c   call cktran(coor(1,kt(1)),kt(1))
c   kt(2) = last + 1
c   do 20 i = 1,4
c     coor(i,kt(2)) = ((coor(i,node(2,jn,loop)) +
$             coor(i,node(3,jn,loop)))/2.0 +

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```

$          coor(i,node(njt,jt,loop)))/2.0
20  continue
    call cktran(coor(1,kt(2)),kt(2))
    kt(3) = last + 1
    do 30 i = 1,4
        coor(i,kt(3)) = ((coor(i,node(2,jm,loop)) +
$          coor(i,node(3,jm,loop)))/2.0 +
$          coor(i,node(3,jt,loop)))/2.0
30  continue
    call cktran(coor(1,kt(3)),kt(3))
    kt(4) = last + 1
    do 40 i = 1,4
        coor(i,kt(4)) = ((coor(i,node(2,jm,loop)) +
$          coor(i,node(3,jm,loop)))/2.0 +
$          coor(i,node(njt,jt,loop)))/2.0
40  continue
    call cktran(coor(1,kt(4)),kt(4))
    kt(5) = last + 1
    last = kt(5)
    do 50 i = 1,4
        coor(i,kt(5)) = ((coor(i,node(2,jo,loop)) +
$          coor(i,node(3,jo,loop)))/2.0 +
$          coor(i,node(3,jt,loop)))/2.0
50  continue
    kt(6) = last + 1
    last = kt(6)
    do 60 i = 1,4
        coor(i,kt(6)) = ((coor(i,node(2,jo,loop)) +
$          coor(i,node(3,jo,loop)))/2.0 +
$          coor(i,node(njt,jt,loop)))/2.0
60  continue
    if(iflag.eq.11)then
        call tran11(loop,jt,jn,jm,jo,kt)
    else if(iflag.eq.12)then
        call tran12(loop,jt,jn,jm,jo,kt)
    else if(iflag.eq.13)then
        call tran13(loop,jt,jn,jm,jo,kt)
    else if(iflag.eq.14)then
        call tran14(loop,jt,jn,jm,jo,kt)
    end if
    call push(loop,iter)
    do 70 i = 1,4
        loop = loop-1
        call push(loop,iter)
70  continue
c
    return
end

```

CNAME TRAN11.FOR

```

C-----
C      subroutine tran11(loop,jt,jn,jm,jo,kt)
C-----
C
C  PROGRAM: TRAN11.FOR
C
C  PURPOSE: do the transition by # 11 method
C
C  WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C  MODIFIED BY:      CHECKED BY:      DATE:
C-----
C
C  ALGORITHM EXPLANATION: transition on line 1,5,9
c
c-----
c
c  called:  gtside -- given the two opposite lines on a surface then
c             get the other two lines.
c
c-----
c
c  called by: tran1.for -- do transition on a line
c
c-----
    include 'para.inc'
    dimension coor(4,maxnod),node(maxsid,12,maxloop),kt(6)
    common/auto/coor,node,last,unit

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c...store them into 5 new loop

```

        loop2 = loop + 1
        loop3 = loop + 2
        loop4 = loop + 3
        loop5 = loop + 4
c
c store the loop 5, the left side
c
    do 10 j = 1,12
        node(1,j,loop5) = 2
10    continue
c store line 1
    node(2,jt,loop5) = node(2,jt,loop)
    node(3,jt,loop5) = node(3,jt,loop)
c store line 2
    node(2,jn,loop5) = node(2,jn,loop)
    node(3,jn,loop5) = kt(1)
c store line 3
    node(2,jm,loop5) = node(2,jm,loop)
    node(3,jm,loop5) = kt(3)
c store line 4
    node(2,jo,loop5) = node(2,jo,loop)
    node(3,jo,loop5) = kt(5)
c store line 5 & 6
    call gtside(jt,jn,left,nrit)
    do 11 i = 1,node(1,left,loop)+1
        node(i,left,loop5) = node(i,left,loop)
11    continue
    node(2,nrit,loop5) = node(3,jt,loop)
    node(3,nrit,loop5) = kt(1)
c store line 7 & 8
    call gtside(jm,jo,left,nrit)
    do 12 i = 1,node(1,left,loop)+1
        node(i,left,loop5) = node(i,left,loop)
12    continue
    node(2,nrit,loop5) = kt(3)
    node(3,nrit,loop5) = kt(5)
c store line 9 & 10
    call gtside(jt,jm,left,nrit)
    do 13 i = 1,node(1,left,loop)+1
        node(i,left,loop5) = node(i,left,loop)
13    continue
    node(2,nrit,loop5) = node(3,jt,loop)
    node(3,nrit,loop5) = kt(3)
c store line 11 & 12
    call gtside(jn,jo,left,nrit)
    do 14 i = 1,node(1,left,loop)+1
        node(i,left,loop5) = node(i,left,loop)
14    continue
    node(2,nrit,loop5) = kt(1)
    node(3,nrit,loop5) = kt(5)
c
c store the loop 4, the back side
c
    do 20 j = 1,12
        node(1,j,loop4) = 2
20    continue
c store line 1
    node(2,jt,loop4) = kt(3)
    node(3,jt,loop4) = kt(4)
c store line 2
    node(2,jn,loop4) = kt(5)
    node(3,jn,loop4) = kt(8)
c store line 3
    do 21 i = 1,node(1,jm,loop)+1
        node(i,jm,loop4) = node(i,jm,loop)
21    continue
c store line 4
    do 22 i = 1,node(1,jo,loop)+1
        node(i,jo,loop4) = node(i,jo,loop)
22    continue
c store line 5 & 6
    call gtside(jt,jn,left,nrit)
    node(2,left,loop4) = kt(3)
    node(3,left,loop4) = kt(5)
    node(2,nrit,loop4) = kt(4)
    node(3,nrit,loop4) = kt(8)
c store line 7 & 8
    call gtside(jm,jo,left,nrit)
    do 23 i = 1,node(1,left,loop)+1
        node(i,left,loop4) = node(i,left,loop)

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```

23     continue
      do 24 i = 1,node(1,nrit,loop)+1
        node(i,nrit,loop4) = node(i,nrit,loop)
24     continue
c store line 9 & 10
      call gtside(jt,jm,left,nrit)
      node(2,left,loop4) = kt(3)
      node(3,left,loop4) = node(node(1,left,loop)+1,left,loop)
      node(2,nrit,loop4) = kt(4)
      node(3,nrit,loop4) = node(node(1,nrit,loop)+1,nrit,loop)
c store line 11 & 12
      call gtside(jn,jo,left,nrit)
      node(2,left,loop4) = kt(5)
      node(3,left,loop4) = node(node(1,left,loop)+1,left,loop)
      node(2,nrit,loop4) = kt(6)
      node(3,nrit,loop4) = node(node(1,nrit,loop)+1,nrit,loop)
c
c store the loop 3, the top side
c
      do 30 j = 1,12
        node(1,j,loop3) = 2
30     continue
c store line 1
      node(2,jt,loop3) = kt(1)
      node(3,jt,loop3) = kt(2)
c store line 2
      do 31 i = 1,node(1,jn,loop)+1
        node(i,jn,loop3) = node(i,jn,loop)
31     continue
c store line 3
      node(2,jm,loop3) = kt(5)
      node(3,jm,loop3) = kt(6)
c store line 4
      do 32 i = 1,node(1,jo,loop)+1
        node(i,jo,loop3) = node(i,jo,loop)
32     continue
c store line 5 & 6
      call gtside(jt,jn,left,nrit)
      node(2,left,loop3) = kt(1)
      node(3,left,loop3) = node(node(1,left,loop)+1,left,loop)
      node(2,nrit,loop3) = kt(2)
      node(3,nrit,loop3) = node(node(1,nrit,loop)+1,nrit,loop)
c store line 7 & 8
      call gtside(jm,jo,left,nrit)
      node(2,left,loop3) = kt(5)
      node(3,left,loop3) = node(node(1,left,loop)+1,left,loop)
      node(2,nrit,loop3) = kt(6)
      node(3,nrit,loop3) = node(node(1,nrit,loop)+1,nrit,loop)
c store line 9 & 10
      call gtside(jt,jm,left,nrit)
      node(2,left,loop3) = kt(1)
      node(3,left,loop3) = kt(5)
      node(2,nrit,loop3) = kt(2)
      node(3,nrit,loop3) = kt(6)
c store line 11 & 12
      call gtside(jn,jo,left,nrit)
      do 33 i = 1,node(1,left,loop)+1
        node(i,left,loop3) = node(i,left,loop)
33     continue
      do 34 i = 1,node(1,nrit,loop)+1
        node(i,nrit,loop3) = node(i,nrit,loop)
34     continue
c
c store the loop 2, the middle bottom side
c
      do 40 j = 2,12
        node(1,j,loop2) = 2
40     continue
c store line 1
      num = node(1,jt,loop)
      node(1,jt,loop2) = num-2
      do 50 j = 1,num-2
        node(j+1,jt,loop2) = node(j+2,jt,loop)
50     continue
c store line 2
      node(2,jn,loop2) = kt(1)
      node(3,jn,loop2) = kt(2)
c store line 3
      node(2,jm,loop2) = kt(3)
      node(3,jm,loop2) = kt(4)

```

```

c store line 4
  node(2,jo,loop2) = kt(5)
  node(3,jo,loop2) = kt(8)
c store line 5 & 8
  call gtside(jt,jn,left,nrit)
  node(2,left,loop2) = node(3,jt,loop)
  node(3,left,loop2) = kt(1)
  node(2,nrit,loop2) = node(num,jt,loop)
  node(3,nrit,loop2) = kt(2)
c store line 7 & 8
  call gtside(jm,jo,left,nrit)
  node(2,left,loop2) = kt(3)
  node(3,left,loop2) = kt(5)
  node(2,nrit,loop2) = kt(4)
  node(3,nrit,loop2) = kt(8)
c store line 9 & 10
  call gtside(jt,jm,left,nrit)
  node(2,left,loop2) = node(3,jt,loop)
  node(3,left,loop2) = kt(3)
  node(2,nrit,loop2) = node(num,jt,loop)
  node(3,nrit,loop2) = kt(4)
c store line 11 & 12
  call gtside(jn,jo,left,nrit)
  node(2,left,loop2) = kt(1)
  node(3,left,loop2) = kt(5)
  node(2,nrit,loop2) = kt(2)
  node(3,nrit,loop2) = kt(8)

c
c store the loop 1, the right side
c
c store line 1
  node(1,jt,loop) = 2
  node(2,jt,loop) = node(num,jt,loop)
  node(3,jt,loop) = node(num+1,jt,loop)
c store line 2, skip 2nd node
  node(1,jn,loop) = 2
  node(2,jn,loop) = kt(2)
c store line 3, skip 2nd node
  node(1,jm,loop) = 2
  node(2,jm,loop) = kt(4)
c store line 4, skip 2nd node
  node(1,jo,loop) = 2
  node(2,jo,loop) = kt(8)
c store line 5, skip 8
  call gtside(jt,jn,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = node(num,jt,loop)
  node(3,left,loop) = kt(2)
c store line 7, skip 8
  call gtside(jm,jo,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(4)
  node(3,left,loop) = kt(8)
c store line 9, skip 10
  call gtside(jt,jm,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = node(num,jt,loop)
  node(3,left,loop) = kt(4)
c store line 11, skip 12
  call gtside(jn,jo,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(2)
  node(3,left,loop) = kt(8)

c
  return
end

```

CNAME TRAN12.FOR

```

C-----
C      subroutine tran12(loop,jt,jn,jm,jo,kt)
C-----
C
C PROGRAM: TRAN12.FOR
C
C PURPOSE: do the transition by # 12 method
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 8, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C-----

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```

C
C ALGORITHM EXPLANATION: transition on line 2,8,10
c
c-----
c
c called:  gtside -- given the two opposite lines on a surface then
c           get the other two lines.
c
c-----
c
c called by: tran1.for -- #1 transition controler
c
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),kt(6)
      common/auto/coor,node,last,unit

c...store them into 5 new loop
      loop2 = loop + 1
      loop3 = loop + 2
      loop4 = loop + 3
      loop5 = loop + 4
c
c store the loop 5, the left side
c
      do 10 j = 1,12
          node(1,j,loop5) = 2
10      continue
c store line 1
      node(2,jm,loop5) = node(2,jm,loop)
      node(3,jm,loop5) = kt(3)
c store line 2
      node(2,jt,loop5) = node(2,jt,loop)
      node(3,jt,loop5) = node(3,jt,loop)
c store line 3
      node(2,jo,loop5) = node(2,jo,loop)
      node(3,jo,loop5) = kt(5)
c store line 4
      node(2,jn,loop5) = node(2,jn,loop)
      node(3,jn,loop5) = kt(1)
c store line 5 & 6
      call gtside(jm,jt,left,nrit)
      do 11 i = 1,node(1,left,loop)+1
          node(i,left,loop5) = node(i,left,loop)
11      continue
      node(2,nrit,loop5) = kt(3)
      node(3,nrit,loop5) = node(3,jt,loop)
c store line 7 & 8
      call gtside(jo,jn,left,nrit)
      do 12 i = 1,node(1,left,loop)+1
          node(i,left,loop5) = node(i,left,loop)
12      continue
      node(2,nrit,loop5) = kt(5)
      node(3,nrit,loop5) = kt(1)
c store line 9 & 10
      call gtside(jm,jo,left,nrit)
      do 13 i = 1,node(1,left,loop)+1
          node(i,left,loop5) = node(i,left,loop)
13      continue
      node(2,nrit,loop5) = kt(3)
      node(3,nrit,loop5) = kt(5)
c store line 11 & 12
      call gtside(jt,jn,left,nrit)
      do 14 i = 1,node(1,left,loop)+1
          node(i,left,loop5) = node(i,left,loop)
14      continue
      node(2,nrit,loop5) = node(3,jt,loop)
      node(3,nrit,loop5) = kt(1)
c
c store the loop 4, the back side
c
      do 20 j = 1,12
          node(1,j,loop4) = 2
20      continue
c store line 1
      node(2,jm,loop4) = kt(5)
      node(3,jm,loop4) = kt(6)
c store line 2
      node(2,jt,loop4) = kt(1)
      node(3,jt,loop4) = kt(2)

```

```

c store line 3
  do 21 i = 1,node(1,jo,loop)+1
    node(i,jo,loop4) = node(i,jo,loop)
21  continue
c store line 4
  do 22 i = 1,node(1,jn,loop)+1
    node(i,jn,loop4) = node(i,jn,loop)
22  continue
c store line 5 & 6
  call gtside(jm,jt,left,nrit)
  node(2,left,loop4) = kt(5)
  node(3,left,loop4) = kt(1)
  node(2,nrit,loop4) = kt(6)
  node(3,nrit,loop4) = kt(2)
c store line 7 & 8
  call gtside(jo,jn,left,nrit)
  do 23 i = 1,node(1,left,loop)+1
    node(i,left,loop4) = node(i,left,loop)
23  continue
  do 24 i = 1,node(1,nrit,loop)+1
    node(i,nrit,loop4) = node(i,nrit,loop)
24  continue
c store line 9 & 10
  call gtside(jm,jo,left,nrit)
  node(2,left,loop4) = kt(5)
  node(3,left,loop4) = node(node(1,left,loop)+1,left,loop)
  node(2,nrit,loop4) = kt(6)
  node(3,nrit,loop4) = node(node(1,nrit,loop)+1,nrit,loop)
c store line 11 & 12
  call gtside(jt,jn,left,nrit)
  node(2,left,loop4) = kt(1)
  node(3,left,loop4) = node(node(1,left,loop)+1,left,loop)
  node(2,nrit,loop4) = kt(2)
  node(3,nrit,loop4) = node(node(1,nrit,loop)+1,nrit,loop)
c
c store the loop 3, the bottom side
c
  do 30 j = 1,12
    node(1,j,loop3) = 2
30  continue
c store line 1
  do 31 i = 1,node(1,jm,loop)+1
    node(i,jm,loop3) = node(i,jm,loop)
31  continue
c store line 2
  node(2,jt,loop3) = kt(3)
  node(3,jt,loop3) = kt(4)
c store line 3
  do 32 i = 1,node(1,jo,loop)+1
    node(i,jo,loop3) = node(i,jo,loop)
32  continue
c store line 4
  node(2,jn,loop3) = kt(5)
  node(3,jn,loop3) = kt(6)
c store line 5 & 6
  call gtside(jm,jt,left,nrit)
  node(2,left,loop3) = node(2,left,loop)
  node(3,left,loop3) = kt(3)
  node(2,nrit,loop3) = node(2,nrit,loop)
  node(3,nrit,loop3) = kt(4)
c store line 7 & 8
  call gtside(jo,jn,left,nrit)
  node(2,left,loop3) = node(2,left,loop)
  node(3,left,loop3) = kt(5)
  node(2,nrit,loop3) = node(2,nrit,loop)
  node(3,nrit,loop3) = kt(6)
c store line 9 & 10
  call gtside(jm,jo,left,nrit)
  do 33 i = 1,node(1,left,loop)+1
    node(i,left,loop3) = node(i,left,loop)
33  continue
  do 34 i = 1,node(1,nrit,loop)+1
    node(i,nrit,loop3) = node(i,nrit,loop)
34  continue
c store line 11 & 12
  call gtside(jt,jn,left,nrit)
  node(2,left,loop3) = kt(3)
  node(3,left,loop3) = kt(5)
  node(2,nrit,loop3) = kt(4)
  node(3,nrit,loop3) = kt(6)

```

```

c
c store the loop 2, the middle top side
c
  do 40 j = 1,12
    node(1,j,loop2) = 2
40  continue
c store line 1
  node(2,jm,loop2) = kt(3)
  node(3,jm,loop2) = kt(4)
c store line 2
  num = node(1,jt,loop)
  node(1,jt,loop2) = num-2
  do 50 j = 1,num-2
    node(j+1,jt,loop2) = node(j-2,jt,loop)
50  continue
c store line 3
  node(2,jo,loop2) = kt(5)
  node(3,jo,loop2) = kt(6)
c store line 4
  node(2,jn,loop2) = kt(1)
  node(3,jn,loop2) = kt(2)
c store line 5 & 6
  call gtside(jm,jt,left,nrit)
  node(2,left,loop2) = kt(3)
  node(3,left,loop2) = node(3,jt,loop)
  node(2,nrit,loop2) = kt(4)
  node(3,nrit,loop2) = node(num,jt,loop)
c store line 7 & 8
  call gtside(jo,jn,left,nrit)
  node(2,left,loop2) = kt(5)
  node(3,left,loop2) = kt(1)
  node(2,nrit,loop2) = kt(6)
  node(3,nrit,loop2) = kt(2)
c store line 9 & 10
  call gtside(jm,jo,left,nrit)
  node(2,left,loop2) = kt(3)
  node(3,left,loop2) = kt(5)
  node(2,nrit,loop2) = kt(4)
  node(3,nrit,loop2) = kt(6)
c store line 11 & 12
  call gtside(jt,jn,left,nrit)
  node(2,left,loop2) = node(3,jt,loop)
  node(3,left,loop2) = kt(1)
  node(2,nrit,loop2) = node(num,jt,loop)
  node(3,nrit,loop2) = kt(2)
c
c store the loop 1, the right side
c
c store line 1
  node(1,jm,loop) = 2
  node(2,jm,loop) = kt(4)
c store line 2
  node(1,jt,loop) = 2
  node(2,jt,loop) = node(num,jt,loop)
  node(3,jt,loop) = node(num+1,jt,loop)
c store line 3
  node(1,jo,loop) = 2
  node(2,jo,loop) = kt(6)
c store line 4
  node(1,jn,loop) = 2
  node(2,jn,loop) = kt(2)
c store line 5, skip 6
  call gtside(jm,jt,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(4)
  node(3,left,loop) = node(num,jt,loop)
c store line 7, skip 8
  call gtside(jo,jn,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(6)
  node(3,left,loop) = kt(2)
c store line 9, skip 10
  call gtside(jm,jo,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(4)
  node(3,left,loop) = kt(6)
c store line 11, skip 12
  call gtside(jt,jn,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = node(num,jt,loop)

```

```

      node(3,left,loop) = kt(2)
c
      return
      end

CNAME TRAN13.FOR
C-----
      subroutine tran13(loop,jt,jn,jm,jo,kt)
C-----
C
C PROGRAM: TRAN13.FOR
C
C PURPOSE: do the transition by # 13 method
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C ALGORITHM EXPLANATION: transition on line 3,7,11
c
c-----
c called:  gtside -- given the two opposite lines on a surface then
c           get the other two lines.
c-----
c called by: tran1.for -- #1 transition controller
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),kt(8)
      common/auto/coor,node,last,unit

c...store them into 5 new loop
      loop2 = loop + 1
      loop3 = loop + 2
      loop4 = loop + 3
      loop5 = loop + 4
c
c store the loop 5, the left side
c
      do 10 j = 1,12
         node(1,j,loop5) = 2
10      continue
c store line 1
      node(2,jm,loop5) = node(2,jm,loop)
      node(3,jm,loop5) = kt(3)
c store line 2
      node(2,jo,loop5) = node(2,jo,loop)
      node(3,jo,loop5) = kt(5)
c store line 3
      node(2,jt,loop5) = node(2,jt,loop)
      node(3,jt,loop5) = node(3,jt,loop)
c store line 4
      node(2,jn,loop5) = node(2,jn,loop)
      node(3,jn,loop5) = kt(1)
c store line 5 & 8
      call gtside(jm,jo,left,nrit)
      do 11 i = 1,node(1,left,loop)+1
         node(i,left,loop5) = node(i,left,loop)
11      continue
      node(2,nrit,loop5) = kt(3)
      node(3,nrit,loop5) = kt(5)
c store line 7 & 8
      call gtside(jt,jn,left,nrit)
      do 12 i = 1,node(1,left,loop)+1
         node(i,left,loop5) = node(i,left,loop)
12      continue
      node(2,nrit,loop5) = node(3,jt,loop)
      node(3,nrit,loop5) = kt(1)
c store line 9 & 10
      call gtside(jm,jt,left,nrit)
      do 13 i = 1,node(1,left,loop)+1
         node(i,left,loop5) = node(i,left,loop)
13      continue
      node(2,nrit,loop5) = kt(3)
      node(3,nrit,loop5) = node(3,jt,loop)
c store line 11 & 12

```

```

        call gtside(jo,jn,left,nrit)
        do 14 i = 1,node(1,left,loop)+1
            node(i,left,loop5) = node(i,left,loop)
14      continue
        node(2,nrit,loop5) = kt(5)
        node(3,nrit,loop5) = kt(1)
c
c store the loop 4, the front side
c
        do 20 j = 1,12
            node(1,j,loop4) = 2
20      continue
c store line 1,2
        do 21 i = 1,node(1,jm,loop)+1
            node(i,jm,loop4) = node(i,jm,loop)
21      continue
        do 22 i = 1,node(1,jo,loop)+1
            node(i,jo,loop4) = node(i,jo,loop)
22      continue
c store line 3
        node(2,jt,loop4) = kt(3)
        node(3,jt,loop4) = kt(4)
c store line 4
        node(2,jn,loop4) = kt(5)
        node(3,jn,loop4) = kt(6)
c store line 5 & 6
        call gtside(jm,jo,left,nrit)
        do 23 i = 1,node(1,left,loop)+1
            node(i,left,loop4) = node(i,left,loop)
23      continue
        do 24 i = 1,node(1,nrit,loop)+1
            node(i,nrit,loop4) = node(i,nrit,loop)
24      continue
c store line 7 & 8
        call gtside(jt,jn,left,nrit)
        node(2,left,loop4) = kt(3)
        node(3,left,loop4) = kt(5)
        node(2,nrit,loop4) = kt(4)
        node(3,nrit,loop4) = kt(6)
c store line 9 & 10
        call gtside(jm,jt,left,nrit)
        node(2,left,loop4) = node(2,left,loop)
        node(3,left,loop4) = kt(3)
        node(2,nrit,loop4) = node(2,nrit,loop)
        node(3,nrit,loop4) = kt(4)
c store line 11 & 12
        call gtside(jo,jn,left,nrit)
        node(2,left,loop4) = node(2,left,loop)
        node(3,left,loop4) = kt(5)
        node(2,nrit,loop4) = node(2,nrit,loop)
        node(3,nrit,loop4) = kt(6)
c
c store the loop 3, the top side
c
        do 30 j = 1,12
            node(1,j,loop3) = 2
30      continue
c store line 1
        node(2,jm,loop3) = kt(5)
        node(3,jm,loop3) = kt(6)
c store line 2
        do 31 i = 1,node(1,jo,loop)+1
            node(i,jo,loop3) = node(i,jo,loop)
31      continue
c store line 3
        node(2,jt,loop3) = kt(1)
        node(3,jt,loop3) = kt(2)
c store line 4
        do 32 i = 1,node(1,jn,loop)+1
            node(i,jn,loop3) = node(i,jn,loop)
32      continue
c store line 5 & 6
        call gtside(jm,jo,left,nrit)
        node(2,left,loop3) = kt(5)
        node(3,left,loop3) = node(node(1,left,loop)+1,left,loop)
        node(2,nrit,loop3) = kt(6)
        node(3,nrit,loop3) = node(node(1,nrit,loop)+1,nrit,loop)
c store line 7 & 8
        call gtside(jt,jn,left,nrit)
        node(2,left,loop3) = kt(1)

```

```

node(3, left, loop3) = node(node(1, left, loop)+1, left, loop)
node(2, nrit, loop3) = kt(2)
node(3, nrit, loop3) = node(node(1, nrit, loop)+1, nrit, loop)
c store line 9 & 10
call gtside(jm, jt, left, nrit)
node(2, left, loop3) = kt(5)
node(3, left, loop3) = kt(1)
node(2, nrit, loop3) = kt(8)
node(3, nrit, loop3) = kt(2)
c store line 11 & 12
call gtside(jo, jn, left, nrit)
do 33 i = 1, node(1, left, loop)+1
  node(i, left, loop3) = node(i, left, loop)
33 continue
do 34 i = 1, node(1, nrit, loop)+1
  node(i, nrit, loop3) = node(i, nrit, loop)
34 continue
c
c store the loop 2, the middle bottom side
c
do 40 j = 1, 12
  node(1, j, loop2) = 2
40 continue
c store line 1
node(2, jm, loop2) = kt(3)
node(3, jm, loop2) = kt(4)
c store line 2
node(2, jo, loop2) = kt(5)
node(3, jo, loop2) = kt(8)
c store line 3
num = node(1, jt, loop)
node(1, jt, loop2) = num - 2
do 50 j = 1, num - 2
  node(j+1, jt, loop2) = node(j+2, jt, loop)
50 continue
c store line 4
node(2, jn, loop2) = kt(1)
node(3, jn, loop2) = kt(2)
c store line 5 & 6
call gtside(jm, jo, left, nrit)
node(2, left, loop2) = kt(3)
node(3, left, loop2) = kt(5)
node(2, nrit, loop2) = kt(4)
node(3, nrit, loop2) = kt(8)
c store line 7 & 8
call gtside(jt, jn, left, nrit)
node(2, left, loop2) = node(3, jt, loop)
node(3, left, loop2) = kt(1)
node(2, nrit, loop2) = node(num, jt, loop)
node(3, nrit, loop2) = kt(2)
c store line 9 & 10
call gtside(jm, jt, left, nrit)
node(2, left, loop2) = kt(3)
node(3, left, loop2) = node(3, jt, loop)
node(2, nrit, loop2) = kt(4)
node(3, nrit, loop2) = node(num, jt, loop)
c store line 11 & 12
call gtside(jo, jn, left, nrit)
node(2, left, loop2) = kt(5)
node(3, left, loop2) = kt(1)
node(2, nrit, loop2) = kt(8)
node(3, nrit, loop2) = kt(2)
c
c store the loop 1, the right side
c
c store line 1
node(1, jm, loop) = 2
node(2, jm, loop) = kt(4)
c store line 2
node(1, jo, loop) = 2
node(2, jo, loop) = kt(8)
c store line 3
node(1, jt, loop) = 2
node(2, jt, loop) = node(num, jt, loop)
node(3, jt, loop) = node(num+1, jt, loop)
c store line 4
node(1, jn, loop) = 2
node(2, jn, loop) = kt(2)
c store line 5 , skip 6
call gtside(jm, jo, left, nrit)

```

```

        node(1,left,loop) = 2
        node(2,left,loop) = kt(4)
        node(3,left,loop) = kt(8)
c store line 7, skip 8
        call gtside(jt,jn,left,nrit)
        node(1,left,loop) = 2
        node(2,left,loop) = node(num,jt,loop)
        node(3,left,loop) = kt(2)
c store line 9, skip 10
        call gtside(jm,jt,left,nrit)
        node(1,left,loop) = 2
        node(2,left,loop) = kt(4)
        node(3,left,loop) = node(num,jt,loop)
c store line 11, skip 12
        call gtside(jo,jn,left,nrit)
        node(1,left,loop) = 2
        node(2,left,loop) = kt(8)
        node(3,left,loop) = kt(2)
c
        return
        end

```

CNAME TRAN14.FOR

```

C-----
      subroutine tran14 (loop,jt,jn,jm,jo,kt)
C-----
C
C PROGRAM: TRAN14.FOR
C
C PURPOSE: do the transition by # 14 method
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 8, 1988
C MODIFIED BY:           CHECKED BY:           DATE:
C-----
C ALGORITHM EXPLANATION: transition on line 4 8,12
c
c-----
c called: gtside -- given the two opposite lines on a surface then
c get the other two lines.
c-----
c called by: tran1.for -- # 1 transition controller
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),kt(8)
      common/auto/coor,node,last,unit

c...store them into 5 new loop
      loop2 = loop + 1
      loop3 = loop + 2
      loop4 = loop + 3
      loop5 = loop + 4
c
c store the loop 5, the left side
c
      do 10 j = 1,12
          node(1,j,loop5) = 2
10 continue
c store line 1
      node(2,jo,loop5) = node(2,jo,loop)
      node(3,jo,loop5) = kt(5)
c store line 2
      node(2,jn,loop5) = node(2,jn,loop)
      node(3,jn,loop5) = kt(1)
c store line 3
      node(2,jm,loop5) = node(2,jm,loop)
      node(3,jm,loop5) = kt(3)
c store line 4
      node(2,jt,loop5) = node(2,jt,loop)
      node(3,jt,loop5) = node(3,jt,loop)
c store line 5 & 8
      call gtside(jo,jn,left,nrit)
      do 11 i = 1,node(1,left,loop)+1
          node(i,left,loop5) = node(i,left,loop)
11 continue

```

```

        node(2,nrit,loop5) = kt(5)
        node(3,nrit,loop5) = kt(1)
c store line 7 & 8
        call gtside(jm,jt,left,nrit)
        do 12 i = 1,node(1,left,loop)+1
            node(i,left,loop5) = node(i,left,loop)
12         continue
            node(2,nrit,loop5) = kt(3)
            node(3,nrit,loop5) = node(3,jt,loop)
c store line 9 & 10
        call gtside(jo,jm,left,nrit)
        do 13 i = 1,node(1,left,loop)+1
            node(i,left,loop5) = node(i,left,loop)
13         continue
            node(2,nrit,loop5) = kt(5)
            node(3,nrit,loop5) = kt(3)
c store line 11 & 12
        call gtside(jn,jt,left,nrit)
        do 14 i = 1,node(1,left,loop)+1
            node(i,left,loop5) = node(i,left,loop)
14         continue
            node(2,nrit,loop5) = kt(1)
            node(3,nrit,loop5) = node(3,jt,loop)
c
c store the loop 4, the front side
c
        do 20 j = 1,12
            node(1,j,loop4) = 2
20         continue
c store line 1
        do 21 i = 1,node(1,jo,loop)+1
            node(i,jo,loop4) = node(i,jo,loop)
21         continue
c store line 2
        do 22 i = 1,node(1,jn,loop)+1
            node(i,jn,loop4) = node(i,jn,loop)
22         continue
c store line 3
        node(2,jm,loop4) = kt(5)
        node(3,jm,loop4) = kt(8)
c store line 4
        node(2,jt,loop4) = kt(1)
        node(3,jt,loop4) = kt(2)
c store line 5 & 6
        call gtside(jo,jn,left,nrit)
        do 23 i = 1,node(1,left,loop)+1
            node(i,left,loop4) = node(i,left,loop)
23         continue
            do 24 i = 1,node(1,nrit,loop)+1
                node(i,nrit,loop4) = node(i,nrit,loop)
24         continue
c store line 7 & 8
        call gtside(jm,jt,left,nrit)
        node(2,left,loop4) = kt(5)
        node(3,left,loop4) = kt(1)
        node(2,nrit,loop4) = kt(8)
        node(3,nrit,loop4) = kt(2)
c store line 9 & 10
        call gtside(jo,jm,left,nrit)
        node(2,left,loop4) = node(2,left,loop)
        node(3,left,loop4) = kt(5)
        node(2,nrit,loop4) = node(2,nrit,loop)
        node(3,nrit,loop4) = kt(8)
c store line 11 & 12
        call gtside(jn,jt,left,nrit)
        node(2,left,loop4) = node(2,left,loop)
        node(3,left,loop4) = kt(1)
        node(2,nrit,loop4) = node(2,nrit,loop)
        node(3,nrit,loop4) = kt(2)
c
c store the loop 3, the bottom side
c
        do 30 j = 1,12
            node(1,j,loop3) = 2
30         continue
c store line 1
        do 31 i = 1,node(1,jo,loop)+1
            node(i,jo,loop3) = node(i,jo,loop)
31         continue
c store line 2

```



```

        node(2,jn,loop3) = kt(5)
        node(3,jn,loop3) = kt(6)
c store line 3
    do 32 i = 1,node(1,jm,loop)+1
        node(i,jm,loop3) = node(i,jm,loop)
32    continue
c store line 4
    node(2,jt,loop3) = kt(3)
    node(3,jt,loop3) = kt(4)
c store line 5 & 6
    call gtside(jo,jn,left,nrit)
    node(2,left,loop3) = node(2,left,loop)
    node(3,left,loop3) = kt(5)
    node(2,nrit,loop3) = node(2,nrit,loop)
    node(3,nrit,loop3) = kt(6)
c store line 7 & 8
    call gtside(jm,jt,left,nrit)
    node(2,left,loop3) = node(2,left,loop)
    node(3,left,loop3) = kt(3)
    node(2,nrit,loop3) = node(2,nrit,loop)
    node(3,nrit,loop3) = kt(4)
c store line 9 & 10
    call gtside(jo,jm,left,nrit)
    do 33 i = 1,node(1,left,loop)+1
        node(i,left,loop3) = node(i,left,loop)
33    continue
    do 34 i = 1,node(1,nrit,loop)+1
        node(i,nrit,loop3) = node(i,nrit,loop)
34    continue
c store line 11 & 12
    call gtside(jn,jt,left,nrit)
    node(2,left,loop3) = kt(5)
    node(3,left,loop3) = kt(3)
    node(2,nrit,loop3) = kt(6)
    node(3,nrit,loop3) = kt(4)
c
c store the loop 2, the middle top side
c
    do 40 j = 1,12
        node(1,j,loop2) = 2
40    continue
c store line 1
    node(2,jo,loop2) = kt(5)
    node(3,jo,loop2) = kt(6)
c store line 2
    node(2,jn,loop2) = kt(1)
    node(3,jn,loop2) = kt(2)
c store line 3
    node(2,jm,loop2) = kt(3)
    node(3,jm,loop2) = kt(4)
c store line 4
    num = node(1,jt,loop)
    node(1,jt,loop2) = num-2
    do 50 j = 1,num-2
        node(j+1,jt,loop2) = node(j+2,jt,loop)
50    continue
c store line 5 & 6
    call gtside(jo,jn,left,nrit)
    node(2,left,loop2) = kt(5)
    node(3,left,loop2) = kt(1)
    node(2,nrit,loop2) = kt(6)
    node(3,nrit,loop2) = kt(2)
c store line 7 & 8
    call gtside(jm,jt,left,nrit)
    node(2,left,loop2) = kt(3)
    node(3,left,loop2) = node(3,jt,loop)
    node(2,nrit,loop2) = kt(4)
    node(3,nrit,loop2) = node(num,jt,loop)
c store line 9 & 10
    call gtside(jo,jm,left,nrit)
    node(2,left,loop2) = kt(5)
    node(3,left,loop2) = kt(3)
    node(2,nrit,loop2) = kt(6)
    node(3,nrit,loop2) = kt(4)
c store line 11 & 12
    call gtside(jn,jt,left,nrit)
    node(2,left,loop2) = kt(1)
    node(3,left,loop2) = node(3,jt,loop)
    node(2,nrit,loop2) = kt(2)
    node(3,nrit,loop2) = node(num,jt,loop)

```

```

c
c store the loop 1, the right side
c
c store line 1
  node(1,jo,loop) = 2
  node(2,jo,loop) = kt(8)
c store line 2
  node(1,jn,loop) = 2
  node(2,jn,loop) = kt(2)
c store line 3
  node(1,jm,loop) = 2
  node(2,jm,loop) = kt(4)
c store line 4
  node(1,jt,loop) = 2
  node(2,jt,loop) = node(num,jt,loop)
  node(3,jt,loop) = node(num+1,jt,loop)
c store line 5, skip 6
  call gtside(jo,jn,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(8)
  node(3,left,loop) = kt(2)
c store line 7 , skip 8
  call gtside(jm,jt,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(4)
  node(3,left,loop) = node(num,jt,loop)
c store line 9, skip 10
  call gtside(jo,jm,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(8)
  node(3,left,loop) = kt(4)
c store line 11, skip 12
  call gtside(jn,jt,left,nrit)
  node(1,left,loop) = 2
  node(2,left,loop) = kt(2)
  node(3,left,loop) = node(num,jt,loop)
c
  return
  end

```

CNAME TRAN2.FOR

```

-----
subroutine tran2(loop,iter,ibeg)
-----

```

C PROGRAM: TRAN2.FOR

C PURPOSE: DO TRANSITION BY #2 METHOD

C WRITTEN BY: SSUTA HSU CHECKED BY: DATE: OCT. 6, 1988

C MODIFIED BY: CHECKED BY: DATE:

```

-----
C ALGORITHM EXPLANATION: do the transition on two opposit line of
c                    a surface in a volume
c
c
-----

```

c called:

```

c    cktran    -- check the nodes on the line to see if the node have
c                    been used or not
c    tran21    -- do the transition on line 1-3,1-2,...etc.
c    tran22    -- do the transition on line 2-4,3-4,...etc.
c    push      -- store the data into new loop
c
-----

```

c called by: transi.for

```

-----
include 'para.inc'
dimension coor(4,maxnod),node(maxsid,12,maxloop),isur(4,12)
dimension kt(4)
common/auto/coor,node,last,unit
data isur/9,7,11,5, 9,3,10,1, 5,2, 8,1, 5,11,7,9,
$            1,8, 2,5, 1,10,3,9, 10,8,12,8, 11,4,12,2,
$            7,4, 8,3, 8,12,8,10, 3,8, 4,7, 2,12,4,11/
c

```

 iend = ibeg+8

```

num1 = node(1, isur(1, ibeg), loop)
num2 = node(1, isur(3, ibeg), loop)
if(num1.ge.4 .and. num2.eq.2) then
  j1 = 1
  j2 = 3
else if(num2.ge.4 .and. num1.eq.2) then
  j1 = 3
  j2 = 1
end if
kt(1) = last+1
do 20 i = 1,4
  $
  $
  $ coor(i,kt(1)) =((coor(i,node(2, isur(j2, ibeg), loop)) +
  $ coor(i,node(3, isur(j2, ibeg), loop)))/2.0 +
  $ coor(i,node(3, isur(j1, ibeg), loop)))/2.0
20 continue
call cktran(coor(1,kt(1)),kt(1))
kt(2) = last+1
numj1 = node(1, isur(j1, ibeg), loop)
do 30 i = 1,4
  $
  $
  $ coor(i,kt(2)) =((coor(i,node(2, isur(j2, ibeg), loop)) +
  $ coor(i,node(3, isur(j2, ibeg), loop)))/2.0 +
  $ coor(i,node(numj1, isur(j1, ibeg), loop)))/2.0
30 continue
call cktran(coor(1,kt(2)),kt(2))
kt(3) = last+1
do 40 i = 1,4
  $
  $
  $ coor(i,kt(3)) =((coor(i,node(2, isur(j2, iend), loop)) +
  $ coor(i,node(3, isur(j2, iend), loop)))/2.0 +
  $ coor(i,node(3, isur(j1, iend), loop)))/2.0
40 continue
call cktran(coor(1,kt(3)),kt(3))
kt(4) = last+1
numj2 = node(1, isur(j1, iend), loop)
do 50 i = 1,4
  $
  $
  $ coor(i,kt(4)) =((coor(i,node(2, isur(j2, iend), loop)) +
  $ coor(i,node(3, isur(j2, iend), loop)))/2.0 +
  $ coor(i,node(numj2, isur(j1, iend), loop)))/2.0
50 continue
call cktran(coor(1,kt(4)),kt(4))
if(num1.ge.4 .and. num2.eq.2) then
  call tran21(loop, ibeg, iend, kt)
else if(num2.ge.4 .and. num1.eq.2) then
  call tran22(loop, ibeg, iend, kt)
end if
call push(loop, iter)
do 80 i = 1,3
  loop = loop+1
  call push(loop, iter)
80 continue
c
return
end

```

CNAME TRAN21.FOR

```

c-----
c      subroutine tran21(loop, ibeg, iend, kt)
c-----
C
C PROGRAM: TRAN21.FOR
C
C PURPOSE: do the bottom transition for # 2
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C-----
C ALGORITHM EXPLANATION:
c
c-----
c called:  gtside -- given the two opposite lines on a surface then
c           get the other two lines.
c-----
c called by: tran2.for
c-----
include 'para.inc'

```

```

dimension coor(4,maxnod),node(maxsid,12,maxloop),isur(4,12)
dimension kt(4)
logical oline
common/auto/coor,node,last,unit
common/oldln/nnum(20,maxsid=6),lold(maxsid=6),nall
data isur/9,7,11,5, 9,3,10,1, 5,2, 8,1, 5,11,7,9,
$          1,6, 2,5, 1,10,3,9, 10,8,12,8, 11,4,12,2,
$          7,4, 8,3, 8,12,8,10, 3,8, 4,7, 2,12,4,11/
c
c...store them into 4 new loop
c
      loop4 = loop + 1
      loop3 = loop + 2
      loop2 = loop + 3
      call gtside(isur(1,ibeg),isur(1,iend),n9,n10)
      call gtside(isur(3,ibeg),isur(3,iend),n11,n12)
c
c for transition on top line 2,4
c
c store the loop 4, at the top
      nn = 0
      nm = -1
      do 10 i = ibeg,iend,(iend-ibeg)
        do 20 j = 1,4
          node(1,isur(j,i),loop4) = 2
20      continue
c store line 2,4,1,3
        do 30 j = 1,2
          node(j-1,isur(3,i),loop4) = node(j+1,isur(3,i),loop)
          nn = nn + 1
          node(j+1,isur(1,i),loop4) = kt(nn)
30      continue
c store line 5 & 7
          nm = nm + 2
          node(2,isur(4,i),loop4) = kt(nm)
          node(3,isur(4,i),loop4) = node(2,isur(3,i),loop)
c store line 6 & 8
          node(2,isur(2,i),loop4) = kt(1-nm)
          nc = node(1,isur(2,i),loop)+1
          node(3,isur(2,i),loop4) = node(nc,isur(2,i),loop)
10      continue
c
c store line 11,12
c
      do 40 i = n11,n12,(n12-n11)
        do 50 j = 1,node(1,i,loop)+1
          node(j,i,loop4) = node(j,i,loop)
50      continue
40      continue
c
c store line 9,10
c
      nm = 0
      line1 = isur(1,ibeg)
      line2 = isur(1,iend)
      left = n11
      do 60 j = n9,n10,(n10-n9)
        k = last
        nm = nm-1
        oline = .false.
        call fnwnd(line1,line2,nm,nm,left,loop4,nnod,oline,nold)
        if(oline)then
          nnod = lold(nold)
          node(1,j,loop4) = nnod
          do 70 i=1,nnod
            node(i+1,j,loop4) = nnum(i,nold)
70      continue
          last = k
        else
          if(j.eq.n9 .and. nnod.ge.node(1,n11,loop))then
            nnod = node(1,n11,loop)
          else if(j.eq.n10 .and. nnod.ge.node(1,n12,loop))then
            nnod = node(1,n12,loop)
          end if
          node(1,j,loop4) = nnod
          node(2,j,loop4) = node(nm+1,line1,loop4)
          do 80 i = 2,nnod-1
            node(i+1,j,loop4) = k+i-1
80      continue
          node(nnod+1,j,loop4) = node(nm+1,line2,loop4)

```

```

        last = k + nnod - 2
        end if
        left = n12
80      continue
c
c store the loop 2, at the left side
c
      do 110 i = ibeg, iend, (iend-ibeg)
c store line 1,3
        node(1, isur(1,i), loop2) = 2
        node(2, isur(1,i), loop2) = node(2, isur(1,i), loop)
        node(3, isur(1,i), loop2) = node(3, isur(1,i), loop)
c store line 2,4
        node(1, isur(3,i), loop2) = 2
        node(2, isur(3,i), loop2) = node(3, isur(4,i), loop4)
        node(3, isur(3,i), loop2) = node(2, isur(4,i), loop4)
c store line 5,7
        node(1, isur(4,i), loop2) = node(1, isur(4,i), loop)
        do 100 j = 1, node(1, isur(4,i), loop)
            node(j+1, isur(4,i), loop2) = node(j-1, isur(4,i), loop)
100      continue
c store line 6 & 8
        node(1, isur(2,i), loop2) = 2
        node(2, isur(2,i), loop2) = node(3, isur(1,i), loop)
        node(3, isur(2,i), loop2) = node(2, isur(1,i), loop4)
110     continue
c
c store line 9 & 11
c
      do 120 i = n9, n11, (n11-n9)
        do 130 j = 1, node(1, i, loop)+1
            node(j, i, loop2) = node(j, i, loop)
130      continue
120     continue
c
c store line 12
c
      do 131 i = 1, node(1, n9, loop4)+1
        node(i, n12, loop2) = node(i, n9, loop4)
131     continue
c
c store line 10
c
      k = last
      line1 = isur(1, ibeg)
      line2 = isur(1, iend)
      oline = .false.
      call fnewnd(line1, line2, 2, 2, left, loop, nnod, oline, nold)
c
c store split line into right line of loop two
c
      if(oline) then
        nnod = lold(nold)
        node(1, n10, loop2) = nnod
        do 132 i=1, nnod
            node(i+1, n10, loop2) = nnum(i, nold)
132      continue
          last = k
c
c if the old line was found then delete that line in nnum to save space
c
          do 133 i = nold, nall-1
            lold(i) = lold(i+1)
            do 134 jj = 1, lold(i)
                nnum(jj, i) = nnum(jj, i+1)
134      continue
133      continue
          nall = nall - 1
        else
          if(nnod.ge.node(1, n9, loop)) nnod = node(1, n9, loop)
          node(1, n10, loop2) = nnod
          node(2, n10, loop2) = node(3, line1, loop)
          do 135 i = 2, nnod-1
            node(i+1, n10, loop2) = k+i-1
135      continue
          node(nnod+1, n10, loop2) = node(3, line2, loop)
          last = k + nnod - 2
c
c store this new line data for old line reference check
c

```

```

        nall = nall + 1
        do 136 i = 1, nnod
            nnum(i,nall) = node(i+1,n10,loop2)
136      continue
        lold(nall) = nnod
    end if

c
c store the loop 3, the bottom side
c
    do 170 i = ibeg,iend,(iend-ibeg)
c store line 2,4
        do 140 j = 1,3
            node(j,isur(3,i),loop3) = node(j,isur(1,i),loop4)
140      continue
c store line 1,3
        nct = node(1,isur(1,i),loop)-2
        node(1,isur(1,i),loop3) = nct
        do 150 j = 1,nct
            node(j+1,isur(1,i),loop3) = node(j+2,isur(1,i),loop)
150      continue
c store line 5,7
        node(1,isur(4,i),loop3) = 2
        node(2,isur(4,i),loop3) = node(3,isur(1,i),loop)
        node(3,isur(4,i),loop3) = node(2,isur(1,i),loop4)
c store line 8,8
        node(1,isur(2,i),loop3) = 2
        node(2,isur(2,i),loop3) = node(nct+2,isur(1,i),loop)
        node(3,isur(2,i),loop3) = node(3,isur(1,i),loop4)
170      continue
c
c store line 9,11,12
c
        do 180 j = 1,node(1,n10,loop2)+1
            node(j,n9,loop3) = node(j,n10,loop2)
180      continue
        do 190 j = 1,node(1,n9,loop4)+1
            node(j,n11,loop3) = node(j,n9,loop4)
190      continue
        do 191 i = 1,node(1,n10,loop4)+1
            node(i,n12,loop3) = node(i,n10,loop4)
191      continue
c
c store line 10
c
        k = last
        num1 = node(1,line1,loop)-1
        num2 = node(1,line2,loop)-1
        oline = .false.
        left = n10
        call fnewnd(line1,line2,num1,num2,left,loop,nnod,oline,nold)
        k = last
        if(oline)then
            nnod = lold(nold)
            node(1,n10,loop3) = nnod
            do 192 i=1,nnod
                node(i+1,n10,loop3) = nnum(i,nold)
192          continue
            last = k
c
c if the old line was found then delete that line to save space
c
        do 195 i = nold,nall-1
            lold(i) = lold(i+1)
            do 196 jj = 1,lold(i)
                nnum(jj,i) = nnum(jj,i+1)
196          continue
195          continue
            nall = nall - 1
        else
            if(nnod.ge.node(1,n10,loop)) nnod = node(1,n10,loop)
            node(1,n10,loop3) = nnod
            node(2,n10,loop3) = node(num1+1,line1,loop)
            do 193 i = 2,nnod-1
                node(i+1,n10,loop3) = k+i-1
193          continue
            node(nnod+1,n10,loop3) = node(num2+1,line2,loop)
            last = k + nnod - 2
c
c store this new line data for old line reference check
c

```

```

        nall = nall + 1
        do 194 i = 1, nnod
            nnum(i,nall) = node(i+1,n10,loop3)
194      continue
        lold(nall) = nnod
    end if

c
c store the loop 1, the right side
c
        do 230 i = ibeg,iend,(iend-ibeg)
c store line 2,4
            do 200 j = 1,3
                node(j,isur(3,i),loop) = node(j,isur(2,i),loop4)
200      continue
c store line 1,3
            node(1,isur(1,i),loop) = 2
            node(2,isur(1,i),loop) = node(2+nct,isur(1,i),loop)
            node(3,isur(1,i),loop) = node(3+nct,isur(1,i),loop)
c store line 5,7
            do 220 j = 1,3
                node(j,isur(4,i),loop) = node(j,isur(2,i),loop3)
220      continue
c
c skips line 6,8
c
230      continue
c
c store line 9,11
c
        do 240 i = 1,node(1,n10,loop3)+1
            node(i,n9,loop) = node(i,n10,loop3)
240      continue
        do 241 i = 1,node(1,n12,loop3)+1
            node(i,n11,loop) = node(i,n12,loop3)
241      continue
c
c skips line n10,n12
c
        return
    end

```

CNAME TRAN22.FOR

```

c-----
c          subroutine tran22(loop,ibeg,iend,kt)
c-----
C
C PROGRAM: TRAN22.FOR
C
C PURPOSE: do the top transition for # 2
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C
c-----
C ALGORITHM EXPLANATION: do transition for line 2,4, or 6,8, or 10,12
c
c
c called:  gtside -- given the two opposite lines on a surface then
c           get the other two lines.
c
c-----
c called by: tran2.for
c
c-----
        include 'para.inc'
        dimension coor(4,maxnod),node(maxsid,12,maxloop),isur(4,12)
        dimension kt(4)
        logical oline
        common/auto/coor,node,last,unit
        common/oldln/nnum(20,maxsid-6),lold(maxsid-6),nall
        data isur/9,7,11,5, 9,3,10,1, 5,2, 8,1, 5,11,7,9,
$           1,6, 2,5, 1,10,3,9, 10,8,12,6, 11,4,12,2,
$           7,4, 8,3, 8,12,8,10, 3,8, 4,7, 2,12,4,11/
c
c...store them into 4 new loop
c
        loop3 = loop + 1

```

```

loop4 = loop + 2
loop2 = loop + 3
call gtside(isur(1,ibeg),isur(1,iend),n9,n10)
call gtside(isur(3,ibeg),isur(3,iend),n11,n12)
c
c store the loop 4, at the bottom side
c
  nn = 0
  nm = -1
  do 10 i = ibeg,iend,(iend-ibeg)
    do 20 j = 1,4
      node(1,isur(j,i),loop4) = 2
20    continue
c store line 1 & 3
    do 30 j = 1,2
      node(j+1,isur(1,i),loop4) = node(j+1,isur(1,i),loop)
c store line 2 & 4
      nn = nn + 1
      node(j+1,isur(3,i),loop4) = kt(nn)
30    continue
c store line 5 & 7
      nm = nm + 2
      node(2,isur(4,i),loop4) = node(2,isur(1,i),loop)
      node(3,isur(4,i),loop4) = kt(nm)
c store line 6 & 8
      node(2,isur(2,i),loop4) = node(2,isur(2,i),loop)
      node(3,isur(2,i),loop4) = kt(1+nm)
10    continue
c store line 9 & 10
c
c find the lines on both sides
c
  do 40 i = n9,n10,(n10-n9)
    nod = node(1,i,loop)
    do 50 j = 1,nod+1
      node(j,i,loop4) = node(j,i,loop)
50    continue
40    continue
c
c store line 11,12
c
  nm = 0
  line1 = isur(3,ibeg)
  line2 = isur(3,iend)
  left = n9
  do 60 j = n11,n12,(n12-n11)
    k = last
    nm = nm+1
    oline = .false.
    call fnewnd(line1,line2,nm,nm,left,loop4,nnod,oline,nold)
    if(oline)then
      nnod = lold(nold)
      node(1,j,loop4) = nnod
      do 70 i=1,nnod
        node(i+1,j,loop4) = nnum(i,nold)
70      continue
      last = k
    else
      if(j.eq.n11 .and. nnod.ge.node(1,n9,loop))then
        nnod = node(1,n9,loop)
      else if(j.eq.n12 .and. nnod.ge.node(1,n10,loop))then
        nnod = node(1,n10,loop)
      end if
      node(1,j,loop4) = nnod
      node(2,j,loop4) = node(nm+1,line1,loop4)
      do 73 i = 2,nnod-1
        node(i+1,j,loop4) = k+i-1
73      continue
      node(nnod+1,j,loop4) = node(nm+1,line2,loop4)
      last = k + nnod - 2
    end if
    left = n10
60    continue
c
c store the loop 2, at the left side
c
  do 110 i = ibeg,iend,(iend-ibeg)
c store line 1,3
    do 80 j = 1,3
      node(j,isur(1,i),loop2) = node(j,isur(4,i),loop4)

```



```

80      continue
c store line 2 & 4
      node(1,isur(3,i),loop2) = 2
      do 90 j = 1,2
        node(j+1,isur(3,i),loop2) = node(j+1,isur(3,i),loop)
90      continue
c store line 5 & 7
      node(1,isur(4,i),loop2) = node(1,isur(4,i),loop)
      do 100 j = 1,node(1,isur(4,i),loop)
        node(j+1,isur(4,i),loop2) = node(j+1,isur(4,i),loop)
100     continue
c store line 6 & 8
      node(1,isur(2,i),loop2) = 2
      node(2,isur(2,i),loop2) = node(2,isur(3,i),loop4)
      node(3,isur(2,i),loop2) = node(3,isur(3,i),loop)
110    continue
c store line 9 & 11
c
      do 120 i = n9,n11,(n11-n9)
        nod = node(1,i,loop)
        do 130 j = 1,nod+1
          node(j,i,loop2) = node(j,i,loop)
130     continue
120    continue
c
c store line 10
c
      do 140 i = 1,node(1,n11,loop4)+1
        node(i,n10,loop2) = node(i,n11,loop4)
140    continue
c
c store line 12
c
      k = last
      line1 = isur(3,ibeg)
      line2 = isur(3,iend)
      left = n11
      oline = .false.
      call fnewnd(line1,line2,2,2,left,loop,nnod,oline,nold)
      if(oline)then
        nnod = lold(nold)
        node(1,n12,loop2) = nnod
        do 150 i=1,nnod
          node(i+1,n12,loop2) = nnum(i,nold)
150     continue
        last = k
c
c if the old line was found then delete that line to save space
c
      do 151 i = nold,nall-1
        lold(i) = lold(i+1)
        do 152 jj = 1,lold(i)
          nnum(jj,i) = nnum(jj,i+1)
152     continue
151     continue
        nall = nall - 1
      else
        if(nnod.ge.node(1,n11,loop)) nnod = node(1,n11,loop)
        node(1,n12,loop2) = nnod
        node(2,n12,loop2) = node(3,line1,loop)
        do 160 i = 2,nnod-1
          node(i+1,n12,loop2) = k+i-1
160     continue
        node(nnod+1,n12,loop2) = node(3,line2,loop)
        last = k + nnod - 2
c
c store this new line data for old line reference check
c
      nall = nall + 1
      do 161 i = 1, nnod
        nnum(i,nall) = node(i+1,n12,loop2)
161     continue
        lold(nall) = nnod
      end if
c
c store the loop 3, the top side
c
      do 170 i = ibeg,iend,(iend-ibeg)
c store line 1,3
        node(1,isur(1,i),loop3) = 2

```

```

        node(2,isur(1,i),loop3) = node(2,isur(3,i),loop4)
        node(3,isur(1,i),loop3) = node(3,isur(3,i),loop4)
c store line 2,4
        nct = node(1,isur(3,i),loop)-2
        node(1,isur(3,i),loop3) = nct
        do 180 j = 1,nct
            node(j+1,isur(3,i),loop3) = node(j+2,isur(3,i),loop)
180        continue
c store line 5,7
        node(1,isur(4,i),loop3) = 2
        node(2,isur(4,i),loop3) = node(2,isur(3,i),loop4)
        node(3,isur(4,i),loop3) = node(3,isur(3,i),loop)
c store line 6,8
        node(1,isur(2,i),loop3) = 2
        node(2,isur(2,i),loop3) = node(3,isur(3,i),loop4)
        node(3,isur(2,i),loop3) = node(nct+2,isur(3,i),loop)
170        continue
c
c store line 9,10,11
c
        do 190 j = 1,node(1,n11,loop4)+1
            node(j,n9,loop3) = node(j,n11,loop4)
190        continue
        do 191 j = 1,node(1,n12,loop4)+1
            node(j,n10,loop3) = node(j,n12,loop4)
191        continue
        do 192 i = 1,node(1,n12,loop2)+1
            node(i,n11,loop3) = node(i,n12,loop2)
192        continue
c
c store line 12
c
        k = last
        num1 = node(1,line1,loop)-1
        num2 = node(1,line2,loop)-1
        left = n12
        oline = .false.
        call fnewnd(line1,line2,num1,num2,left loop,nnod,oline,nold)
        k = last
        if(oline)then
            nnod = lold(nold)
            node(1,n12,loop3) = nnod
            do 194 i=1,nnod
                node(i+1,n12,loop3) = nnum(i,nold)
194            continue
            last = k
c
c if the old line was found then delete that line to save space
c
            do 195 i = nold,nall-1
                lold(i) = lold(i+1)
                do 196 jj = 1,lold(i)
                    nnum(jj,i) = nnum(jj,i+1)
196                continue
195            continue
            nall = nall - 1
        else
            if(nnod.ge.node(1,n12,loop)) nnod = node(1,n12,loop)
            node(1,n12,loop3) = nnod
            node(2,n12,loop3) = node(num1+1,line1,loop)
            do 197 i = 2,nnod-1
                node(i+1,n12,loop3) = k+i-1
197            continue
            node(nnod+1,n12,loop3) = node(num2+1,line2,loop)
            last = k + nnod - 2
c
c store this new line data for old line reference check
c
            nall = nall + 1
            do 198 i = 1, nnod
                nnum(i,nall) = node(i+1,n12,loop3)
198            continue
            lold(nall) = nnod
        end if
c
c store the loop 1, the right side
c
        do 230 i = ibeg,iend,(iend-ibeg)
c store line 1,3
            node(1,isur(1,i),loop) = 2

```

```

        node(2,isur(1,i),loop) = node(3,isur(2,i),loop4)
        node(3,isur(1,i),loop) = node(2,isur(2,i),loop4)
c store line 2,4
        node(1,isur(3,i),loop) = 2
        node(2,isur(3,i),loop) = node(nct+2,isur(3,i),loop)
        node(3,isur(3,i),loop) = node(nct+3,isur(3,i),loop)
c store line 5,7
        do 220 j = 1,3
            node(j,isur(4,i),loop) = node(j,isur(2,i),loop3)
220      continue
c
c skips line 6,8
c
230      continue
c
c store line 9,11
c
        do 240 i = 1,node(1,n10,loop3)+1
            node(i,n9,loop) = node(i,n10,loop3)
240      continue
        do 241 i = 1,node(1,n12,loop3)+1
            node(i,n11,loop) = node(i,n12,loop3)
241      continue
c
c skips line 10,12
c
        return
        end

```

CNAME TRAN5.FOR

```

c-----
      subroutine tran5(loop,iter)
c-----
C
C PROGRAM: TRAN5.FOR
C
C PURPOSE: DO TRANSITION BY #5 METHOD
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C MODIFIED BY:           CHECKED BY:           DATE:
C
c-----
C
C ALGORITHM EXPLANATION: If model was too complex to subdivided by
c either tran1 or tran2 method, then add a cube inside the old solid
c
c-----
c
c called:
c   tran51 -- add a new hexahedron in the center of old hexahedron
c   push   -- store the data into new loop
c
c-----
c
c called by: transi.for -- transition flag checker
c
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),cent(4)
      common/auto/coor,node,last,unit
c
      k = last
c
c find the 8 point of center cube
c
      num = node(1,4,loop)
      do 10 i = 1,4
          diff = coor(i,node(num+1,4,loop)) - coor(i,node(2,1,loop))
          coor(i,k+1) = coor(i,node(2,1,loop)) + diff/3.0
          coor(i,k+8) = coor(i,node(2,1,loop)) + diff*2.0/3.0
10      continue
      num = node(1,1,loop)
      do 20 i = 1,4
          diff = coor(i,node(2,4,loop)) - coor(i,node(num+1,1,loop))
          coor(i,k+2) = coor(i,node(num+1,1,loop)) + diff/3.0
          coor(i,k+7) = coor(i,node(num+1,1,loop)) + diff*2.0/3.0
20      continue
      num = node(1,3,loop)
      do 30 i = 1,4
          diff = coor(i,node(num+1,3,loop)) - coor(i,node(2,2,loop))

```

```

        coor(i,k+3) = coor(i,node(2,2,loop)) + diff/3.0
        coor(i,k+6) = coor(i,node(2,2,loop)) + diff*2.0/3.0
30      continue
        num = node(1,2,loop)
        do 40 i = 1,4
            diff = coor(i,node(2,3,loop)) - coor(i,node(num+1,2,loop))
            coor(i,k+4) = coor(i,node(num+1,2,loop)) + diff/3.0
            coor(i,k+5) = coor(i,node(num+1,2,loop)) + diff*2.0/3.0
40      continue
        call tran51(loop)
        last = k+8
        call push(loop,iter)
        do 50 i = 1,6
            loop = loop+1
            call push(loop,iter)
50      continue
c
c      return
c      end

```

CNAME TRAN51.FOR

```

C-----
c      subroutine tran51(loop)
C-----
C
C      PROGRAM: TRAN51.FOR
C
C      PURPOSE: DO TRANSITION BY #51 METHOD
C
C      WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C      MODIFIED BY:                CHECKED BY:           DATE:
C-----
C
C      ALGORITHM EXPLANATION: add a cube to the center of the complex
c      transition solid
c
c-----
c
c      called by: tran5.for -- special number of nodes controller
c
c-----
        include 'para.inc'
        dimension coor(4,maxnod),node(maxsid,12,maxloop),isur(4,6)
        common/auto/coor,last,unit
        data isur/1,6,2,5, 1,10,3,9, 10,8,12,6, 2,12,4,11,
          $      5,11,7,9, 3,8,4,7/
c...store them into 5 new loop
        k = last
        loop2 = loop + 1
        loop3 = loop + 2
        loop4 = loop + 3
        loop5 = loop + 4
        loop6 = loop + 5
        loop7 = loop + 6
c initialize the node number
        do 20 j = 1,6
            do 10 i = 1,12
                node(1,i,loop+j) = 2
10          continue
20      continue
c
c      store one surface for each loop first
c
        do 50 m = 1,6
            do 40 j = 1,4
                do 30 i = 1,node(1,isur(j,m),loop)+1
                    node(i,isur(j,m),loop+m) = node(i,isur(j,m),loop)
30          continue
40          continue
50      continue
c store the loop 2, the front side
c store line 3,4,7,8
        inc = -1
        do 60 i = 3,4
            inc = inc + 2
            node(2,i,loop2) = k+inc
            node(3,i,loop2) = k+inc+1
60      continue
        inc = 0

```

```

do 70 i = 7,8
  inc = inc + 1
  node(2,i,loop2) = k+inc
  node(3,i,loop2) = k+inc+2
70  continue
c store line 9,10,11,12
  inc = 0
  do 80 i = 9,12
    inc = inc + 1
    node(2,i,loop2) = node(2,i,loop)
    node(3,i,loop2) = k+inc
80  continue
c store the loop 3, the bottom side
c store line 2,4,11,12
  inc = -3
  do 90 i = 2,4,2
    inc = inc + 4
    node(2,i,loop3) = k+inc
    node(3,i,loop3) = k+inc+1
90  continue
  inc = 0
  do 100 i = 11,12
    inc = inc + 1
    node(2,i,loop3) = k+inc
    node(3,i,loop3) = k+inc+4
100 continue
c store line 5,6,7,8
  inc = 0
  do 110 i = 5,6
    inc = inc + 1
    node(2,i,loop3) = node(2,i,loop)
    node(3,i,loop3) = k+inc
110 continue
  inc = 4
  do 120 i = 7,8
    inc = inc + 1
    node(2,i,loop3) = node(2,i,loop)
    node(3,i,loop3) = k+inc
120 continue
c
c store the loop 4, the right side
c store line 9,11,5,7
  inc = 0
  do 130 i = 9,11,2
    inc = inc + 2
    node(2,i,loop4) = k+inc
    node(3,i,loop4) = k+inc+4
130 continue
  inc = -2
  do 140 i = 5,7,2
    inc = inc + 4
    node(2,i,loop4) = k+inc
    node(3,i,loop4) = k+inc+2
140 continue
c store line 1,2,3,4
  inc = 0
  do 150 i = 1,4
    num = node(1,i,loop)
    inc = inc + 2
    node(2,i,loop4) = k+inc
    node(3,i,loop4) = node(num+1,i,loop)
150 continue
c
c store the loop 5, the top side
c store line 1,3,9,10
  inc = -1
  do 160 i = 1,3,2
    inc = inc + 4
    node(2,i,loop5) = k+inc
    node(3,i,loop5) = k+inc+1
160 continue
  inc = 2
  do 170 i = 9,10
    inc = inc + 1
    node(2,i,loop5) = k+inc
    node(3,i,loop5) = k+inc+4
170 continue
c store line 5,6,7,8
  inc = 2
  do 180 i = 5,6

```

```

        num = node(1,i,loop)
        inc = inc + 1
        node(2,i,loop5) = k+inc
        node(3,i,loop5) = node(num+1,i,loop)
180    continue
        inc = 8
        do 190 i = 7,8
            num = node(1,i,loop)
            inc = inc + 1
            node(2,i,loop5) = k+inc
            node(3,i,loop5) = node(num+1,i,loop)
190    continue
c
c store the loop 6, the left side
c store line 10,12,6,8
        inc = -1
        do 200 i = 10,12,2
            inc = inc + 2
            node(2,i,loop6) = k+inc
            node(3,i,loop6) = k-inc+4
200    continue
        inc = -3
        do 210 i = 6,8,2
            inc = inc + 4
            node(2,i,loop6) = k+inc
            node(3,i,loop6) = k+inc+2
210    continue
c store line 1,2,3,4
        inc = -1
        do 220 i = 1,4
            inc = inc + 2
            node(2,i,loop6) = node(2,i,loop)
            node(3,i,loop6) = k+inc
220    continue
c
c store the loop 7, the back side
c store line 1,2,5,8
        inc = 3
        do 240 i = 1,2
            inc = inc + 2
            node(2,i,loop7) = k+inc
            node(3,i,loop7) = k+inc+1
240    continue
        inc = 4
        do 250 i = 5,8
            inc = inc + 1
            node(2,i,loop7) = k+inc
            node(3,i,loop7) = k+inc+2
250    continue
c store line 9,10,11,12
        inc = 4
        do 260 i = 9,12
            num = node(1,i,loop)
            inc = inc + 1
            node(2,i,loop7) = k+inc
            node(3,i,loop7) = node(num+1,i,loop)
260    continue
c
c store loop1, the center cube
c store line 1,2,3,4
        inc = -1
        do 270 i = 1,4
            node(1,i,loop) = 2
            inc = inc + 2
            node(2,i,loop) = k+inc
            node(3,i,loop) = k+inc+1
270    continue
c store line 5,6,7,8
        inc = 0
        do 280 i = 5,8
            node(1,i,loop) = 2
            inc = inc + 1
            node(2,i,loop) = k+inc
            node(3,i,loop) = k+inc+2
280    continue
        inc = 4
        do 290 i = 7,8
            node(1,i,loop) = 2
            inc = inc + 1
            node(2,i,loop) = k+inc

```

```

      node(3,i,loop) = k+inc+2
290  continue
c store line 9,10,11,12
      inc = 0
      do 300 i = 9,12
          node(1,i,loop) = 2
          inc = inc + 1
          node(2,i,loop) = k+inc
          node(3,i,loop) = k+inc+4
300  continue
c
      return
      end

```

CNAME TRANSI.FOR

```

c-----
      subroutine transi(loop,iter,iflagx,iflagy,iflagz)
c-----
C
C PROGRAM: TRANSI.FOR
C
C PURPOSE: CHOOSE THE PROPER TRANSITION
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C ALGORITHM EXPLANATION: check the transition flag and assign the proper
C transition
c-----
c
c called:
c   tran1.for -- transition method 1
c   tran2.for -- transition method 2
c   tran5.for -- transition method 5
c-----
c
c called by:
c   vmesh.for -- volume meshing controler
c-----
c
c do #5 transition --- add cube in the center
c
c   if(iflagx.eq.5 .or. iflagy.eq.5 .or. iflagz.eq.5)then
c       call tran5(loop,iter)
c
c do z-direction transition
c
c   else if(iflagz.eq.21 .and. iflagx.ne.21)then
c       call tran2(loop,iter,2)
c   else if(iflagz.eq.22 .and. iflagy.ne.21)then
c       call tran2(loop,iter,1)
c   else if(iflagz.eq.11 .and. iflagx.ne.21 .and. iflagy.ne.21)then
c       call tran1(loop,iter,9,iflagz)
c   else if(iflagz.eq.12 .and. iflagx.ne.21 .and. iflagy.ne.21)then
c       call tran1(loop,iter,10,iflagz)
c   else if(iflagz.eq.13 .and. iflagx.ne.21 .and. iflagy.ne.21)then
c       call tran1(loop,iter,11,iflagz)
c   else if(iflagz.eq.14 .and. iflagx.ne.21 .and. iflagy.ne.21)then
c       call tran1(loop,iter,12,iflagz)
c
c do y-direction transition
c
c   else if(iflagy.eq.21 .and. iflagx.ne.22)then
c       call tran2(loop,iter,3)
c   else if(iflagy.eq.22 .and. iflagz.ne.21)then
c       call tran2(loop,iter,4)
c   else if(iflagy.eq.11 .and. iflagx.ne.22 .and. iflagz.ne.21)then
c       call tran1(loop,iter,5,iflagy)
c   else if(iflagy.eq.12 .and. iflagx.ne.22 .and. iflagz.ne.21)then
c       call tran1(loop,iter,6,iflagy)
c   else if(iflagy.eq.13 .and. iflagx.ne.22 .and. iflagz.ne.21)then
c       call tran1(loop,iter,7,iflagy)
c   else if(iflagy.eq.14 .and. iflagx.ne.22 .and. iflagz.ne.21)then
c       call tran1(loop,iter,8,iflagy)
c

```

```

c do x-direction transition
c
  else if(iflagx.eq.21 .and. iflagy.ne.22)then
    call tran2(loop,iter,5)
  else if(iflagx.eq.22 .and. iflagz.ne.22)then
    call tran2(loop,iter,6)
  else if(iflagx.eq.11 .and. iflagy.ne.22 .and. iflagz.ne.22)then
    call tran1(loop,iter,1,iflagx)
  else if(iflagx.eq.12 .and. iflagy.ne.22 .and. iflagz.ne.22)then
    call tran1(loop,iter,2,iflagx)
  else if(iflagx.eq.13 .and. iflagy.ne.22 .and. iflagz.ne.22)then
    call tran1(loop,iter,3,iflagx)
  else if(iflagx.eq.14 .and. iflagy.ne.22 .and. iflagz.ne.22)then
    call tran1(loop,iter,4,iflagx)
  end if
c
  return
  end

CNAME VGTNOD.FOR
c-----
c      subroutine vgtnod(loop,next,linel,line2,num1,num2)
c-----
C
C PROGRAM: VGTNOD.FOR
C
C PURPOSE: FIND THE BEST SPLIT NODES BETWEEN TWO LINES
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C-----
C
C ALGORITHM EXPLANATION: CHECK THROUGH THE OLD LINE TO FIND THE BEST SPLIT
C NODES, IF NOT FOUND THEN CREATE A BEST LINE
c-----
c
c called:
c   vnxt1.for -- find the next best node
c   gtside.for -- get other pair of lines in a surface
c   vnewds.for -- get new distance
c-----
c
c called by:
c   xmesh.for -- x-direction mesher
c   ymesh.for -- y-direction mesher
c   zmesh.for -- z-direction mesher
c-----
c
c   include 'para.inc'
c   dimension coor(4,maxnod),node(maxsid,12,maxloop),xo(maxsid),
c   $           yo(maxsid),zo(maxsid),iorder(2,12)
c   common/auto/coor,node,last,unit
c   logical found
c   data iorder/2,3, 4,1, 1,4, 3,2,  6,7, 8,5, 5,8, 7,8,
c   $           10,11, 12,9, 9,12, 11,10/
c
c check for the existing line, if not found then make a new one
c
c   line2 = iorder(next,linel)
c   call vnxt1(loop,linel,line2,num1,num2,found)
c   if(found) return
c
c find the best split line
c
c   nnod = node(1,line2,loop)
c   nod1 = node(num1+1,linel,loop)
c   dmin = 1.e10
c   dist = 1.e9
c   k = 0
c   k = 1
c   do while (dist.lt.dmin .and. k.lt.nnod)
c     dmin = dist
c     k = k+2
c     k = k+1
c     if(k.lt.nnod)then
c       nod2 = node(k+1,line2,loop)
c       dist = sqrt((coor(1,nod1)-coor(1,nod2))**2+

```



```

      $          (coor(2,nod1)-coor(2,nod2))**2+
      $          (coor(3,nod1)-coor(3,nod2))**2)
c
c   get the deformed distance -- better result
c       call vnewds(nod1,nod2,num,xo,yo,zo,dist)
c       end if
c       end do
c       num2 = k-2
C       num2 = k-1
c
c       return
c       end

CNAME VNEWDS.FOR
c-----
      subroutine vnewds(nod1,nod2,num,xo,yo,zo,dist)
c-----
C
C   PROGRAM: VNEWDS.FOR
C
C   PURPOSE: find the distance of a deformed lines
C
C   WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C   MODIFIED BY:          CHECKED BY:          DATE:
C-----
C   ALGORITHM EXPLANATION:
c-----
c
c   called: deform.for -- deformed the nodes
c-----
c
c   called by:
c   vgtnod.for -- find the best node for split line
c   fnewnd.for -- find the new nodes
c-----
      include 'para.inc'
      dimension coor(4,maxnod),node(maxsid,12,maxloop),xo(maxsid),
      $          yo(maxsid),zo(maxsid)
      common/auto/coor,node,last,unit
c
c   assume the number of nodes require same as the left line to start with
c
      sum=0.0
      do 10 j = 1,num
        xi = coor(1,nod1) + (coor(1,nod2)-coor(1,nod1))*(j-1)/(num-1)
        yi = coor(2,nod1) + (coor(2,nod2)-coor(2,nod1))*(j-1)/(num-1)
        zi = coor(3,nod1) + (coor(3,nod2)-coor(3,nod1))*(j-1)/(num-1)
        call deform(xi,yi,zi,xo(j),yo(j),zo(j))
        if(j.gt.1)then
          $          sum = sum + sqrt((xo(j)-xo(j-1))**2 + (yo(j)-yo(j-1))**2 +
          $          (zo(j)-zo(j-1))**2)
        end if
      continue
      dist = sum
c
c       return
c       end

```

```

CNAME VNEWLN.FOR
c-----
      subroutine vnewln(line,num,loop)
c-----
C
C   PROGRAM: VNEWLN.FOR
C
C   PURPOSE: store the old line into two new loop lines at subdivision node
C
C   WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C   MODIFIED BY:          CHECKED BY:          DATE:
C-----
C   ALGORITHM EXPLANATION:
c-----
c

```

```

c
c  called by:
c      xmesh.for -- x-direction mesher
c      ymesh.for -- y-direction mesher
c      zmesh.for -- z-direction mesher
c
c-----
c      include 'para.inc'
c      dimension coord(4,maxnod),node(maxsid,12,maxloop),mtemp(maxsid)
c      common/auto/coord,node,last,unit
c
c      do 10 i = 1,node(1,line,loop)
c          mtemp(i) = node(i+1,line,loop)
10  continue
c
c...store data into new loop one
c
c      node(1,line,loop) = node(1,line,loop)-num+1
c      do 20 i = 1,node(1,line,loop)
c          node(i+1,line,loop) = mtemp(i+num-1)
20  continue
c
c...store data into new loop two
c
c      loop2 = loop + 1
c      node(1,line,loop2) = num
c      do 30 i = 1,num
c          node(i+1,line,loop2) = mtemp(i)
30  continue
c
c      return
c      end

CNAME VNEWND.FOR
c-----
c      subroutine vnewnd(line1,line2,nnod1,nnod2,loop)
c-----
C
C  PROGRAM: VNEWND.FOR
C
C  PURPOSE:  find the new nodes for the split line
C
C  WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C  MODIFIED BY:          CHECKED BY:          DATE:
C
C-----
C  ALGORITHM EXPLANATION: CHECK THE EXISTENCE OF THE OLD LINE IF NOT EXIST
C  THEN ADD NEW NODE TO THAT LINE, OTHERWISE COPY THE OLD LINE NODES
c-----
c
c  called:
c      gtside.for -- get other pair of lines in a surface
c      fnewnd.for -- find the new nodes
c-----
c
c  called by:
c      xmesh.for -- x-direction mesher
c      ymesh.for -- y-direction mesher
c      zmesh.for -- z-direction mesher
c
c-----
c      include 'para.inc'
c      dimension coord(4,maxnod),node(maxsid,12,maxloop)
c      logical oline
c      common/auto/coord,node,last,unit
c      common/oldln/nnum(20,maxsid*8),lold(maxsid*8),nall
c
c      if(line2.lt.line1)then
c          mbeg = line2
c          next = line1
c          num1 = nnod2
c          num2 = nnod1
c      else
c          mbeg = line1
c          next = line2
c          num1 = nnod1
c          num2 = nnod2

```

```

      end if
c
c find the lines on both sides
c
      call gtside(mbeg,next,left,nrit)
c
c find the require nodes in the new line and adjust the next line num
c to fit them into even number.
c
      call fnewnd(mbeg,next,num1,num2,left,loop,nnod,oline,nold)
c
c...store old left line into new left line loop two
c
      k=last
      numrit = node(1,nrit,loop)
      numlft = node(1,left,loop)
      loop2 = loop + 1
      do 30 i = 1,numlft+1
        node(i,left,loop2) = node(i,left,loop)
30      continue
c
c store split line into right line of loop two
c
      if(oline)then
        nnod = lold(nold)
        node(1,nrit,loop2) = nnod
        do 31 i=1,nnod
          node(i+1,nrit,loop2) = nnum(i,nold)
31      continue
          last = k
        else
          node(1,nrit,loop2) = nnod
          node(2,nrit,loop2) = node(num1+1,mbeg,loop)
          do 40 i = 2,nnod-1
            node(i+1,nrit,loop2) = k+i-1
40      continue
            node(nnod+1,nrit,loop2) = node(num2+1,next,loop)
            last = k + nnod - 2
          end if
c
c...store split line into left line of loop one
c* no need to store the old right line(it doesn't change loop number)
c
      do 50 i = 1,nnod+1
        node(i,left,loop) = node(i,nrit,loop2)
50      continue
c
c if the old line was found then delete that line to save space
c
      if(oline)then
        do 80 i = nold,nall-1
          lold(i) = lold(i+1)
          do 90 j = 1,lold(i)
            nnum(j,i) = nnum(j,i+1)
90      continue
80      continue
          nall = nall - 1
c
c if this was a new line then store for old line reference check
c
c      if(line1.ne.1 .and. line1.ne.5 .and. line1.ne.9 .and.
c      $ line2.ne.1 .and. line2.ne.5 .and. line2.ne.9)then
c        else
c          nall = nall + 1
c          do 51 i = 1, nnod
c            nnum(i,nall) = node(i+1,left,loop)
51      continue
c          lold(nall) = nnod
c        end if
c
c...plot the new splitting line **** remove it after done
c
c      x1 = coord(1,node(2,nrit,loop2))
c      y1 = coord(2,node(2,nrit,loop2))
c      z1 = coord(3,node(2,nrit,loop2))
c      do 60 i = 2,nnod
c        x2 = coord(1,node(i+1,nrit,loop2))
c        y2 = coord(2,node(i+1,nrit,loop2))
c        z2 = coord(3,node(i+1,nrit,loop2))
c      call pipe(x1,y1,z1,x2,y2,z2)
c

```



```

        end do
        call choln(nod1,nod2,found,nold)
        if(found)then
            num1 = i - 1
            num2 = n - 1
            return
        end if
10      continue
20      continue
        return
        end
CNAME VSTORE.FOR
c-----
c      subroutine vstore(loop)
c-----
C
C PROGRAM: VSTORE.FOR
C
C PURPOSE: STORE THE NODAL NUMBER FOR PLOT
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 8, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C-----
C
C ALGORITHM EXPLANATION: store the node number for draw polygone by
C counter clockwise order
c-----
c
c called by: vmesh.for -- volume mesh controler
c-----
c
c      include 'para.inc'
c      dimension x(0:maxdeg,maxid),y(0:maxdeg,maxid),
$          z(0:maxdeg,maxid),ideg(maxid),isurf(4,8)
c      common/data/x,y,z,ideg
c      dimension coor(4,maxnod),node(maxsid,1:maxloop)
c      common/auto/coor,node,last,unit
c      common/fin/ncon(0:10,maxid)
$      data isurf/1,8,2,5, 9,3,10,1, 10,8,12,6, 2,12,4,11,
$          5,11,7,9, 7,4,8,3/
c
c      save id
c
c      if(id.gt.maxid)then
c          write(*,*)'ERROR -- Increase <maxid> to larger than ',id
c          return
c      end if
c      do 10 k = 1,8
c          num = -1
c          id=id+1
c          do 20 l = 1,2
c          do 20 i = 1,node(1,isurf(l,k),loop)-1
c              num=num+1
c              ncon(num,id) = node(i+1,isurf(l,k),loop)
20      continue
c          do 30 l = 3,4
c          do 30 i = node(1,isurf(l,k),loop),2,-1
c              num=num+1
c              ncon(num,id) = node(i+1,isurf(l,k),loop)
30      continue
c          ideg(id) = num
c
c      store the coordinate into common block for drawing
c
c          do 40 i = 0,ideg(id)
c              x(i,id)=coor(1,ncon(i,id))
c              y(i,id)=coor(2,ncon(i,id))
c              z(i,id)=coor(3,ncon(i,id))
40      continue
10      continue
c
c          return
c          end
CNAME WRTMOV.FOR
c-----
c      subroutine wrtmov(cname)
c-----
C

```

```

C PROGRAM: WRTMOV.FOR
C
C PURPOSE: WRITE THE MOVIE SURFACE GEOMETRY FILE
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:                   CHECKED BY:          DATE:
C
C-----

```

```

C ALGORITHM EXPLANATION:
C
C-----

```

```

c called by: vmesh.for -- volume mesh controller
c
C-----

```

```

      include 'para.inc'
      common/auto/coor,node,last,unit
      common/fin/ncon(0:10,maxid)
      common/bsur/isurf,icont
      dimension coor(4,maxnod),node(maxsid,12,maxloop),jp(4*maxid)
      $           ,isurf(maxid/4)
      character cname=(*)
c
      write(*,*)' Write movie geometry surface file...'
      imov = 14
      open(imov,file=cname//'.geo',status='new')
c
c check to see if the node was in boundary surface or not
c
      np = 1
      ic = 0
      do 10 id = 1,icont
        if(ncon(0,id).gt.0)then
          do 20 k = 0,3
            ic = ic+1
            jp(ic) = ncon(k,isurf(id))
            if(mod(ic,4).eq.0)jp(ic) = -jp(ic)
20          continue
        end if
10      continue
      nedge = ic
      nj = last
      npt = icont
c
c write movie polygon file
c
      write(imov,150) np,nj,npt,nedge
      write(imov,150) np,npt
      write(imov,180) ((coor(i,j),i=1,3),j=1,nj)
      write(imov,150) (jp(i),i=1,nedge)
c
150      format(10i8)
180      format(1p8e12.5)
c
      return
      end
CNAME WRTNAV.FOR
c
c-----

```

```

      subroutine wrtnav(cname)
C-----
C
C PROGRAM: WRTNAV.FOR
C
C PURPOSE: write out the movie and navgraph solid geometry files
C
C WRITTEN BY: SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:                   CHECKED BY:          DATE:
C
C-----

```

```

C ALGORITHM EXPLANATION:
C
C-----

```

```

c called by:
c vmesh.for -- volume meshing controller
c
C-----

```

```

      include 'para.inc'

```

```

dimension jp(4*maxid), invt(4)
dimension coor(4,maxnod), node(maxsid,12,maxloop)
character cname*(*)
common/auto/coor,node,last,unit
common/fin/ncon(0:10,maxid)
data invt/0,3,2,1/

c
write(*,*) ' Write MOVIESTAR solid CADOS file'
imov = 15
inav = 18
open(imov,file=cname//'.mov',status='new')
open(inav,file=cname//'.nav',status='new')

c
np = 1
nj = last
do 10 id = 1,maxid
  if(ncon(0,id).le.0)then
    npt = id-1
    go to 20
  end if
10 continue
20 continue
ic = 0
do 30 id = 1,npt
  if(mod(id,6) .eq. 1)then
    do 40 j = 0,3
      ic = ic+1
      jp(ic) = ncon(j,id)
40 continue
  else if(mod(id,6) .eq. 0)then
    do 50 j = 0,3
      ic = ic + 1
      jp(ic) = ncon(invt(j+1),id)
50 continue
  end if
30 continue
nedge = ic
npt = npt/6

c
c write movie solid element file
c
write(imov,150) -np,nj,npt,nedge
write(imov,150) np,npt
write(imov,180) ((coor(i,j),i=1,3),j=1,nj)
write(imov,150) (jp(i),i=1,nedge)

c
c write navgraph neutral file
c
ntrnum = 0
nmpnum = 1
ngpnum = 0
do 60 j = 1,nj
  write(inav,180) j,ntrnum,(coor(i,j),i=1,3),const
60 continue
ibeg=1
do 70 j = 1,npt
  iend = ibeg+7
  write(inav,190) j,nmpnum,ngpnum,(jp(i),i=ibeg,iend)
  ibeg = iend+1
70 continue

c
150 format(10i8)
180 format(1p8e12.5)
180 format('NODE','2(I5,',')',3(G14.7,',')',G14.7,',')
190 format('ELEM','I5,',LS0,',I5,',',I5,8(',','I5)',',')
c
return
end

```

```

.....
-   The following files are for the   -
-   use of free-form deformation     -
.....

```

```

CNAME ALLDEF.FOR
c-----
      subroutine alldef
c-----
C
C PROGRAM: ALLDEF.FOR
C
C PURPOSE: DEFORMATION PROCESSOR
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C MODIFIED BY:                CHECKED BY:           DATE:
C-----
C ALGORITHM EXPLANATION: call 'deform' to perform the deformation
c-----
c
c called: deform.for -- deform the nodes
c-----
c
c called by: vmesh.for -- volume mesher
c-----
      include 'para.inc'
      dimension x(0:maxdeg,maxid),y(0:maxdeg,maxid),
      $          z(0:maxdeg,maxid),ideg(maxid)
      common/data/x,y,z,ideg
c
c deformed all the data
c
      write(*,*)' Freeform deformation running, please wait...'
      id = 1
      do while (ideg(id).gt.0)
        do 20 i = 0,ideg(id)
          xi = x(i,id)
          yi = y(i,id)
          zi = z(i,id)
          call deform(xi,yi,zi,xo,yo,zo)
          x(i,id) = xo
          y(i,id) = yo
          z(i,id) = zo
20        continue
          id = id + 1
        end do
      c
      return
      end
CNAME FFDEFO.FOR
c-----
      subroutine ffdefo
c-----
C
C PROGRAM: FFDEFO.FOR
C
C PURPOSE: DEFORMATION PROCESSOR
C
C WRITTEN BY:  SSUTA HSU           CHECKED BY:           DATE: OCT. 6, 1988
C MODIFIED BY:                CHECKED BY:           DATE:
C-----
C ALGORITHM EXPLANATION: call 'deform' to perform the deformation
c-----
c
c called: deform.for -- deform the nodes
c-----
c
c called by: vmesh.for -- volume mesher
c-----
      include 'para.inc'

```



```

      dimension xx(0:maxdeg,maxid),yy(0:maxdeg,maxid),
      $          zz(0:maxdeg,maxid),ideg(maxid)
      common/data/xx,yy,zz,ideg
      dimension coor(4,maxnod),node(maxsid,12,maxloop)
      common/auto/coor,node,last,unit
      common/fin/ncon(0:10,maxid)
c
c deformed after meshing the cube surfaces
c
c      write(*,*)' Freeform deformation running, please wait...'
      do 10 j = 1,last
          xi = coor(1,j)
          yi = coor(2,j)
          zi = coor(3,j)
          call deform(xi,yi,zi,xo,yo,zo)
          coor(1,j) = xo
          coor(2,j) = yo
          coor(3,j) = zo
10      continue
c
c store the coordinate into common block for drawing
c
      id = 1
      do while (ideg(id).ne.0)
          do 20 i = 0,ideg(id)
              xx(i,id)=coor(1,ncon(i,id))
              yy(i,id)=coor(2,ncon(i,id))
              zz(i,id)=coor(3,ncon(i,id))
20          continue
          id = id+1
      end do
c
      return
      end
CNAME FFEDIT.FOR
C-----
      subroutine ffedit
c-----
      common/defor/xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
      common/delta/dxyz(400)
c
      call editff(dxyz,l,m,n)
c
      return
      end
C-----
      subroutine editff(dxyz,l,m,n)
c-----
      dimension dxyz(4,0:l,0:m,0:n)
c
      write(*,100)l,m,n
      write(*,*)' Enter control point location number (I,J,K) '
      read(*,*)i,j,k
      write(*,101)dxyz(1,i,j,k),dxyz(2,i,j,k),dxyz(3,i,j,k),dxyz(4,i,j,k)
      write(*,*)' Enter new value (X,Y,Z,W) '
      read(*,*)dxyz(1,i,j,k),dxyz(2,i,j,k),dxyz(3,i,j,k),dxyz(4,i,j,k)
c
100      format(1x,'Control point dimension I =',i3,'J =',i3,'K =',i3)
101      format(1x,'Current control point value is ( ',4(f8.2,1x),')')
c
      return
      end
CNAME FRFORM.FOR
C-----
      subroutine frform(cname,ierr)
C-----
C
C PROGRAM: FRFORM.FOR
C
C PURPOSE: store data for free form deformation.
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 8, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C
C-----
C ALGORITHM EXPLANATION:
c
c-----

```

```

c
c  called: ffreed.for -- readin the deformation data
c
-----
c
c  called by: gtsoli.for -- solid mesh procesor
c
-----
c
c      common/defor/xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
c      common/delta/dxyz
c      dimension dxyz(400)
c      character cname*(*)
c      ierr = 0
c
c  read in the deformation control points data
c
c      open(2,file= cname//'.ffd',status='old',err=10)
c      read (2,*)xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
c      kall = 4*(l+1)*(m+1)*(n+1)
c      if(kall .gt. 400)then
c          write(*,*)'...Error in FRFORM increase dxyz() to',kall
c          stop
c      end if
c      call ffreed(l,m,n,dxyz)
c      go to 11
10      ierr=1
11      close(2)
c
c  perform free form deformation to all the nodes for further processing
c
c      call ffdefo
c
c      return
c      end
CNAME FFREAD.FOR
-----
c      subroutine ffreed(l,m,n,d)
-----
C
C  PROGRAM: FRFORM.FOR
C
C  PURPOSE: read in data for free form deformation.
C
C  WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C  MODIFIED BY:      CHECKED BY:      DATE:
C
-----
C
C  ALGORITHM EXPLANATION:
c
-----
c
c  called: none
c
-----
c
c  called by: frform.for -- deformation processor
c
-----
c
c      dimension d(4,0:l,0:m,0:n)
c
c      do 10 k = 0,n
c      do 10 j = 0,m
c          read (2,*)(d(1,i,j,k),d(2,i,j,k),d(3,i,j,k),
10      $          d(4,i,j,k),i=0,l)
c      continue
c
c      return
c      end
CNAME NODEFO.FOR
-----
c      subroutine nodefo
-----
C
C  PROGRAM: FRFORM.FOR
C
C  PURPOSE: deleted data of free-form deformation.
C
C  WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988

```

```

C   MODIFIED BY:                CHECKED BY:                DATE:
C
C-----
C   ALGORITHM EXPLANATION:
C-----
C   called: none
C-----
C   called by: main.for -- user's interface
C-----
C       common/defor/xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
C       xmin = 0.0
C       xmax = 0.0
C       ymin = 0.0
C       ymax = 0.0
C
C   deform all the nodes for further processing
C
C       call ffdefo
C
C       return
C       end
C
CNAME PDEF.FOR
C-----
C       subroutine pdef
C-----
C   PROGRAM: PDEF.FOR
C   PURPOSE: plot the free form deformation control points
C
C   WRITTEN BY:  SSUTA HSU          CHECKED BY:                DATE: OCT. 8, 1988
C   MODIFIED BY:                CHECKED BY:                DATE:
C-----
C   ALGORITHM EXPLANATION:
C-----
C   called: none
C-----
C   called by: deform.for -- deformation processor
C-----
C       common/defor/xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
C       common/delta/dxyz
C       dimension dxyz(400)
C
C       call pdefom(l,m,n,dxyz,xmin,ymin,zmin,xmax,ymax,zmax)
C
C       return
C       end
CNAME PDEFOM.FOR
C-----
C       subroutine pdefom(l,m,n,d,xmin,ymin,zmin,xmax,ymax,zmax)
C-----
C   PROGRAM: PDEFOM.FOR
C   PURPOSE: calculate the free form deformation control points convex hull
C
C   WRITTEN BY:  SSUTA HSU          CHECKED BY:                DATE: OCT. 8, 1988
C   MODIFIED BY:                CHECKED BY:                DATE:
C-----
C   ALGORITHM EXPLANATION:
C-----
C

```

```

c called: circ.for -- draw control point
c pipe.for -- draw convex hull
c
c-----
c
c called by: main.for -- user's interface
c
c-----
common/window/wxmin,wymin,wzmin,wxmax,wymax,wzmax
common/vport/vxmin,vymin,vzmin,vxmax,vymax,vzmax
dimension d(4,0:l,0:m,0:n),xc(0:200),yc(0:200),zc(0:200)
c
c for contral point circle
c
r=5.0*(wxmax-wxmin)/(vxmax-vxmin)
xdiff=(xmax-xmin)
ydiff=(ymax-ymin)
zdiff=(zmax-zmin)
if(xdiff.ne.0.0 .and. ydiff.ne.0 .and. zdiff.ne.0)then
do 30 k=0,n
do 30 j=0,m
do 30 i=0,l
ijk = i + j*(l+1) + k*(m+1)*(l+1)
xc(ijk) = (xmin + i*(xdiff)/l + d(1,i,j,k))
yc(ijk) = (ymin + j*(ydiff)/m + d(2,i,j,k))
zc(ijk) = (zmin + k*(zdiff)/n + d(3,i,j,k))
call circ(xc(ijk),yc(ijk),zc(ijk),r,5)
if(i.gt.0)call pipe(xc(ijk-1),yc(ijk-1),zc(ijk-1),
30 $ xc(ijk),yc(ijk),zc(ijk))
continue
do 40 k=0,n
do 40 j=1,m
do 40 i=0,l
ijk = i + j*(l+1) + k*(m+1)*(l+1)
lst = i + (j-1)*(l+1) + k*(m+1)*(l+1)
$ call pipe(xc(lst),yc(lst),zc(lst),
40 $ xc(ijk),yc(ijk),zc(ijk))
continue
do 50 k=1,n
do 50 j=0,m
do 50 i=0,l
ijk = i + j*(l+1) + k*(m+1)*(l+1)
lst = i + j*(l+1) + (k-1)*(m+1)*(l+1)
$ call pipe(xc(lst),yc(lst),zc(lst),
50 $ xc(ijk),yc(ijk),zc(ijk))
continue
end if
c
return
end

c-----
subroutine deform(x,y,z,xdu,ydu,zdu)
c-----
c
c this subroutine calculate the free form deformation.
c
c-----
C
C PROGRAM: FRFORM.FOR
C
C PURPOSE: calculate the free form deformation.
C
C WRITTEN BY: SSUTA HSU CHECKED BY: DATE: OCT. 6, 1988
C MODIFIED BY: CHECKED BY: DATE:
C
C-----
C
C ALGORITHM EXPLANATION: perform free-form deformation by Casteljau
c algorithm
c
c-----
c
c called: none
c
c-----
c
c called by: deform.for -- deformation processor
c

```

```

-----
c      common/defor/xmin,ymin,zmin,xmax,ymax,zmax,l,m,n
c      common/delta/dxyz
c      dimension dxyz(400)
c
c      $ call exdefm(x,y,z,xdu,ydu,zdu,xmin,ymin,zmin,
c          $      xmax,ymax,zmax,l,m,n,dxyz)
c
c      return
c      end
CNAME EXDEFM.FOR
-----
c      subroutine exdefm(x,y,z,xdu,ydu,zdu,xmin,ymin,zmin,
c          $      xmax,ymax,zmax,l,m,n,d)
-----
C
C PROGRAM: FRFORM.FOR
C
C PURPOSE: calculate the free form deformation.
C
C WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C MODIFIED BY:      CHECKED BY:      DATE:
C
-----
C ALGORITHM EXPLANATION: perform free-form deformation by Casteljau
c      algorithm
c
-----
c      called: none
c
-----
c      called by: deform.for -- deformation processor
c
-----
c      dimension d(4,0:l,0:m,0:n)
c      dimension xd(0:10),yd(0:10),zd(0:10),wd 0:10)
c      dimension xds(0:10),yds(0:10),zds(0:10) wds(0:10)
c      dimension xdt(0:10),ydt(0:10),zdt(0:10) wdt(0:10)
c
c      xdif=(xmax-xmin)
c      ydif=(ymax-ymin)
c      zdif=(zmax-zmin)
c
c if there is no deformation boundary then return the same values
c
c      if(xdif.eq.0.0 .or. ydif.eq.0.0 .or. zdif.eq.0.0)then
c          xdu=x
c          ydu=y
c          zdu=z
c      else
c          s=(x-xmin)/xdif
c          t=(y-ymin)/ydif
c          u=(z-zmin)/zdif
c
c if the point in the deformation boundary then deform it
c
c      if(s.ge.-0.005 .and.s.le.1.005 .and. t.ge.-0.005.and.t.le.1.005
c          $      .and. u.ge.-0.005 .and. u.le.1.005)then
c          do 9 k=0,n
c              do 10 j=0,m
c                  do 20 i=0,l
c                      wd(i) = d(4,i,j,k)
c                      xd(i) = (xmin + i*(xdif)/l + d(1,i,j,k))*d(4,i,j,k)
c                      yd(i) = (ymin + j*(ydif)/m + d(2,i,j,k))*d(4,i,j,k)
c                      zd(i) = (zmin + k*(zdif)/n + d(3,i,j,k))*d(4,i,j,k)
20                  continue
c                      wds(j)=point(wd,l,s)
c                      xds(j)=point(xd,l,s)
c                      yds(j)=point(yd,l,s)
c                      zds(j)=point(zd,l,s)
10                  continue
c                      wdt(k)=point(wds,m,t)
c                      xdt(k)=point(xds,m,t)
c                      ydt(k)=point(yds,m,t)
c                      zdt(k)=point(zds,m,t)
9                  continue
c                      wdu=point(wdt,n,u)

```

```

        xdu=point(xdt,n,u)
        ydu=point(ydt,n,u)
        zdu=point(zdt,n,u)
        xdu=xdu/wdu
        ydu=ydu/wdu
        zdu=zdu/wdu
c
c if the point out of the deformation boundary then no deformation
c
        else
            xdu=x
            ydu=y
            zdu=z
        end if
    end if
c
    return
end

CNAME CIRC.FOR
-----
      subroutine circ(xc,yc,zc,r,n)
-----
C
C PROGRAM: CIRC.FOR
C
C PURPOSE: DRAW CIRCLE ON THE GIVEN POINT
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C
-----
C ALGORITHM EXPLANATION:
C
-----
c parameter:
c xc      x-coordinate center of the circle
c yc      y-coordinate center of the circle
c r       radius of the circle
c n       segments needs to form a circle
-----
c called: pipe.for -- graphics line ploter
c
-----
c called by: pdefor.for -- plot control point for free-form deformation
c
-----
      dimension x(50),y(50)
      real r
c
      pi=acos(-1.0)
      if(n.le.3) n=3
      a=2.0*pi/n
      do 10 i=1,n
          x(i)=xc + r*cos(a*i)
          y(i)=yc + r*sin(a*i)
10      continue
c
      do 20 j=1,n-1
          call pipe (x(j),y(j),zc,x(j+1),y(j+1),zc)
20      continue
          call pipe (x(n),y(n),zc,x(1),y(1),zc)
c
      return
      end
CNAME PUSH.FOR
-----
      subroutine push(loop,ipt)
-----
C
C PROGRAM: PUSH.FOR
C
C PURPOSE: PUT THE DATA INTO STACK
C
C WRITTEN BY:  SSUTA HSU          CHECKED BY:          DATE: OCT. 6, 1988
C MODIFIED BY:          CHECKED BY:          DATE:
C

```

```

C-----
C
C  ALGORITHM EXPLANATION:
C-----
C
C  CALLED: NONE
C-----
C
C  CALLED BY: MOST OF THE ROUTINE
C-----
C
C      include 'para.inc'
C      common/stack/iarray(maxloop)
C
C      ipt = ipt + 1
C      iarray(ipt)=loop
C
C      return
C      end
C
CNAME POP.FOR
C-----
C      subroutine pop(loop,ipt)
C-----
C
C  PROGRAM: PUSH.FOR
C
C  PURPOSE: POP THE DATA OUT OF STACK
C
C  WRITTEN BY:  SSUTA HSU      CHECKED BY:      DATE: OCT. 6, 1988
C  MODIFIED BY:      CHECKED BY:      DATE:
C-----
C
C  ALGORITHM EXPLANATION:
C-----
C
C  CALLED: NONE
C-----
C
C  CALLED BY: MOST OF THE ROUTINE
C-----
C
C      include 'para.inc'
C      common/stack/iarray(maxloop)
C
C      loop=iarray(ipt)
C      ipt = ipt - 1
C
C      return
C      end

```

```
c
c maximum internal node number on each surface for 2-D case
c   parameter (intnod=1000)
c
c maximum allowable node number on each line
c   parameter (maxsid=100)
c
c maximum degree of a curve or line segments
c   parameter (maxdeg=20) for 2-D case
c   parameter (maxdeg=4)
c
c maximum number of polygons = maxele=8
c   parameter (maxid=12000)
c
c maximum number of elements
c   parameter (maxele=2000)
c
c maximum number of nodes in any mesh
c   parameter (maxnod=100000)
c
c maximum number of stack necessary for the loops
c   parameter (maxloop=20)
c
c set pi for processing 2-D case
c   parameter (pi=3.141593)
c   parameter (twopi=6.283185)
c   parameter (halfpi=1.570795)
c
c set zero for error measurement
c   parameter (zero=0.0001)
```


AUTOMATIC MESHING OF FREE-FORM DEFORMATION SOLIDS

Ssuta S. Hsu

Department of Civil Engineering

Ph.D. Degree, April 1989

ABSTRACT

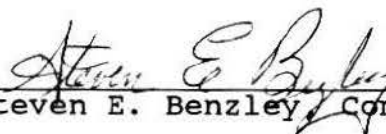
Development of computer models and subsequent finite element analysis, are important aspects of modern engineering design. In this process, the geometry creation and finite element analysis software are well developed; however, the process of discretizing a geometry into a proper finite element model is time consuming and tedious.

The work presented here uses the free-form deformation method to create smooth solid models, and invokes a solid subdivision and transition method to generate the hexahedron finite elements. The combination of these two techniques provides an automatic mesh generator that is easy to use, creates acceptable hexahedron elements for finite element analysis, and can model basically any complex shape.

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